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# Scheduling multimedia services in cloud computing environment

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## ABSTRACT

Currently, security is a critical factor for multimedia services running in the cloud computing environment. As an effective mechanism, trust can improve security level and mitigate attacks within cloud computing environments. Unfortunately, existing scheduling strategy for multimedia service in the cloud computing environment do not integrate trust mechanism when making scheduling decisions. In this paper, we propose a scheduling scheme for multimedia services in multi clouds. At first, a novel scheduling architecture is presented. Then, We build a trust model including both subjective trust and objective trust to evaluate the trust degree of multimedia service providers. By employing Bayesian theory, the subjective trust degree between multimedia service providers and users is obtained. According to the attributes of QoS, the objective trust degree of multimedia service providers is calculated. Finally, a scheduling algorithm integrating trust of entities is proposed by considering the deadline, cost and trust requirements of multimedia services. The scheduling algorithm heuristically hunts for reasonable resource allocations and satisfies the requirement of trust and meets deadlines for the multimedia services. Detailed simulated experiments demonstrate the effectiveness and feasibility of the proposed trust scheduling scheme.

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## KEYWORDS

Trust mechanism; cloud computing; multimedia services; QoS

## 1. Introduction

To meet the requirements of users, many of multimedia service providers (MSPs) provision various multimedia services such as, WebTV, VoD, mobile instant message, voice and video services delivered over Internet (Tan and Su 2011). Multimedia services demand power computational capacity and large storage space to process rich multimedia content. Recently, cloud computing has broad prospect to provide desired resources and provisioning quality of service (QoS) for multimedia services due to its features of on-demand supply and cost advantage (Tan and Su 2011; Zhu et al. 2011). Many of multimedia services in cloud computing environment spring up, such as video surveillance (Ren and Van Der Schaar 2013), cloud-based healthcare (Shini, Thomas, and Chithraranjan 2012), real-time access of medical images and video (An et al. 2014; Yang et al. 2015), cloud based photo and video sharing and so on.

Currently, security is a critical factor for multimedia services running in cloud computing environment (Lu, Varna, and Wu 2010; Wang et al. 2014b; Zhu and Lu 2009). Due to inherent distributed deployment of cloud computing environment, it is difficult to ensure each user can recognize a huge number of MSPs and select appropriate and trustworthy MSPs for trade.

Furthermore, uncertainty and unreliability of multimedia cloud pose threats to the operation of multimedia services. For example, multimedia cloud services sometimes may be offline unexpectedly due to power outage or other reasons. And the abuses of service resources by malicious users will greatly deplete MSP's resources and decrease its service capabilities (Zhou et al. 2011). On the other hand, some of selfish MSPs focus on augmenting their benefit and do not worry the benefit of users. In the context, it has contributed to the lack of confidence in the information discovered by the MSPs (Fard, Prodan, and Fahringer 2013). For instance, the MSPs might overstate the capacity of their resources in order to achieve more profits.

As an effective mechanism, trust has recently been suggested to enhance security and alleviate attacks in cloud computing environments (Dong and Dai 2008). Executing the jobs on the trusty MSPs can guarantee the security of sensitive data, reduce the services completion time and drop the ratio of failure execution. Unfortunately, most existing scheduling algorithms for multimedia services in cloud computing environment are merely devised to guarantee QoS demands while possibly ignoring the trust.

In this paper, we propose a novel scheduling and provisioning architecture for multimedia services in cloud computing environment. Within the proposed architecture, we build a trust model including both subjective trust and objective trust to evaluate the trust degree of a MSP. By employing Bayesian theory, the subjective trust degree between multimedia service providers and users is obtained. According to the attributes of QoS, the objective trust degree of a MSP is calculated. Considering QoS demands in terms of cost, complete time and the requirement of trust for multimedia services in cloud environment, we propose Trust Multimedia Scheduling Strategy (TMSS), a novel trust mechanism that counters the sensitive security of multimedia services. TMSS can heuristically find trusted resource allocations and achieve minimum completion time while meeting the constraints of cost and trust.

The contributions of this paper include the following aspects:

- (1) A novel scheduling and provisioning architecture for multimedia services running in cloud environment is designed.
- (2) TMSS is proposed for security-sensitive multimedia services with independence tasks in such cloud computing environment.
- (3) Detailed simulation experiments illustrate the validity and practicable of the proposed scheme.

The rest of the paper is organized as follows: [Section 2](#) summarizes the related work. In [Section 3](#), we describe the scheduling architecture for multimedia services in cloud computing environment. In [Section 4](#), we describe the formulation and solution of multimedia service scheduling and provisioning model, and present a novel trust model including subjective and objective trust that can compute the trust degree of MSPs and users. The proposed scheduling algorithm called TMSS is discussed in [Section 5](#). Simulated experiment results of the proposed scheme are analyzed in [Section 6](#). [Section 7](#) makes the conclusion and illustrates the future work.

## 2. Related works

In this section, we briefly summarize the related work about cost optimal scheduling, load balancing scheduling, resource failure and trust mechanism for multimedia services in cloud computing environment.

A number of scheduling algorithms for multimedia service have been developed to satisfy industrial requirements. Zhou et al.(2011) proposed a security-critical heuristic strategy to schedule multimedia applications on heterogeneous networks. Security requirements for multimedia applications is integrated into a real-time scheduling algorithm. Jyh-Cheng et al. (1998) investigated how to schedule of multimedia services in a low-power wireless ATM networks. The proposed algorithm is a

priority round robin with dynamic reservation update scheduling. Dong and Dai (2008) studied how to schedule non-real-time or real-time multimedia services in IEEE 802.16 e OFDM system. Considering the performances of efficiency, fairness and complexity, they proposed a new scheduling algorithm improving the efficiency of air interface resource at the expense of the short-time unfairness, but ensures the long-time fairness. Wu et al. (2010) built a model to realize the e-customer relationship building in buying multimedia product, suggest watermarks can help relationship building and retain e-customer to obtain multimedia products from legal channel.

Cloud resource providers provide on-demand resource provisioning with a pay-per-use model. Due to stronger resource demands of multimedia services, there is an upsurge of research interest in cost optimal for multimedia services in cloud computing environment. Li, Zhuo, and Shen (2013) presented a QoS based resource allocation method for handling multimedia tasks in the cloud. Considering completion time, cost and energy consumption, a QoS-based resource allocation method is proposed. Based on queuing model, Nan, Yifeng, and Ling (2011) proposed a novel optimized resource allocation schema for multimedia cloud. The response time minimization problem and resource cost minimization problem is formulated and solved in single service and multiple services, respectively. Zhipiao Liu et al. (Dong and Dai 2008) established a cloud service request model with SLA constraints, and then presented a cost-aware service request scheduling approach based on genetic algorithm. Wang et al. (2014c) proposed an accurate evaluation approach of Quality of Cloud Service (QoCS) in service-oriented cloud computing. The authors employed fuzzy synthetic decision to evaluate cloud service providers according to cloud users' preferences and then adopt cloud model to computing the uncertainty of cloud services based on monitored QoCS data. Wang et al. (2016) explored an alternative virtual machine placement approach to minimize energy consumption during the provision of data-intensive services with a global QoS guarantee in national cloud data centers. They used an improved particle swarm optimization algorithm to develop an optimal virtual machine placement approach involving a tradeoff between energy consumption and global QoS guarantee for data-intensive services. Song et al. (2016) presented a two-stage approach: media task management and cloud resource management for multimedia services in a cloud computing environment. The concept of allocation deadline is introduced in both of these approaches which makes the proposed solution unique from existing methods and benefits both multimedia service provider and cloud operator. For the media task management, a queuing based approach is proposed. In this approach, an online and efficient task management algorithm using Markov analysis and prediction was developed. For the resource management, an efficient heuristic algorithm was presented that can select and allocate task in an active and dynamic way through a dynamic controller. Lu et al. (2015) studied how to maximize an MSP's profit from provisioning multimedia services to geographically distributed users with a hybrid cloud. They first designed a service provisioning model to manage the resources in the hybrid cloud. Considering different time granularities for resource reservations, they leveraged the Lyapunov optimization technique to maximize the profit of MSP and proposed an online algorithm that can manage the hybrid cloud in the distributed manner. Briefly, some approaches described above assume that the resource information published by the MSPs is correct, which is doubtful to selfish MSPs.

Load balancing for multimedia services in cloud computing environment have been widely studied. Wen et al. (2011) investigated how to make each multimedia service task to obtain the required resources in the shortest time. They proposed an effective load balancing approach for cloud-based multimedia system. Its main advantage is fully considering the load of all servers and the network conditions, and thus achieving reasonable resource allocation and scheduling. Lin, Chin, and Deng (2014) took into account a more practical dynamic multi service scenario in which each server cluster only handles a specific type of multimedia task, and each client requests a different type of multimedia service at a different time. The scenario can be modeled as an integer linear programming problem. In (Lin et al. 2013), cloud computing based multi-tasking video transcoding is investigated and a parallelizing video transcoding framework is proposed based

on load balancing strategy. The multi-resources transcoding process in the cloud as a NP-hard problem was formulated. A heuristic algorithm called MLFT is proposed to minimize the entire finish time. In summary, some of aforementioned scheduling algorithms do not consider the dynamic change of resources in a real cloud computing.

Due to the increased functionality and complexity of the cloud computing system, resource failures are inevitable. Wang et al. (2014a) studied the performance of a cloud-based multimedia system using retrying for fault tolerance. The response time of the cloud-based multimedia system is modeled and the probability distribution of the response time is derived considering the effect imposed by retrying tasks in case of service failures happening. Javadi, Abawajy, and Buyya (2012) proposed a scalable hybrid Cloud infrastructure as well as resource provisioning policies to assure QoS targets of the users. The proposed policies take into account the workload model and the failure correlations to redirect users' requests to the appropriate Cloud providers. Though the works consider the resource failure and provision reliable multimedia services for user, they ignore security or trust requirement in scheduling schema and decrease the efficiency of media task scheduling.

At present, trust has been regarded as a valid mechanism to promote security within networked environments such as grid, cloud computing. Based on cognitive trust model, Wang et al. (2012) evaluated the trust of entities employing Bayesian theory, and integrated the trust model into existing DLS (dynamic level scheduling) algorithm to schedule the directed acyclic graph (DAG)-structure applications. In (Qu and Buyya 2014), Qu C. and Buyya R. proposed a system evaluating trust of clouds with regard to users' fuzzy QoS requirements and services' dynamic performances to promote service selection. Employing differential equations, Xiaoyong et al. (2011) designed a security scheduling schema to evaluate the trust degree of each node in heterogeneous system. It employed task priority rank to figure out security overhead of critical tasks. A security scheduling algorithm for DAG is presented to meet security requirement of workflow applications. Li et al. (2014) considered cloud user's trust and QoS and focused on a trustful workflow scheduling schema in multi cloud environment. The proposed schema describes the trust mechanism between cloud service providers and users. The trust mechanism can assure the successful execution of critical tasks. Two-level based scheduling mode maybe easily handles with the problem of workflow scheduling. Fard, Prodan, and Fahringer (2013) presented a pricing model and a truthful strategy to schedule single tasks. In order to minimize the completion time and monetary cost, this paper developed a trust scheduling algorithm for scientific workflows. Tan et al. (2014) proposed a trust service-oriented workflow scheduling algorithm. The scheduling algorithm uses Fuzzy Model to formulate the scheduling problem. It provided various strategies to help users balance different requirements in terms of cost, deadline and trust. Compared with our research, few of work focus on the problem of trust scheduling for multimedia services in cloud computing environment.

### 3. Multimedia services scheduling and provisioning architecture in cloud environment

In this section, we first present the multimedia services scheduling and provisioning architecture in cloud computing environment, and then we describe the scheduling process for multimedia services.

As depicted in Figure 1, the entities in cloud computing environment consist of MSPs, users and credible intermediate service center (CISC).

MSPs provision various multimedia services to end users for profit. MSPs offering the similar multimedia services are classified as one class. MSPs can resort to other MSPs for service composition in order to satisfy users' demands. Users access and obtain multimedia services provisioned by MSPs. Users are categorized according to their social contexts, preference for sharing services, and demand of available resource. CISC is the service register center of multimedia cloud environment. MSPs must register and publish a variety of multimedia services in

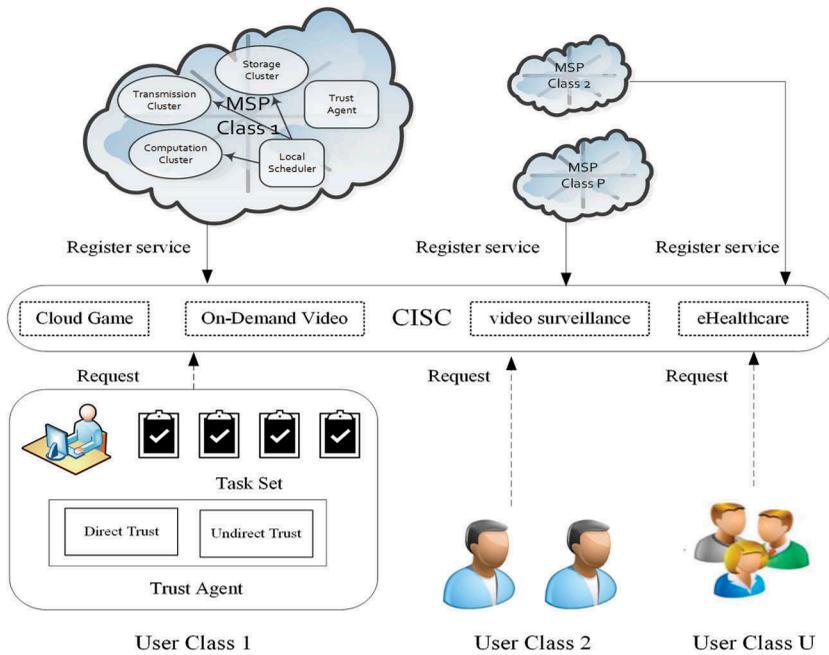


Figure 1. Multimedia services scheduling architecture in cloud computing environment.

CISC. In addition, CISC is the trustful third party intermediaries. Users resort CISC to find suitable MSPs according to the information provided by CISC in term of current users, resource and system load.

Users' trust agent evaluates the integrity of the MSPs according to the trust degree and informs user which is the most suitable one and execute the trade. Trust Agent uses Bayesian trust model, Eigen Trust (Kamvar, Schlosser, and Garcia-Molina 2003), differential equation (Xiaoyong et al. 2011) to compute subjective trust degree based on the information from direct experience or recommend reply. In addition, user can obtain the objective trust degree of a MSP from the feedback information of CISC. Finally, according to the subjective and objective trust degree of a MSP, users calculate the comprehensive trust value and determine whether they can trade with the special MSP. MSPs like users also judge the users' trust degree when processing services. MSPs can select trustworthy MSPs for service composition to meet personalized requirements.

The multimedia services scheduling process is detailed as followed.

When one user wants to trade with one MSP, he (she) submit its requests to the CISC, then CISC will help the user to select trustful MSPs. The user trade with MSPs and obtain the direct trust. The direct trust varies by MSPs and user's direct interaction, environment changing, and declines with time. The user first checks transaction history with the MSP. If the direct trust threshold is smaller than the effective number, he (she) take trust decision according to the direct trust information, else request recommendations from the same type of other users and feedback trust from CISC.

If a MSP processes service requests submitted by one user, he uses its trust agent to determine whether the user is trustful. If the requester is not included the untruthful list, the trust agent suggest the MSP to trade with the user. The trust relation among MSPs is same to the one that exists between users and MSPs.

Figure 2 describes the multimedia services scheduling process in cloud computing environment.

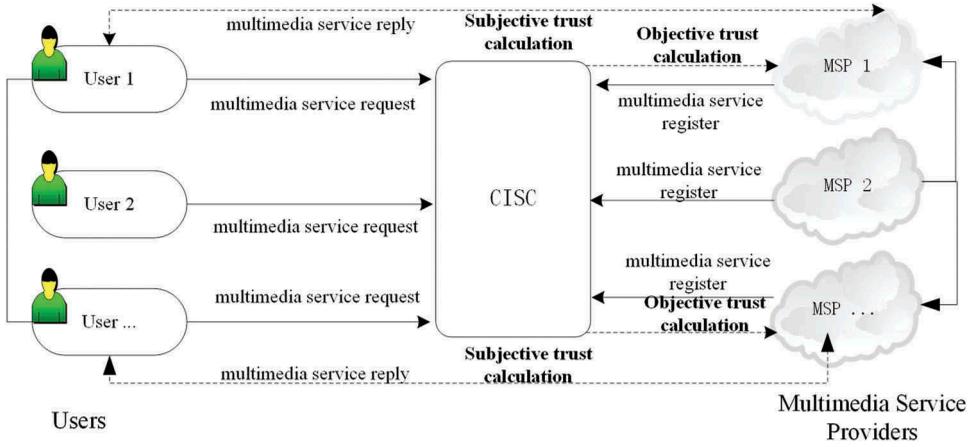


Figure 2. Multimedia services scheduling process in cloud computing environment.

#### 4. Formulation and solution of multimedia services scheduling and provisioning model

In this section, we present the formulation and solution of multimedia services scheduling problem of independent tasks across multi clouds. Key notations are listed in Table 1.

##### 1) Performance Model

There are some independent MSPs in cloud computing environment. Each of them is connected via high-speed network. Let  $P = \{p_1, p_2, \dots, p_m\}$  represent a set of MSPs, available to thousands of users at the same time. Each MSP  $p_j$  is characterized by a four tuple, i.e.  $p_j = \{id_j, name_j, Tr_j, resList_j\}$ .  $id$ ,  $name$ ,  $Tr$ ,  $resList$  represent the ID, Name, trust degree and resource set of MSPs, respectively. Each resource is presented by  $res_j = \{com_j, ep_j, st_j, tj_j, bw_j\}$ ,  $com_j$ ,  $ep_j$ ,  $st_j$ ,  $tj_j$ ,  $bw_j$  represent the processing capacity, computational price, storage price, traffic price, bandwidth of MSP $p_j$ , respectively.

In this paper, we focus on a class of parallel security-sensitive cloud-based multimedia services, such as video transform, medical image and video render. Each service request consists of a set of

Table 1. The notations.

Notation	Description
$P$	set of MSPs
$p_j$	the $j$ th MSP
$com_j$	the processing capacity of MSP $p_j$
$ep_j$	computational price of MSP $p_j$
$st_j$	storage price of MSP $p_j$
$tj_j$	traffic price of MSP $p_j$
$bw_j$	bandwidth of MSP $p_j$
$T$	set of multimedia tasks
$t_i$	the $i$ th multimedia task
$a$	the arrival time of a multimedia service request
$f$	the completion time of a multimedia service request
$D$	the deadline of executing the multimedia service
$M$	the cost constraint of executing the multimedia service
$length_i$	the task length of task $t_i$ ,
$datasize_i$	the data size of task $t_i$ ,
$tr_i$	the trust constraint of task $t_i$ ,
$est_j$	the remaining execution time of the running task on MSP $p_j$
$ft_{ij}$	the finish time of task $t_i$ on MSP $p_j$

independent tasks. Let a five tuple  $(T, a, f, D, M)$  represent such a service request. A set  $T = \{t_1, t_2, \dots, t_n\}$  represents dynamical independent tasks. Task  $t_i$  can be described by a collection of parameters, i.e.,  $t_i = \{length_i, datasize_i, tr_i\}$ , where  $length_i$ ,  $datasize_i$  and  $tr_i$  are the task length, data size and trust constraint of task  $t_i$ , respectively. In this paper, we only focus on inbound traffic, hence  $datasize_i$  is the input data of the task. It affects the data transfer time.  $a$  and  $f$  are the arrival time and completion time of a service.  $D$  represents the deadline of executing the multimedia service.  $M$  represents the cost constraint of executing the multimedia service.

Let  $est_j$  represents the remaining execution time of the running task on  $MSP p_j$ . It can be determined by the remaining workload size of all tasks that are running on  $p_j$ . Let  $dti$  represents the data transmission time. Data transfer occurs when a  $MSP p_j$  does not have data of task  $t_i$ . The data transmission time  $dti$  is defined as follows:

$$dti_{ij} = \frac{datasize_i}{bw_j}, \quad (1)$$

let  $eti_{ij}$  be the execution time of task  $t_i$  on  $MSP p_j$ , We get

$$eti_{ij} = \frac{length_i}{com_j} + dti_{ij}. \quad (2)$$

Let  $fti_{ij}$  be the finish time of task  $t_i$  on  $MSP p_j$ , then:

$$fti_{ij} = eti_{ij} + est_j, \quad (3)$$

Moreover, we use  $x_{ij}$ , a binary variable to reflect a mapping task  $t_i$  to  $MSP p_j$  in multimedia cloud. If  $x_{ij}$  is 1, task  $t_i$  is allocated to  $MSP p_j$ , and is 0 otherwise.

The finish time is also used to determine whether the task's timing constrain can be guaranteed, i.e.

$$x_{ij} = \begin{cases} 0 & \text{if } fti_{ij} > D, \\ 1 \text{ or } 0, & \text{if } fti_{ij} \leq D. \end{cases} \quad (4)$$

Generally, the service charge includes three items: computing, storage and data transfer. Hence, the cost of processing task  $t_i$  can be computed as followed:

$$CostF(t_i, p_j) = \frac{l_i}{c_j} \times ep_j + sd_i \times st_j + sd_i \times tj. \quad (5)$$

To meet deadline constraint of a service running in the multimedia cloud, we use a picking variable  $b_j$  to indicate whether a  $MSP$  is used, where  $b_j = 1$  means  $MSP j$  is used and  $b_j = 0$  otherwise. That is, the following must hold:

$$\max_{j=1-m} \left( \left( \sum_{i=1}^n eti_{ij} \times x_{ij} + est_j \right) \times b_j \right) \leq D. \quad (6)$$

The cost constraint can be defined as follows:

$$\sum_{i=1}^n \sum_{j=1}^m CostF(t_i, p_j) \times x_{ij} \leq M. \quad (7)$$

## 2) Trust Model

In general, trust is classified into two categories: subjective trust and objective trust. Subjective trust is related to the users' perception/belief of a  $MSP$  based on their own experiences and other

relevant knowledge (Li et al. 2014; Fan and Perros 2014). Objective trust is related to reliability, availability, integrity, and other parameters associated with a service (Li et al. 2014, 2015). In this paper, we build a trust model including both subjective trust and objective trust to evaluate the trust of a MSP.

### (1) Subjective trust evaluation

To meet the trust constraint of a multimedia service running in a MSP, we apply Bayesian theory to evaluate the subjective trust of MSPs. Bayesian theory is based on subjective probability, which is an opinion and degree of rational belief. It is the probability that someone believes something will happen. Subjective trust relationships usually include direct trust and indirect trust.

The direct trust is achieved by direct transaction history between cloud entities. The transaction history can be described by users' past experience with MSPs.

Assume A and B are two entities in the multi clouds, and their transaction results are expressed by binomial events. Let N denote times trades between them, S is successful trade times; F unsuccessful trade times; and  $Tr_{dt}$  denote the probability of successful trade at N + 1 times. Hence,

$$Tr_{dt} = E(\text{Beta}(Tr|S + 1, F + 1)) = \frac{S + 1}{S + F + 2}, \quad (8)$$

where  $0 < Tr_{dt} < 1$ , and  $F, S > 0$ .

If cloud entities trade with unfamiliar ones, they cloud adopt or combine indirect trust to help the trust decision. To compute the indirect trust, an entity makes use of its own experience of interactions, but also recommendation information from other entities. Using these recommendations, the user (MSP) computes the number of successful trade and the number of unsuccessful ones for each response.

Let  $q$  be the available recommendation number,  $n_s^m$  and  $n_f^m$  be the successful trade number, the unsuccessful trade number respectively. Those numbers are calculated according to the recommendations responded by an entity. Let  $n_s^r$ ,  $n_f^r$  be the total successful trade number and the unsuccessful trade number in the recommendations,  $\hat{\theta}_r$  as the probability of successful trade, we get

$$\hat{\theta}_r = \frac{n_s + \sum_{k=1}^q n_s^k + 1}{\sum_{k=1}^q n_s^k + \sum_{l=1}^q n_f^l + 2}, \quad (9)$$

Entities maybe not collect adequate experience from own history of trade to compute a direct trust. So, we employ a metric called Conf (level of confidence) (Deno and Sun 2008) to aid an entity determine whether the information collected is sufficient for computing the direct trust.

The Conf is achieved as follow:

$$\text{Conf} = 1 - \frac{(n_s + 1)(n_f + 1)}{(n_s + n_f + 2)^2(n_s + n_f + 3)}. \quad (10)$$

If Conf is low, it indicates that the confidence in the current trust information is not sufficient for computing the direct trust and the entity should collect more information.

Since the trust degree varies by time, we adopt a decay factor to represent the importance of the history information dropping with time. The trade of entities consists of time sequences. Let  $i$  represent a certain sequence,  $s_i$  and  $f_i$  is the successful trade number and

the unsuccessful trade number, the formula with the decay factor can be modeled (Wang et al. 2012):

$$s(n) = \sum_{i=1}^n s_i \cdot \eta^{(n-i)}, f(n) = \sum_{i=1}^n f_i \cdot \eta^{(n-i)}. \quad (11)$$

where  $s(n)$ ,  $f(n)$  is the number of successful and failure trade after  $n$ th sequence separately, and  $0 \leq \eta \leq 1$ . When  $\eta = 1$ , history interactions is abandoned;  $\eta = 0$ , the latest history record is considered. It should record the whole history interactions and can be solved by applying the following recursive algorithm (Wang et al. 2012).

$$s(i) = s(i-1) \cdot \eta + s_i, f(i) = f(i-1) \cdot \eta + f_i. \quad (12)$$

where  $s(i)$  and  $f(i)$  are the number of successful and failure trade at the  $i$ th sequence. The direct and recommendation trust at time  $i$  can be evaluated by formula (8) and (11) with  $s(i)$  and  $f(i)$  introduced into formula (12).

For two entities A and B in multimedia cloud, the successful trade probability between them is denoted by  $Tr_{su}$ . Direct interaction between A and B is denoted by  $Tr_{dt}$ . Indirect interaction is represented by  $Tr_{rt}$ . Thus, the subjective trust degree can be described as followed:

$$\begin{cases} Tr_{su} = wd \times Tr_{dt} + wr \times Tr_{rt} \\ wd + wr = 1, wd, wr \geq 0 \end{cases}. \quad (13)$$

where  $wd$ ,  $wr$  is the weight of direct trust  $Tr_{dt}$  and indirect trust  $Tr_{rt}$ .

## (2) Objective trust evaluation

From the perspective of objective trust, we mainly focus on reliability, availability, and integrity, three trusted factors of multimedia services.

Reliability is the probability of trusted service for a given duration, and we apply the average task success ratio to indicate this metric. It can be measured by

$$RE = \frac{renum_s}{renum_t} \quad (14)$$

where  $RE$  represents reliability;  $renum_s$  and  $renum_t$  denote the number of successful and total trades that occurred in a time interval, respectively. The average task success ratio is a statistical value for a special MSP. For example, within a given time, if a user asks for twenty services and seventeen are responded and finished. The average task success ratio is  $17/20 = 0.85$ .

Availability is the uptime percentage of multimedia services during a time interval, which can be measured by

$$AV = \frac{avnum_s}{avnum_t} \quad (15)$$

where  $0 \leq AV \leq 1$  represents availability;  $avnum_s$  and  $avnum_t$  denote the uptime and the total time of an trade period, respectively.

Integrity is the data accuracy of the result data, which can be measured by

$$IN = \frac{inum_s}{inum_t} \quad (16)$$

where  $inum_s$  is the number of service requests successfully completed by the MSP,  $inum_t$  is the number of service requests integrity completed by the MSP.

The objective trust of a MSP can be expressed by

$$Tr_{ob} = w_{RE}RE + w_{AV}AV + w_{IN}IN \quad (17)$$

where  $w_{RE}$ ,  $w_{AV}$ , and  $w_{IN}$  denote the weight of  $RE$ ,  $AV$  and  $IN$  while evaluating the objective trust of a MSP respectively,  $w_{RE} + w_{AV} + w_{IN} = 1$ .

### (3) Comprehensive trust calculation

According to the above description, the comprehensive trust calculation of a MSP can be given by

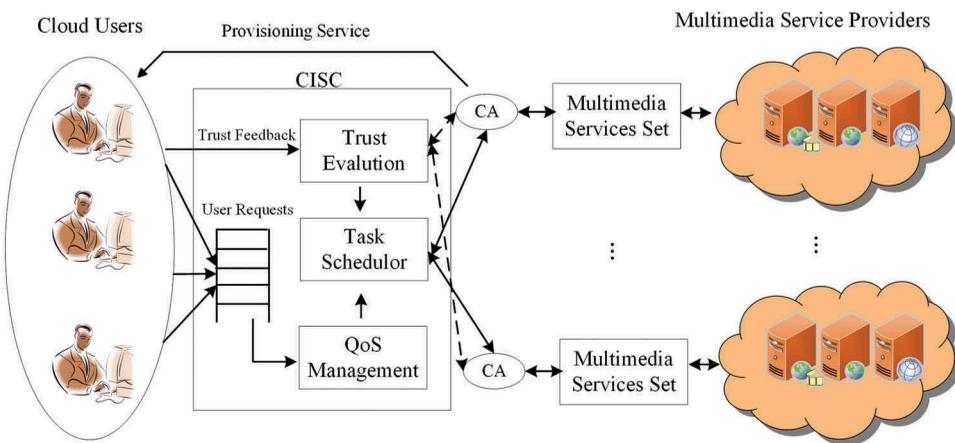
$$CT = \alpha Tr_{su} + \beta Tr_{ob} \quad (18)$$

where  $\alpha$  and  $\beta$  is the weight of  $Tr_{su}$ ,  $Tr_{ob}$  respectively,  $\alpha + \beta = 1$ .

### (4) Obtain the trust metrics

In general, each MSP must register its multimedia services through the CISC. The user negotiates with the CISC on QoS requirements. To acquire trust metrics, we designed computing agents (CAs), which is a software agent and responsible for collecting and computing objective trust metrics. These metrics can be obtained through calculation and statistics. Such metrics include the average task success ratio, the availability, the integrity of a service etc. To monitor multimedia services, we apply CAs to monitoring the average response time, and the average task success ratio etc. A CA captures data regarding the interaction process between a CISC and a MSP, and responds to requests for monitoring data from CISC. The CA collects monitoring information from the server side continuously, which involves all the performance parameters. The process of obtaining trust metrics is shown in [Figure 3](#).

### 3) Formalization of optimization problem



**Figure 3.** The process of obtaining trust metrics in multimedia cloud computing environment.

Mathematically, a typical model for multimedia applications in the multi clouds can be described as follows:

$$\text{Minimize : } Z = \max_{j=1-m} \left( \left( \sum_{i=1}^n et_{ij} \times x_{ij} + est_j \right) \times b_j \right), \quad (19)$$

$$\text{s.t. } \begin{cases} \max_{j=1-m} \left( \left( \sum_{i=1}^n et_{ij} \times x_{ij} + est_j \right) \times b_j \right) \leq D \\ \sum_{i=1}^n \sum_{j=1}^m \text{Cost}F(t_i, p_j) \times x_{ij} \leq M \\ tr_i \geq MSPTr_j \\ x_{ij}, b_j \in \{0, 1\} \end{cases} . \quad (20)$$

where  $Z$  is the object for minimization like time, cost, etc.  $X_d$  is the set of feasible solutions satisfying the QoS requirements. The aim of the trust-driven scheduling is to map every  $t_i$  onto some  $p_j$  in order to achieve minimum execution time while meeting the constraints in terms of budget and trust.

Because many of cloud-based multimedia services are not risk-free during the execution, it is essential for these MSPs to have security awareness. Therefore, we derive the risk probability to quantitatively analyze the risk rate for a multimedia service scheduled on MSPs. We suppose the risk rate is a function of trust levels. The distribution of the risk for any fixed time interval follows a Poisson probability distribution. The risk rate model is used for illustration purpose only. Thus, the service's risk rate on a specific MSP can be presented as follows:

$$pr(t_i, p_j) = \begin{cases} 0, & \text{if } t_i.Tr \leq p_j.trust \\ 1 - e^{-(t_i.Tr - p_j.trust)}, & \text{if } t_i.Tr > p_j.trust \end{cases} . \quad (21)$$

To find the optimal solution for a given problem that schedules  $n$  tasks on  $m$  MSPs with several restrictions, some heuristics must be employed to solve this problem in a real multimedia cloud environment since the number of MSPs and services is very large.

## 5. Multimedia service scheduling and provisioning algorithm in cloud computing environment

This section presents the proposed trust multimedia scheduling algorithm (TMSS) that allocates multimedia cloud resources to meet the demand in term of cost and trust of the submitted multimedia tasks.

The proposed TMSS algorithm is outlined as followed: In the pseudo code of TMSS optimization, a multimedia service enters the waiting queue according to its deadline, its tasks are sorted by their size (see line 1–2). Next, the trust degree of cloud entities is calculated (see line 3–9). Then, it adds all the available MSPs in the set  $trustP$  (see lines 10–13). The execution cost and time of each available MSP is calculated (see lines 16–17). After that, the tasks in multimedia service can be scheduled to fit MSP to meet the demand of cost and deadline (see lines 20–23).

## Algorithm TMSS

1. Sort media service according to the deadline, select the earliest deadline one in the arrival queue;
2. initialize the cost overhead of service, TotalCost = 0;
3. **for** MSP  $p_j$  of MSPs **do**
4.     calculating the direct trust degree of MSPs using Equation (8)
5.     calculating the indirect trust degree of MSPs using Equation (9)
6.     calculating the subjective trust degree of MSPs using Equation (13)
7.     calculating the objective trust degree of MSPs using Equation (17)
8.     calculating the comprehensive trust degree of MSPs using Equation (18)
9. **end for**
10. **for** task  $t_i$  of the service selected **do**
11.     **for** each MSP  $p_k$  of multimedia cloud
12.         computing  $pr(t_i, p_k)$  using Equation (20);
13.     **end for**
14.     search the trust MSP set  $P'$  with  $pr(t_i, p_j) < \theta$ ,  $\theta$  is the risk threshold;
15.     **While** (trustP not NULL) **do**
16.         **if**  $f_i \leq D$  and TotalCost < M **then**
17.             **for** each MSP  $p_j \in trustP$  **do**
18.                 calculate and save the cost  $c_{ij}$  for executing  $t_i$  on MSP  $p_j$
19.                 calculate and save  $et_{ij}$
20.             **end for**
21.             find the  $p_{best}$  that has the shortest completion time
22.             **if**  $p_{best}$  exists
23.                 assign task  $t_i$  to  $p_{best}$
24.                 add the cost of running  $t_i$  on  $p_{best}$  to TotalCost
25.             **end if**
26.         **end if**
27.     **end while**
28. **end for**

## 6. Performance evaluation

In this section, we first illustrate the simulated experiment environment, and then introduce the performance metrics. Next, we introduce the related algorithms. Finally, we detail the experiment results.

### 6.1. Simulated experiment environment

In this paper, we assume the multimedia services are provisioned through the type of virtual machines. Using the Java programming tools, we implemented a simulation environment, which is a cross-clouds platform including 10 MSPs, 100 virtual resources and some of users (SRs). Two assumptions are made: (1) a SR can trade with other SRs and offer the recommendations; (2) a SR can trade with the same type of MSPs.

The simulations run on a machine with an Intel Core 2 Duo CPU and 4 GB RAM. Each incoming task is set to have 200/400 MB input data size,  $10^6$ – $10^7$  Million Instructions and 0–1.0 trust requirement. The processor speed follows a uniform distribution in the interval [800, 1000] MIPS. The pricing model of resources based on Amazon prices, is summarized in Table 2. The failure rate follows a uniform distribution in the interval [0, 0.050].

**Table 2.** Resource pricing model.

Resource	Price(\$)
Computation	1–2.00(per hour)
Network(ln)	0.001(per GB)
Storage	0.001 (per GB per hour)

## 6.2. Performance metrics

To evaluate the trust integration performance, we use the following metrics:

**Completion time:** the duration time needed for the system to complete the execution of all tasks belonging to the same service.

**Cost:** the total cost in monetary units accumulated for an entire service execution.

**Success rate:** it is calculated for an service execution as a ratio between the number of successfully executed tasks and the number of the total scheduled tasks. MSPs that are error prone or those who fail, lead to smaller execution efficiency since the tasks they received must be rescheduled.

## 6.3. Related algorithms

To demonstrate the merit of TMSS, we compare it with three famous algorithms, i.e., MCT, Greedy, and TWFS (Tan et al. 2014). MCT: sorts MSPs by their computational capacity and selects the MSPs with the minimum completion time for the corresponding task. On the contrary, the Greedy approach sorts MSPs by their cost and selects the MSP with the minimum cost. TWFS achieve minimum execution cost and time while meeting the constraints of trust. The key idea of the TWFS algorithm is to find the optimum solution with the deadline constraint by adjusting the weights of time and cost.

## 6.4. Simulation experiments

### 6.4.1. Effects of task number

The experiment is to test the MCT, TWFS, Greedy, TMSS algorithms with different task number. The number of tasks varies from 100 to 500 with an increment of 100. This time we set the every job deadline to be a reasonable integer. The results of execution time are figured in Figure 4. Figure 4 show that the TMSS can averagely perform 89.3 % less time than the MCT in terms of the completion time. The completion time of TMSS is shortest among the four algorithms. Obviously, the TMSS performs significantly better than other algorithms in latency comparisons.

The experiment result shows that the proposed scheduling integrating trust mechanism can shorten the completion time of services compared to the MCT and Greedy algorithm, which do not load trust mechanisms. MCT shorten the services execution time by assigning them to the MSPs

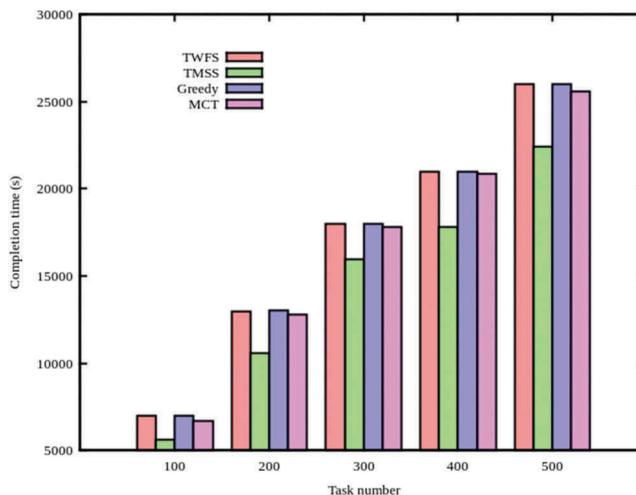


Figure 4. Comparison of completion time varying task number.

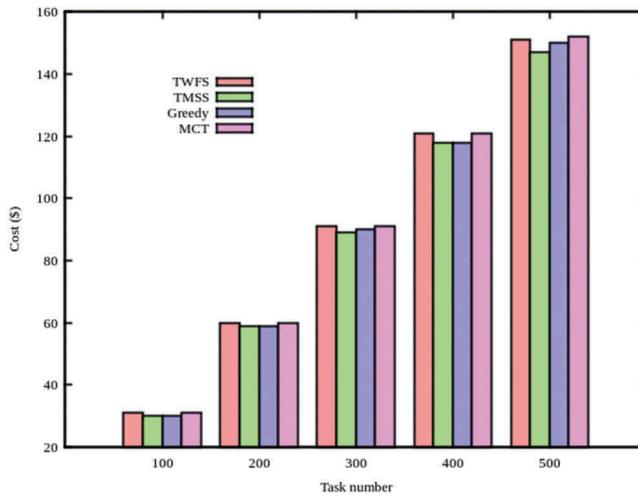


Figure 5. Comparison of cost varying task number.

which has the shortest completion time. But they ignore resources integrity and sensitive tasks may be sent to untruthful MSPs. Hence the execution is delayed and necessary to be re-scheduled. The services completion time is postponed. While our schema considers the trust degree of the candidate MSPs in the scheduling process and reduces the probability of failure scheduling. The whole service execution time is shortened. With task number rising, trust-driven scheduling schema describes more advantage.

Figure 5 shows the result of the cost for executing jobs of 100, 200, 300, 400 and 500 independent tasks by TMSS, MCT, TWFS and Greedy. From Figure 5, we can judge that the execution cost of TMSS algorithm is close to that of Greedy's which is the cheapest. It indicates that the TMSS algorithm outperforms other two algorithms. In this experiment, the TMSS spends about 20.7 % less money than the MCT in average. While the cost of MCT, Greedy and TWFS is close to each other. The reason is that the TMSS considers the trust constraint as the most important constraint, while the MCT focuses on lowering the execution time and spends more money.

#### 6.4.2. Effect of varying deadline

In this experiment, we investigate the effect of deadline on the performance of the trustful scheduling algorithm. Deadline is computed according to the minimum execution time produced by MCT. We change deadline from 10000 seconds to 14000 seconds for 200 tasks. The deadline for 500 tasks range from 26000 seconds to 30000 seconds. Figure 6 depicts the completion time result of our experiment. In Figure 6(a), it is observed that for task number 200, an increase in deadline also gives rise to the completion time. We can see that the completion of TMSS is constant while deadline is large than 11000 seconds. This is TMSS takes the shortest time to execute tasks. From Figure 6(b), the same takes place.

Figure 7(a) shows the results of scheduling several jobs consisting of 200 tasks totally. Given a very strict deadline such as 10000 seconds, all schedulers present a similar, low satisfaction rate because of the limitation of available resources. With a loose deadline such as 12000 seconds, the TMSS can achieve better success ratio than the other three schedulers. Figure 7(b) shows that, for scheduling several jobs consisting of totally 500 tasks, the TMSS can reach 100 % success ratio for a deadline of 27000 seconds while the others cannot even reach 87 % success ratio for a deadline of 27000 seconds. The major reason why the TMSS is better in

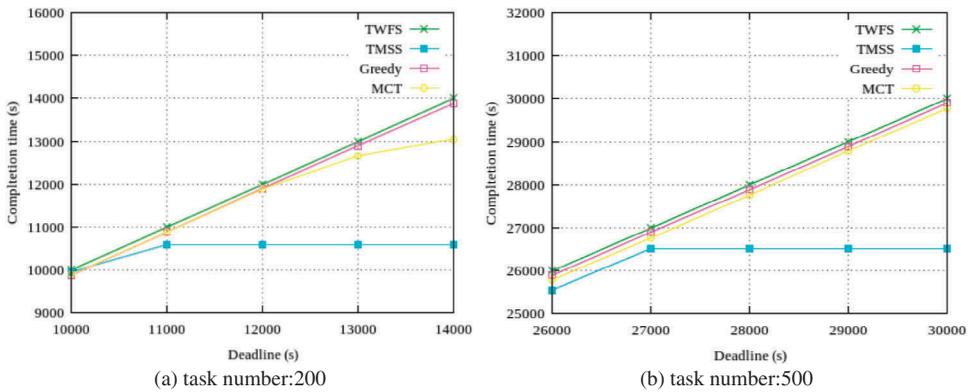


Figure 6. Comparison of completion time varying deadline.

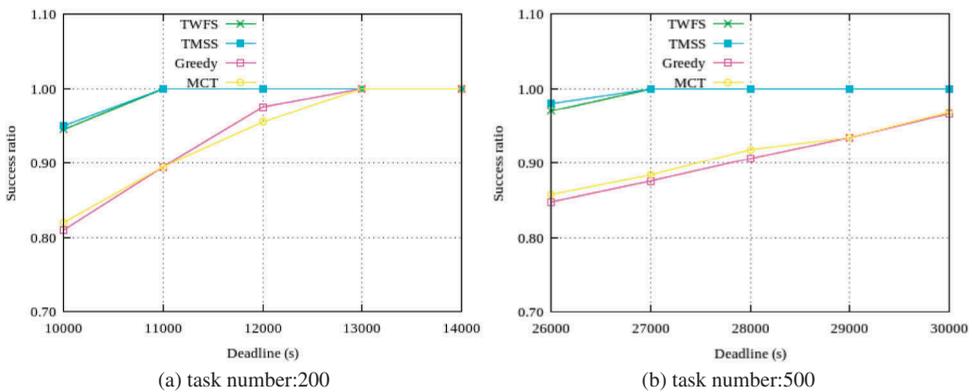


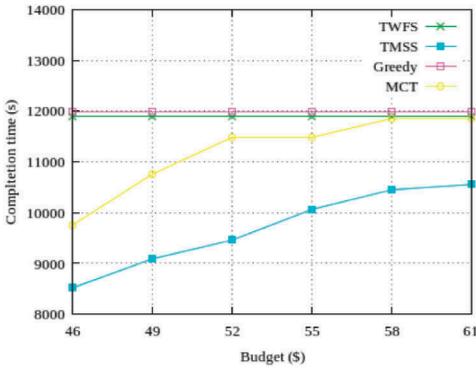
Figure 7. Comparison of success ratio varying deadline.

success ratio is that the TMSS tries to send tasks to trust MSP and assure the normal execution of tasks.

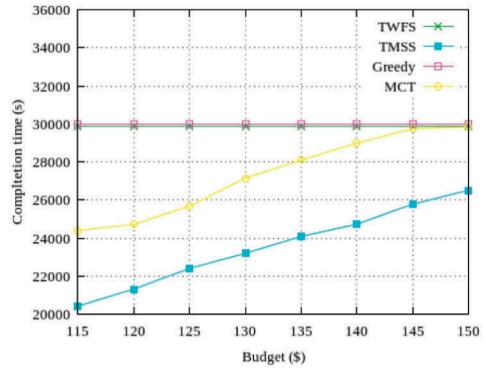
#### 6.4.3. Effect of varying budget

In this experiment, the budget is computed according to the minimized cost produced by Greedy. For 200 tasks, the budget varies from 46\$ to 61\$. The result of completion time is presented by Figure 8(a). For 500 tasks, the budget ranges from 115\$ to 150\$. The result of completion time is presented by Figure 8(b). From Figure 8(a,b), we can see that the completion time of tasks produced by TWFS is closer to Greedy. This is because both scheduling algorithms try to schedule tasks to cheapest or maximized utility resource. The completion time produced by TMSS is shortest among the four algorithms. TMSS spends less than 20 percent than that of MCT. The experiment result shows that the new scheduling model including trust requirement can shorten the services completion time while compared to the MCT and Greedy which do not provide trust mechanisms.

From Figure 9, we can see that: in multimedia cloud environment with some untruthful MSPs, scheduling model including trust mechanism can assure higher success rate because it efficiently eliminates untruthful MSPs according to passed trade information and other trust users' recommendation. The ratio of successful execution of TMSS is much higher than TWFS's and MCT's. This demonstrates that the trustful scheduling algorithm can guarantee the successful execution of tasks.

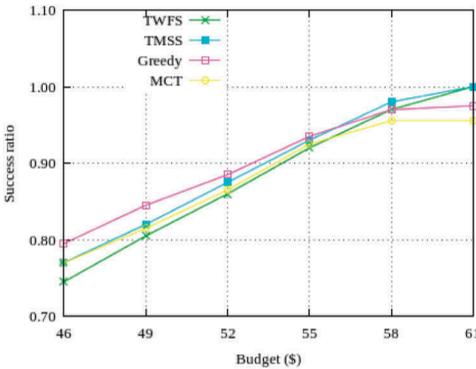


(a) task number:200

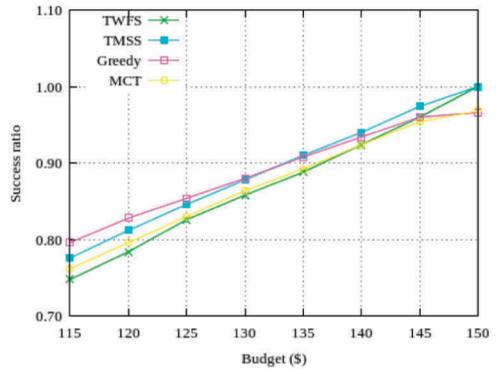


(b) task number:500

Figure 8. Comparison of completion time varying budget.



(a) task number:200



(b) task number:500

Figure 9. Comparison of success ratio varying budget.

## 7. Conclusion and future work

In this paper, we develop and evaluate a novel scheduling and provision scheme for multimedia services in cloud computing environment. Trust mechanism can dynamically assess the trust degree of multimedia service provider by Bayesian method and integrated the trust metric into scheduling, which ensures the execution success rate of security-critical multimedia tasks. The proposed trustful scheduling scheme can heuristically seek approximate optimal resource allocations and meet the requirement of trust and deadline for multimedia services. The performance of TMSS algorithm is compared with three well-known scheduling algorithms for multimedia service: MCT, TWFS and Greedy algorithms. Detailed experiment results demonstrate the feasibility and efficiency of the proposed scheme.

As part of our future work, we plan to extend the TMSS algorithm for dependent multimedia services in cloud computing environments.

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