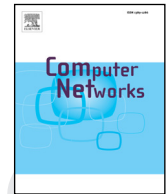




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A two-stage game theoretical approach for interference mitigation in Body-to-Body Networks[☆]

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ABSTRACT

In this paper, we identify and exploit opportunities for cooperation between a group of mobile Wireless Body Area Networks (WBANs), forming a Body-to-Body Network (BBN), through inter-body interference detection and subsequent mitigation. Thus, we consider a dynamic system composed of several BBNs and we analyze the joint mutual and cross-technology interference problem due to the utilization of a limited number of channels by different transmission technologies (i.e., ZigBee and WiFi) sharing the same radio spectrum. To this end, we propose a game theoretical approach to address the problem of Socially-aware Interference Mitigation (SIM) in BBNs, where WBANs are “social” and interact with each other. Our approach considers a two-stage channel allocation scheme: a BBN-stage for inter-WBANs’ communications and a WBAN-stage for intra-WBAN communications. We demonstrate that the proposed BBN-stage and WBAN-stage games admit exact potential functions, and we develop a Best-Response (BR-SIM) algorithm that converges to Nash Equilibrium points. A second algorithm, named Sub-Optimal Randomized Trials (SORT-SIM), is then proposed and compared to BR-SIM in terms of efficiency and computation time. We further compare the BR-SIM and SORT-SIM algorithms to two power control algorithms in terms of signal-to-interference ratio and aggregate interference, and show that they outperform the power control schemes in several cases. Numerical results, obtained in several realistic mobile scenarios, show that the proposed schemes are indeed efficient in optimizing the channel allocation in medium-to-large-scale BBNs.

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1. Introduction

Body-to-Body Networks (BBNs) have recently emerged as a promising solution for monitoring the people behavior and their interactions with the surrounding environment [2].

The BBN consists of several WBANs, which in turn are composed of sensor nodes that are usually placed in the

clothes, on the body or under the skin [3]. These sensors collect information about the person and send it to the sink (i.e., a Mobile Terminal (MT) or a PDA), in order to be processed or relayed to other networks (Fig. 1).

BBNs are widely adopted in several mission-critical scenarios: (i) rescue teams in a disaster area, (ii) groups of soldiers on the battlefield [4], and (iii) patients in a health-care center, whose Wireless Body Area Networks (WBANs) interact with each other. Yet, the BBN can be implemented in both medical and non-medical applications. Indeed, BBNs represent the novel trend for future, ubiquitous healthcare systems, in which the remote monitoring of patients carrying bodyworn sensors and relaying each others physiological data up to the medical center, could greatly

[☆] Very preliminary results of this work have been presented in [1].

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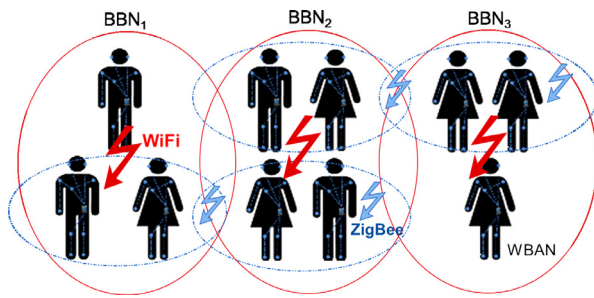


Fig. 1. Three-BBN interfering scenarios: each BBN is composed of several WBANs which use different transmission technologies (i.e., ZigBee and WiFi) sharing the same radio spectrum.

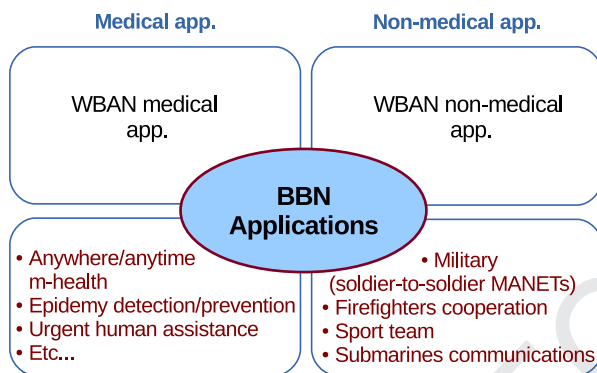


Fig. 2. Application area extensions from WBAN to BBN.

22 reduce the current strain on health budgets and make the
 23 Government's vision of ubiquitous healthcare for distant
 24 patients a reality. For example, when a patient is at home
 25 or far from the medical center, and feels a sudden trouble,
 26 she will be able to broadcast a distress call and bring
 27 out an urgent human assistance from his neighborhood.
 28 Hence, the sensors could be embedded into mobile hand-
 29 sets, portable electronic devices, cars, and clothing. Due
 30 to low-power Body-to-Body Networks, people would no
 31 longer need to be in the range of a cellular tower to make
 32 a call or transmit data. Fig. 2 sorts the different BBN
 33 applications into medical and non-medical classes, and lists
 34 the new intended applications by the deployment of BBN
 35 networks.

36 Due to the scarce wireless resources, many exist-
 37 ing wireless technologies, like IEEE 802.11 (WiFi), IEEE
 38 802.15.1 (Bluetooth) and IEEE 802.15.4 (ZigBee), are forced
 39 to share the same unlicensed 2.4 GHz Industrial, Scien-
 40 tific and Medical (ISM) band. Hence, mutual as well as
 41 cross-technology interference may occur between these
 42 technologies.

43 Indeed, the interference issue is already handled by the
 44 Bluetooth Low Energy (BLE) standard [5], which defines
 45 three channels as advertising channels, used for device dis-
 46 covery and connection establishment, and have been as-
 47 signed center frequencies that minimize overlapping with
 48 IEEE 802.11 channels 1, 6 and 11, which are commonly
 49 used in several countries. Then, an *adaptive frequency hop-*
 50 *ping* mechanism is used on top of the 37 data channels

51 in order to face interference and wireless propagation issues,
 52 such as fading and multipath. This mechanism selects
 53 one of the 37 available data channels for communication
 54 during a given time interval, so as to avoid interference
 55 with neighboring wireless links. Furthermore, a number of
 56 previous works enhanced the existing frequency hopping
 57 mechanism and implemented further schemes, such as the
 58 OverLap Avoidance (OLA) proposed in [6].

59 Coexistence and interference mitigation between
 60 WBANs are also considered by the IEEE 802.15.6 standard.
 61 Three mechanisms are defined: *beacon shifting*, *channel*
 62 *hopping* and *active superframe interleaving* [7]. Yet, our
 63 choice for ZigBee aims at effectively and theoretically tack-
 64 ling the cross-technology interference problem between
 65 WiFi (802.11) and ZigBee (802.15.4) technologies.

66 Since WiFi transmission power can be 10 to 100 times
 67 higher than that of ZigBee, ZigBee communication links
 68 can suffer significant performance degradation in terms
 69 of data reliability and throughput. In addition to the
 70 previously mentioned challenging issues, the mobility of
 71 WBANs in their surrounding environment and their inter-
 72 actions with each other make the interference mitigation
 73 in Body-to-Body networks a very interesting and manda-
 74 tory problem to address. This is indeed the main focus of
 75 our work.

76 In this paper we consider a multi-BBN scenario (an ex-
 77 ample scenario, with 3 BBNs, is illustrated in Fig. 1),
 78 composed of a set of WBANs that share the same ISM band,
 79 and we address the mutual and cross-technology inter-
 80 ference mitigation problem introducing a new game theo-
 81 retical approach. The proposed approach consists of two
 82 nested games. The first game aims to allocate WiFi chan-
 83 nels for inter-WBANs' wireless communications. Specifi-
 84 cally, special players (which are called "delegates" or "lead-
 85 ers") decide the allocation of the needed WiFi channels for
 86 themselves and the underlying subnetworks by maximiz-
 87 ing an utility function, which is a function of mutual and
 88 cross-technology Signal-to-Interference Ratio (SIR) metric.
 89 The second proposed game is a WBAN-stage SIM game
 90 that allows players (or WBANs) to choose the needed Zig-
 91 Bee channels for intra-WBAN communications, taking into
 92 account the allocations performed by the BBN-stage SIM
 93 game.

94 The main contributions of our work are the following:

- We propose a novel game theoretical approach for mutual and cross-technology interference mitigation in BBNs.
- We provide a detailed expression of the *Signal-to-Interference Ratio* to define players' payoff functions, capturing all main interference components, namely the co-channel, the mutual, and the cross-technology interference.
- We demonstrate that our games admit at least one pure strategy Nash Equilibrium (NE) since they are exactly potential, and we develop best response algorithms (BR-SIM) to compute the channel allocations, which converge fast to NE solutions.
- We propose a second algorithm, called Sub-Optimal Randomized Trials (SORT-SIM), that trades-off between efficient channel allocation process and short

111 computation time, and guarantees a sub-optimal solu-
112 tion to the SIM problem.

113 • We perform a thorough performance analysis of the
114 BBN- and WBAN-stage SIM games under different sys-
115 tem parameters, and compare the two proposed al-
116 gorithms, i.e., BR-SIM and SORT-SIM to a distributed
117 power control and a relay-assisted power control al-
118 gorithm. Numerical results show that the proposed
119 schemes are indeed efficient in optimizing the chan-
120 nel allocations in medium-to-large-scale realistic mo-
121 bile BBN scenarios.

122 The paper is structured as follows: Section 2 dis-
123 cusses related work. Section 3 presents the BBN sys-
124 tem model, including the communication and the interfer-
125 ence model. Section 4 details the two-stage Socially-aware
126 Interference Mitigation (SIM) game theoretical approach.
127 Section 5 presents the Best-Response Algorithm (BR-SIM),
128 while Section 6 handles the sub-optimal solution (SORT-
129 SIM) for the SIM problem. Section 7 analyzes numerical
130 results for the proposed solutions in several BBN scenar-
131 ios. Finally, Section 8 concludes this paper.

132 2. Related work

133 In this section, we discuss the most relevant works that
134 deal with the problem of interference mitigation between
135 different technologies (i.e., Bluetooth, ZigBee, WiFi) that
136 share the same frequency spectrum.

137 Whilst a number of previous interference-aware studies
138 have been based upon power considerations [8,9], others
139 have chosen different alternatives [10,11] to deal with this
140 substantial problem which is challenging in WBAN design,
141 and raising even more with the emergence of BBNs.

142 In [8] the authors propose a distributed power control
143 algorithm which converges to the Nash Equilibrium, rep-
144 resenting the best tradeoff between energy and network
145 utility. No transmissions are envisaged among WBANs in
146 [8]; a transmission is either from a WBAN node to its gate-
147 way or vice versa, neither access technology assumption
148 is made, it is rather assumed that only mutual interfer-
149 ence could happen. However, in a BBN context where
150 WBANs communicate with each other, it is mandatory to
151 consider transmissions among WBANs' gateways and thus
152 investigate cross-interference scenarios where different
153 wireless technologies could be used for intra-WBAN and
154 inter-WBANs transmissions scenarios.

155 While most power control models provide interference-
156 aware schemes over power adaptation, authors of [9] opti-
157 mized a transmission scheme given a constant power. They
158 formulated an interference-aware channel access game to
159 deal with the competitive channel usage by different wire-
160 less technologies sharing the ISM band, in both static
161 and dynamic scenarios. Using Game Theory, authors in [9]
162 stated that a decentralized approach is resilient to users'
163 deviation and ensures the robustness of the network, com-
164 pared to a centralized approach where the system cannot
165 be easily protected from a selfish deviation to in-
166 crease, unilaterally, one's throughput. Alike our BBN model,
167 this game considers nodes concurrently transmitting in
168 nearby clusters, incorporating the Signal-to-Interference-

169 plus-Noise Ratio (SINR) model as wireless communication
170 metric. Nonetheless, the game focuses on the channel ac-
171 cess problem under *inter-cluster* interference from nearby
172 APs using the same wireless technology, while the key ad-
173 vantage of our work is to consider both *mutual* and *cross-*
174 *technology* channel interference problems.

175 Game Theory is applied in such distributed problems,
176 such as in [10], where the multi-channel usage problem in
177 Wireless Sensor and Actuator Networks (WSANs) is mode-
178 led as a channel allocation game with the total interfer-
179 ence of the whole network as the social objective to min-
180 imize. In WSANs, communication and control are highly
181 integrated, even though each node (a sensor, actuator or
182 control unit) is equipped exclusively with one simple half-
183 duplex radio transceiver. However, the major difference
184 with our network model is that BBNs are randomly dis-
185 tributed networks where underlying WBANs are mobile
186 and equipped with two radio antennas to ensure on-body
187 and off-body communications. Yet, WBANs may randomly
188 overlap with each other, which makes BBN a highly dy-
189 namic system over time and space, compared to WSNs,
190 apart from the human body environment challenge re-
191 lated to WBANs. Yet, further constraints are to be con-
192 sidered to design an effective channel allocation scheme
193 for BBNs.

194 On the other hand, the main idea in [11] is that using
195 only power control to combat this interference might not
196 be efficient; it could even lead to situations with higher
197 levels of interference in the system. Therefore, the work
198 in [11] proposes several interference mitigation schemes
199 such as adaptive modulation as well as adaptive data
200 rate and adaptive duty cycle. Interference Mitigation Fac-
201 tor is introduced as a metric to quantify the effectiveness
202 of the proposed schemes. Based on SINR measurements,
203 these schemes are likely suitable for small-scale WBANs
204 where SINR is function of the transmit power, such as
205 in [8] which uses the SINR metric as a utility function
206 to model the interference problem between neighboring
207 WBANs considering a power control game. In fact, in [8]
208 the network topology is static and no actual communica-
209 tions among WBANs are considered. However, in [12], an
210 experimental study proved the importance of the impact
211 of human body shadowing in off-body communications.
212 Yet, for relatively complex BBNs, SINR is also highly de-
213 pendent on outdoor conditions and human body effects,
214 and the aforementioned schemes would no longer be ef-
215 ficient, or they should be extended taking into account ad-
216 ditional physiological, physical, and environmental param-
217 eters. Particularly, in dynamic scenarios, when the SINR
218 is varying due to the fast topology changes with neighboring
219 WBANs movements, relying only on the transmit power in
220 order to keep the desired link quality might not be effec-
221 tive. Indeed, in a BBN scenario with high transmit power
222 from other coexisting wireless networks/WBANs, the inter-
223 ference is significant and the desired link quality cannot
224 be achieved unless considering the surrounding conditions
225 (interference) and the wireless channel characteristics in
226 terms of shadowing, fading, etc., which can be incorpo-
227 rated into the channel gain parameters of the SINR.

228 Besides, several works investigated the interference
229 mitigation problem with detailed specifications of wireless

technologies, especially WiFi, ZigBee, and Bluetooth, which are very popular in the WBAN industry. For example, authors in [13] proposed an approach that accurately characterizes the *white space* in WiFi traffic and develop a ZigBee frame control protocol called WISE, which can predict the length of white space in WiFi traffic and achieve desired trade-offs between link throughput and delivery ratio. The empirical study of ZigBee and WiFi coexistence provided by authors in [13] is useful to understand and model the cross-technology problem. Nevertheless, the WiFi-WiFi and ZigBee-ZigBee mutual interference problems still need to be carefully investigated, especially when coupled with mobility, topology changes and other features related to the complexity of BBN networks, which require more intelligent functions at the WBAN coordinator's (MT) level, in order to ensure an effective channel allocation scheme for BBNs. Further studies [14–16] have dealt with the solutions that enable ZigBee links to achieve guaranteed performance in the presence of heavy WiFi interference, but almost all of them propose approaches that assume having already established the ZigBee and WiFi links, and try to implement mechanisms to mitigate the interference between them.

In [17], the authors provided an interesting study that explores the possibility of exploiting Partially Overlapped Channels (POCs) by introducing a game theoretic distributed Channel Assignment (CA) algorithm in Wireless Mesh Networks (WMNs). The proposed CA algorithm aims at increasing the number of simultaneous transmissions in the network while avoiding signal interference among multi-radio nodes. A Cooperative Channel Assignment Game (CoCAG) is implemented, where information is exchanged with neighboring nodes. In fact, by considering neighboring information, nodes can track the instantaneous neighbors' strategies when assigning channels to themselves, which can help in guaranteeing a fair sharing of the frequency band. The major contribution of [17] is that it addresses four different types of interference and their influence on the network capacity: Co-channel Interference, Orthogonal Channels, Adjacent Channel Interference and Self Interference. Nonetheless, one key feature of the WMN is the backbone network composed of Mesh Routers that are usually static and have no constraints on energy consumption, which is not the case for WBANs. Moreover, only IEEE 802.11g was used as wireless technology in [17], and as a consequence no cross-technology scenarios were considered.

Again, in order to cope with the interference issue in WBANs, authors in [18] implemented an intelligent power control game which allows WBANs to improve their performance by learning from history. The proposed power controller implements a genetic algorithm (GA) which enables WBANs to learn from experience and select their power strategies in a distributed manner with no inter-node negotiation or cooperation. Authors state that less inter-node interactions are more attractive for WBANs due to their low overhead and superior scalability. However, such assumption barely adapts to our network model, due to the ever changing topology, the highly dynamic outdoor environment, and the continuously joining and leaving WBANs typical of a BBN scenario.

In [19], we addressed the interference mitigation problem for BBNs considering a centralized approach and we formulated it as an optimization problem. To solve efficiently the problem even for large-scale network scenarios, two heuristic solutions were developed, namely, a customized randomized rounding approach and a tabu search scheme. Our work differs from the work in [19] in two main aspects: (1) we formulate the problem of mutual and cross-technology interference mitigation, considering the Signal-to-Interference-Ratio (SIR) and a noise component related to physical conditions and human body effects, and we therefore allocate WiFi/ZigBee wireless channels to communication links optimizing the SIR ratio, while in [19] the interference was only quantified by the binary decision variables; (2) we address the interference mitigation problem using a distributed approach, with concepts and mathematical tools from Game Theory, while this problem was tackled in [19] in a completely centralized way.

Yet, to the best of our knowledge, this paper is the first to propose a game theoretical approach for an interference-aware channel allocation in BBNs. In our model, multiple WBANs could interact among each other within a BBN, as well as with other coexisting networks/BBNs, involving different access technologies (WiFi, ZigBee, Bluetooth..); this can lead to unavoidable heavy interference environment.

3. System models

In this section, we present the system models, including the network model and the interference model, arising in Body-to-Body Networks.

3.1. Network model

We consider a BBN scenario composed of a set \mathcal{N} of WBANs, which are located in the same geographical area (i.e., a medical center, a rest home or a care home), and share the same unlicensed 2.4 GHz ISM band. Let \mathcal{C}^w and \mathcal{C}^z denote, respectively, the set of WiFi and ZigBee channels in this band.

Each WBAN is equipped with a wearable Mobile Terminal (MT),¹ that uses both the 802.15.4 protocol (i.e., ZigBee) to communicate with the sensor nodes within its WBAN, and the IEEE 802.11 wireless standard (i.e., WiFi) to create a backhaul infrastructure for inter-WBANs' communications.

Since we are assuming that WBANs can move and interact with their surrounding environment, we find ourselves in a quite dynamic BBN scenario, and therefore, we decide to divide the operating time of the whole system into a set T of consecutive epochs, and during each epoch $t \in T$ we suppose that the network topology and environment conditions do not change.

The set $\mathcal{L}^w(t)$ represents all WiFi unidirectional links established by mobile terminals during the epoch $t \in T$;

¹ The WBAN and his corresponding *Mobile Terminal* will be used as synonyms throughout the paper.

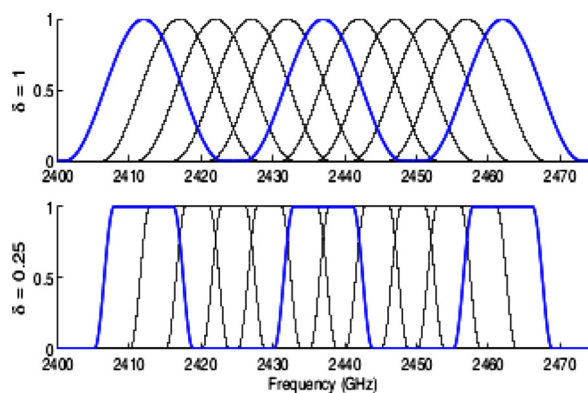


Fig. 3. The 802.11b frequency responses with the raised cosine filter [21].

344 $\mathcal{L}^w(t)$ may vary between two consecutive epochs due to
 345 WBANs' mobility. On the contrary, the set \mathcal{L}^z , which repre-
 346 sents the ZigBee unidirectional links used for intra-WBAN
 347 communication among the sensors, does not change with
 348 time, and for this reason, we omitted the parameter t from
 349 this set.

350 Recent works dealing with interference mitigation have
 351 considered the binary model to represent overlapping be-
 352 tween channels [13–16,19]; i.e. a node is either interfered
 353 or not, however our idea in this work is to quantify the inter-
 354 ference between partially overlapped channels. In [20],
 355 the authors model the overlapping among different WiFi
 356 channels defining a symmetric channel overlapping ma-
 357 trix W , whose element w_{mn} quantifies the degree of in-
 358 terference between channels m and n , and is given as
 359 follows:

$$w_{mn} = \frac{\int_{-\infty}^{+\infty} F_m(w)F_n(w)dw}{\int_{-\infty}^{+\infty} F_m^2(w)dw}, \quad (1)$$

360 where $F_m(w)$ and $F_n(w)$ denote the power spectral den-
 361 sity (PSD) functions of the band-pass filters for channels m
 362 and n , respectively, which can be obtained from the chan-
 363 nels' frequency responses. Yet, we need to know which
 364 channel filter is being used. As in [20], we assume the use
 365 of raised cosine filters, whose principle is explained in [21].
 366 Fig. 3 shows how the PSD function of the IEEE 802.11b de-
 367 pends on the roll-off factor δ , which is a key parameter of
 368 the raised cosine filter, when it is equal to 1 and 0.25, re-
 369 spectively. Hence, [21] gives a simplified expression of the
 370 W matrix:

$$w_{mn} = w_{nm} = \frac{A_o}{A_o + A_{no}} \quad (2)$$

371 where A_o and A_{no} are the overlapping and non-overlapping
 372 areas between the power spectral density (PSD) of chan-
 373 nels m and n , respectively. With expression (2), the W
 374 matrix can be computed off-line and used as a constant ma-
 375 trix in the BR-SIM and SORT-SIM algorithms.

376 Since different wireless technologies use different signal
 377 modulations and access mechanisms, authors in [22]
 378 performed an extensive set of experiments to measure the
 379 partial overlap of the IEEE 802.11b standard, using differ-
 380 ent physical layer modulation methods. First, they consid-
 381 ered 1 and 2 Mbps data-rates for the physical layer, us-

ing the Binary Phase Shift Keying (BPSK) modulation, to
 measure the channel overlap. Then, they reported results
 using the Complementary Code Keying (CCK) modulation,
 with a data-rate of 11 Mbps. It was concluded that partially
 overlapped channels can provide much greater spatial re-
 use if used carefully, depending on the physical separation
 and/or the channel separation between neighboring links,
 whatever the modulation scheme in use.

In this work, we model the channel overlapping prob-
 lem analytically by studying its impact on the signal-
 to-interference ratio. Yet, although the channel overlap-
 ping matrix W has been defined to model the partially-
 overlapped channels for the 802.11b protocol, it does not
 depend, actually, on the technology in use, since the ex-
 pression could involve the PSD functions of any frequency
 responses, provided that the frequency band presents over-
 lapping behaviors, which is actually not the case for Zig-
 Bee and BLE, since both frequency bands present orthogo-
 nal wireless channels.

To summarize, our network model will focus on the fol-
 lowing relevant elements:

- Every single WBAN's MT, equipped with one WiFi antenna and one ZigBee antenna, should dispose of nonoverlapping WiFi and ZigBee channels.
- No interference is present within a WBAN; we assume a TDMA-based medium access control implemented in each WBAN to deal with collisions. Note in addition that there is no interference between adjacent ZigBee channels since there is no overlapping.
- The interference between overlapping WiFi and ZigBee channels is represented by the matrix A , of size $|\mathcal{C}^w| \times |\mathcal{C}^z|$, whose element $a_{c_1c_2}$ is a binary value: $a_{c_1c_2} = 1$ if WiFi channel c_1 overlaps with ZigBee channel c_2 (0 otherwise).
- As in [20], the degree of interference between overlapping WiFi channels is represented by the matrix W , of size $|\mathcal{C}^w| \times |\mathcal{C}^w|$, whose element $w_{c_1c_2} \in [0, 1]$ is a fractional value, defined by the expression in Eq. (1).
- To preserve the network connectivity within the BBN, we assume that all WBANs WiFi interfaces are tuned on the same channel. Therefore, we use the $|\mathcal{L}^w| \times |\mathcal{L}^w|$ matrix $B(t)$, whose element b_{ij} is a binary value: $b_{ij} = 1$ if WiFi links i and j belong to the same BBN at time epoch $t \in \mathcal{T}$ (0 otherwise).
- Finally, WBANs use a higher transmission power on the inter-WBAN channel than on the channel used for intra-WBAN communications (i.e. $p^w \gg p^z$). In particular, data transmissions within ZigBee networks can completely starve due to WiFi communications, which use 10 to 100 times higher transmission power [19].

In order to minimize the total interference within BBNs involving several wireless technologies, it is advantageous to observe every interference component separately, thus we can specify two-kind interference scenarios:

- The Mutual interference:
 - WiFi-WiFi interference at the MT receiver, that occurs while receiving collected data from a nearby WBAN of the same BBN and interfering with adjacent BBNs' WiFi links. Such component includes as well the co-channel interference.

- 442 • ZigBee-ZigBee interference at the MT receiver, that
 443 happens when a ZigBee link of a WBAN interferes
 444 with a ZigBee link of another WBAN belonging to
 445 the same or to a different BBN, when they are allo-
 446 cated the same channel.
- 447 • The Cross-technology interference: WiFi-ZigBee, among
 448 adjacent WBANs, where each WBAN (MT) is communi-
 449 cating with other WBANs over a WiFi link and is sus-
 450 ceptible to interference from nearby ZigBee links, and
 451 vice versa.

452 The *Interference issue* and the *SIR metric* are tightly re-
 453 lated. Thus, in this paper, we would focus on the interfe-
 454 rence metric (SIR) expressed in decibel format by:

$$SIR_i(t)(\text{dB}) = 10\log\left(\frac{g_{ii}(t)p^i}{\sum_{j \neq i} g_{ij}(t)p_j}\right), \quad (3)$$

455 where p^i is the transmission power of transmitter i , $g_{ij}(t)$
 456 is the link gain from transmitter j to receiver i at time
 457 epoch t . Since WBANs can move in their surrounding envi-
 458 ronment, the links' gains $g_{ij}(t)$ vary over time, and the SIR
 459 in turn has been further expressed as a function of time t .

460 The gain parameters are calculated taking into account
 461 the average channel gain evaluated at the reference dis-
 462 tance $d_0 = 1\text{m}$ and with a path loss exponent $n(\alpha)$, ac-
 463 cording to the following formula [23]:

$$g_{ij}(t)|_{\text{dB}} = G(d_0, \alpha)|_{\text{dB}} - 10 \times n(\alpha) \times \log_{10}(d/d_0), \quad \forall i, j \in \mathcal{L}^w(t) \cup \mathcal{L}^z \quad (4)$$

464 Specifically, the average channel gain $G(d_0, \alpha)$, between
 465 WBANs' MTs (Tx Right Hip, Rx Right Hip), significantly
 466 decreases from -37.88 dB to -66.33 dB when switch-
 467 ing from LOS to NLOS conditions, which ensures that our
 468 BBN scenarios are consistent with a realistic human body
 469 environment.

470 3.2. Interference model

471 The interference model defines the set of links that can
 472 interfere with any given link in the network [24]. There
 473 have been various interference models proposed in the lit-
 474 erature; the common concept is that two communication
 475 links $i = (T_i, R_i)$ and $j = (T_j, R_j)$ are interfering if and only
 476 if either T_i or R_i lies within the *interference range* of T_j or
 477 R_j , where T_i, T_j and R_i, R_j designate the transmitter and re-
 478 ceiver interfaces of links i and j , respectively.

479 If modeling the interference characteristics in sensor
 480 networks is challenging, it is more so for BBNs, because
 481 RF characteristics of nodes and environments are neither
 482 known a priori nor computable due to their stochastic,
 483 rapidly changing characteristics [25]. Any routing protocol
 484 working in high interference environment is incapable of
 485 dealing with radio channels suffering from high interfe-
 486 rence ratios. Thus, sharing channels appropriately according
 487 to the interference profiles is mandatory and prior for BBN
 488 networks design.

489 *Interference range* is the range within which nodes in
 490 receive mode will be interfered with an unrelated trans-
 491 mitter and thus suffer from packet loss [26]. For simplic-
 492 ity, ranges are generally assumed concentric which is not

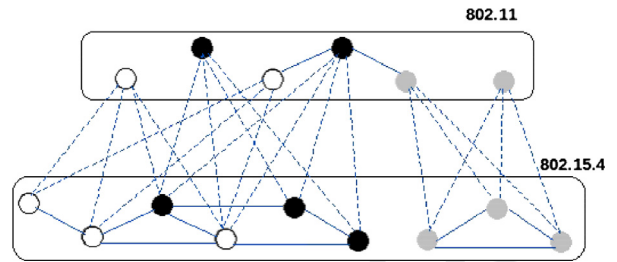


Fig. 4. Cross-technology Conflict graph of the scenario illustrated in Fig. 1.

necessarily given in physical networks. In [26], the inter-
 493 ference range was defined based on SIR, where authors
 494 assume a transmission scenario with transmitter-receiver
 495 distance as d meters and at the same time, an interfering
 496 node r meters away from the receiver, starts another trans-
 497 mission. The received signal is assumed to be successful if
 498 it is above a SIR threshold (SIR_{th}).
 499

Conflict graph: given an interference model, the set
 500 of pairs of communication links that interfere with each
 501 other, assuming mutual and cross-interference in our
 502 model, can be represented using a conflict graph. As done
 503 in [19,27], we depict a conflict graph to model the mutual
 504 and cross-technology interfering wireless links. We adopt
 505 this representation because it will help us in defining the
 506 set of neighbors in next sections for our Socially-aware In-
 507 terference Mitigation game. Therefore, the *cross-technology*
 508 conflict graph $G_c(V_c(t), E_c(t))$ is defined as follows:
 509

- $V_c(t)$: set of vertices corresponding to WiFi and ZigBee
 510 communication links in the network, $V_c(t) = \mathcal{L}^w(t) \cup$
 511 \mathcal{L}^z .
 512
- $E_c(t)$: set of edges corresponding to the interference re-
 513 lationship among pairs of links. Fig. 4 depicts the cross-
 514 technology conflict graph of the three BBN-scenario il-
 515 lustrated in Fig. 1. Solid lines represent conflict edges
 516 between two vertices using the same radio technology,
 517 i.e. $(e_1, e_2) \in E_c(t)$ is a conflict edge if and only if
 518 $e_1, e_2 \in \mathcal{L}^w(t)$ or $e_1, e_2 \in \mathcal{L}^z$, and they are interfering with
 519 each other. Whereas dashed lines correspond to cross-
 520 conflict edges between two vertices using different radio
 521 technologies.
 522

Our goal is to minimize the overall network interfe-
 523 rence. To give an example, let us consider the scenario of
 524 Fig. 1. Each BBN has different interference ranges with its
 525 neighboring BBNs. Assuming that only three WiFi ortho-
 526 gonal channels from the 2.4 GHz band are available (1, 6,
 527 and 11), one trivial solution would be to assign channels 1,
 528 6 and 11 to BBN1, BBN2 and BBN3, respectively. In this
 529 case there would be no interference. Let us assume now
 530 that only two WiFi orthogonal channels 1 and 6 are avail-
 531 able, in addition to channel 2 overlapping with channel 1.
 532 Thus, channels 1, 6 and 2 would be assigned to BBN1, BBN2
 533 and BBN3, respectively. Since BBN1 and BBN3 have dis-
 534 joint interference ranges, they can use overlapping chan-
 535 nels with minimal risk of interference. In practice, the sys-
 536 tem is more complex, with many more BBNs, and/or more
 537 overlapping interference ranges, involving several wire-
 538 less technologies. Therefore a general approach should be
 539

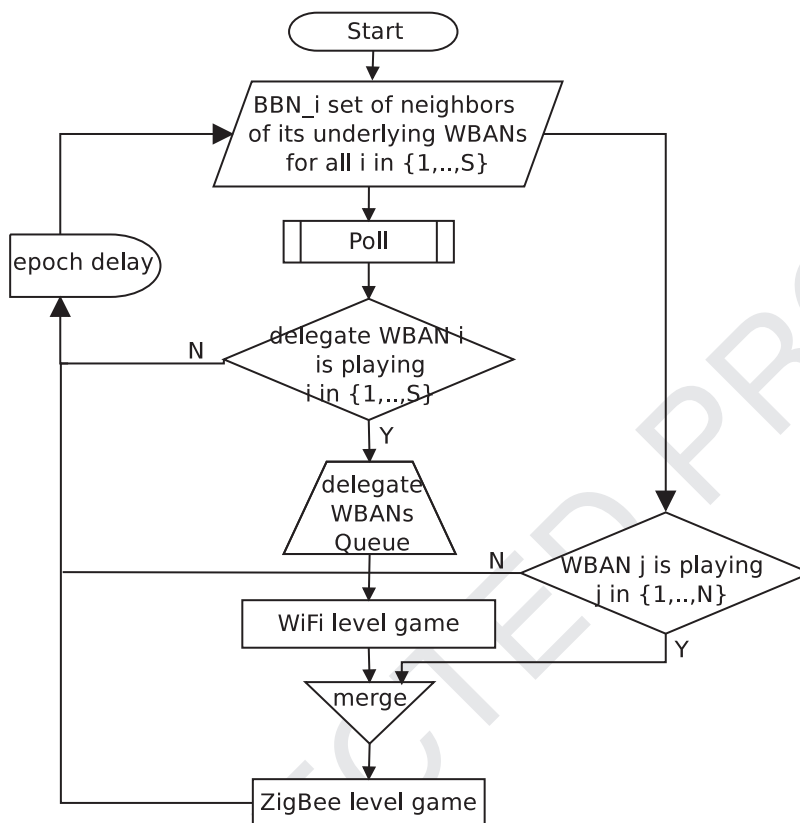


Fig. 5. Flowchart of the two-stage SIM Game: (1) creation of Sub-BBNs and election of the delegates, (2) WiFi level game: allocation of WiFi channels to the set of WiFi links (represented by their delegates), and (3) ZigBee level game: Allocation of ZigBee channels to ZigBee links of WBANs.

540 investigated for an appropriate wireless resource sharing
 541 according to the interference profiles. Likewise, in such
 542 heterogeneous wireless systems, a two-stage channel
 543 allocation scheme is needed; a BBN level game for WiFi
 544 channel allocation step, then a WBAN level game for
 545 ZigBee channel allocation, taking into account the cross-
 546 technology features at each stage.

547 4. Two-stage Socially-aware Interference Mitigation 548 Game (SIM)

549 In this section, we first define the basic notation and
 550 parameters used hereafter, and then we describe in detail
 551 the proposed Socially-aware Interference Mitigation Game
 552 theoretical approach.

553 The lack of a centralized control and prioritization of
 554 access to the radio spectrum, in addition to the restricted
 555 knowledge of network information, motivate us to employ
 556 local interactions for the WiFi and ZigBee level games, in
 557 which players consider their own payoffs as well as those
 558 of their neighbors, so as to optimize their strategies while
 559 relying on their surrounding network information. Besides,
 560 at the BBN-stage game, each group of interacting WBANs
 561 (i.e., each sub-BBN²) is represented by a special player (a

562 delegate or a leader of the group) who decides which WiFi
 563 channel to choose. Indeed, to ensure network connectivity
 564 all WBANs within the same sub-BBN should be tuned to
 565 the same WiFi channel, and we consider this special player
 566 that acts on behalf of the entire sub-BBN. To this end, we
 567 consider in this work a two-stage socially-aware inter-
 568 ference mitigation scheme:

569 (i) At a first stage, each BBN takes a decision on the
 570 WiFi channel that should be assigned to his WiFi trans-
 571 mission links, ensuring minimal interference with his sur-
 572 rounding environment, through a local interaction game
 573 with his neighboring BBNs.

574 (ii) Then, at the second stage, given the WiFi channel
 575 assignment for each BBN, a local interaction game takes
 576 place among the WBANs belonging to the same BBN. After
 577 playing this game, each WBAN (more precisely, each
 578 MT) will be assigned a ZigBee channel to his ZigBee ra-
 579 dio interface, and such assignment guarantees the minimal
 580 interference of the WBAN with his neighboring WBANs.

581 The overall operations for the time epoch $t \in T$ are
 582 represented by the SIM flow chart given in Fig. 5. In this
 583 channel assignment game, the players are the set of links
 584 $\mathcal{L}(t) = \mathcal{L}^W(t) \cup \mathcal{L}^Z$ associated with the set $\mathcal{N} = \{1, \dots, n\}$

² The sub-BBN notation is introduced in order to allow different groups of WBANs, belonging to the same BBN, to communicate on different non-

overlapping WiFi channels. However, when all WBANs (of the same BBN) want to communicate with each other, then the sub-BBNs coincide with their corresponding BBN.

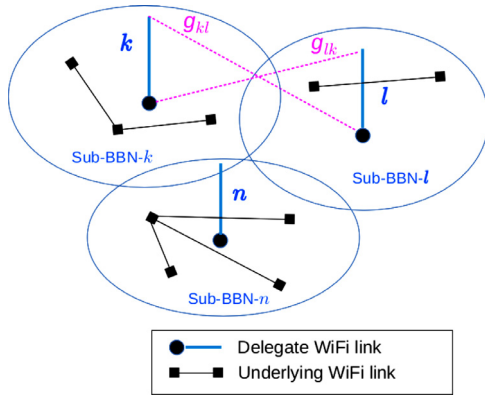


Fig. 6. Delegate and underlying WBANs' WiFi links.

WBANs occupying either the hospital or a care home for old people, and distributed over a set of coexisting BBNs. Each player is represented by a couple of links (l, h) , such that $l \in \mathcal{L}^w(t)$ and $h \in \mathcal{L}^z$ are a WiFi and a ZigBee link corresponding to a given WBAN $i \in \mathcal{N}$ assimilated to its MT. At time epoch $t \in T$, each player chooses a couple of strategies $(s^l(t), s^h(t)) \subset S(t)$, such as $s^l(t)$ is the strategy to allocate a WiFi channel $c_1 \in \mathcal{C}^w$ to the WiFi link $l \in \mathcal{L}^w(t)$ at time epoch $t \in T$, denoted by $x_{c_1}^l$, and $s^h(t)$ is the strategy to allocate a ZigBee channel $c_2 \in \mathcal{C}^z$ to the ZigBee link $h \in \mathcal{L}^z$, denoted by $y_{c_2}^h$. $S(t)$ is obviously the set of the total channel allocation strategies of all players of the BBN scenario. To summarize, the WiFi and ZigBee channel assignment variables are:

$$x_{c_1}^l = \begin{cases} 1, & \text{if WiFi channel } c_1 \text{ is assigned to} \\ & \text{the communication link } l \\ 0, & \text{otherwise} \end{cases}$$

$$y_{c_2}^h = \begin{cases} 1, & \text{if ZigBee channel } c_2 \text{ is assigned to} \\ & \text{the communication link } h \\ 0, & \text{otherwise} \end{cases}$$

Hence, hereafter, we first begin with presenting the first-stage game, to choose a WiFi channel assignment for each MT, and then we describe in detail the second-stage game, where each MT is further assigned a ZigBee channel.

4.1. BBN-stage SIM game

In order to assign a single WiFi channel to each sub-BBN, we opt for a BBN-stage SIM game so that each set of communicating WBANs, forming a sub-BBN, are represented by a specific WiFi link. The representative WiFi link is situated in the center of the sub-BBN and plays the role of the *delegate*, and the other WBANs belonging to the same sub-BBN will be allocated the same WiFi channel (Fig. 6). Our choice of the representative WiFi link is similar to the one made by Govindasamy et al. in [28]. In fact, the work in [28] presents a technique to find the spectral efficiency of an interference-limited representative link with an arbitrary distribution of interference powers, within an ad hoc network with randomly distributed multi-antenna links. This model considers a circular network where the representative receiver is assumed to be at the origin of the circle, and the interferers are links with

other receivers whose locations do not impact the representative link. Of course, there exist a variety of different mechanisms/solutions to select the more appropriate delegate/representative link in the network. However, this issue is not the main concern of this paper and deserves a deep study.

We build the cross-technology conflict graph and we assume that each WBAN has information only about his sub-BBN underlying WBANs, through the exchange of polling messages. Thus, we can identify for each WBAN, the set of interfering neighbors at time epoch $t \in T$ (i.e., the set of edges between a link of such WBAN and transmission links of the others). Let W_l denote the set of links interfering with WiFi link l :

$$W_l(t) = \{k \in \mathcal{L}^w(t) : (l, k) \in E_c(t)\} \cup \{j \in \mathcal{L}^z : (l, j) \in E_c(t)\}$$

Thereby, we can define the BBN-stage game (\mathcal{G}_1) as follows:

- *Players*: the set of BBNs represented by their delegates, such as a delegate player per sub-BBN. For the BBN-stage, the player is assimilated to its WiFi link l .
- *Strategies/actions*: $s^l(t) = x_{c_1}^l(t)$, strategy to choose a WiFi channel c_1 for WiFi link l from the set of available channels in \mathcal{C}^w .
- *Utility function*: to ensure a realistic representation of the game, we use the worst SIR values perceived by the two radio interfaces, WiFi and ZigBee, as utility function.

Hereafter, we describe the SIR given in Eq. (3) that we extend to consider interfering transmitters using different technologies. It is worth noting that Eq. (5) can be easily extended to more than two radio technologies, considering further for example Bluetooth. However, to simplify the analysis we conduct the study with only two components, corresponding to WiFi and ZigBee, respectively. Whence, the SIR of the player $l \in \mathcal{L}^w$, considering the WiFi interface, is given by:

$$SIR^w(x_{c_1}^l)(t) = 10 \log \left(\frac{g_{ll} p_w^l}{I_{c_1}^w(x_{c_1}^l) + I^w(x_{c_1}^l) + I^{wz}(x_{c_1}^l)} \right), \quad (5)$$

where

$I_{c_1}^w(x_{c_1}^l)$: co-channel interference from WiFi links of other sub-BBNs ($b_{kl} = 0$) sharing WiFi channel c_1 with WiFi link l .

$$I_{c_1}^w(x_{c_1}^l) = \sum_{\substack{k \in \mathcal{L}^w \\ b_{kl} = 0}} x_{c_1}^k x_{c_1}^l g_{lk} p_w^k \quad (6)$$

$I^w(x_{c_1}^l)$: mutual interference from WiFi links of other sub-BBNs ($b_{kl} = 0$) using WiFi channels that overlap with c_1 .

$$I^w(x_{c_1}^l) = \sum_{\substack{k \in \mathcal{L}^w \\ b_{kl} = 0}} \left(\sum_{\substack{c \in \mathcal{C}^w \\ c \neq c_1}} w_{c_1 c} x_{c_1}^l x_c^k \right) g_{lk} p_w^k, \quad (7)$$

$I^{wz}(x_{c_1}^l)$: cross-interference from ZigBee links, using ZigBee channels other than c_2 , overlapping with c_1

$$I^{wz}(x_{c_1}^l) = \sum_{\substack{k \in \mathcal{L}^z \\ k \neq h}} \left(\sum_{c \in \mathcal{C}^z} a_{c_1 c} x_{c_1}^l y_c^k \right) g_{lk} p_z^k, \quad (8)$$

665 g_{ll} is the channel gain of link l , g_{lk} the link gain from the
666 transmitter k to the receiver l , p_w^k and p_z^k are the WiFi and
667 ZigBee transmit power, respectively.

668 Note that in expression (8) we use the binary parameter
669 $a_{c_1 c_2}$ to model the cross-technology interference instead
670 of the fractional $w_{c_1 c_2}$ used in Eq. (7) for mutual WiFi interference.
671 In fact, although in the literature the interference of the IEEE 802.11b
672 has been modeled as an additive white Gaussian noise (AWGN) to the
673 ZigBee signal, the experimental results performed in [29] show a significant
674 performance degradation for ZigBee links in the presence of WiFi
675 transmissions. Specifically, the authors measured a packet loss of 99,
676 75% up to 100% in WBANs used for blood analysis and ECG sensing
677 when a video streaming is executed over an interfering WiFi channel.
678 Therefore, due to the tight constraints on WBANs' transmissions reliability,
679 we consider the worst effect caused by WiFi interference on ZigBee
680 communications, using the binary parameter $a_{c_1 c_2} \in \{0, 1\}$.

684 (1) Convergence of BBN-stage game: Nash Equilibrium

685 Having defined the BBN stage of the SIM game, we then
686 demonstrate that such game indeed admits at least one pure-strategy
687 Nash Equilibrium. Thus, we first define the utility function of player
688 l as follows:

$$U_w(x_{c_1}^l) = 10 \log(g_{ll} p_w^l) - 10 \log(IF_l^w(x_{c_1}^l)) \quad (9)$$

689 where $IF_l^w(x_{c_1}^l)$, denoted as the WiFi Interference Function of player
690 l , is the total interference suffered by link l when playing strategy
691 $x_{c_1}^l$, and is expressed as follows:

$$\begin{aligned} IF_l^w(x_{c_1}^l) &= I_c^w(x_{c_1}^l) + I^w(x_{c_1}^l) + I^{wz}(x_{c_1}^l) \\ &= \sum_{k \in W_l \cap L^w} \sum_{c \in C^w} f(x_{c_1}^l, x_c^k) + \sum_{\substack{j \in W_l \cap L^z \\ c \in C^z \\ c \neq c_2}} \sum g(x_{c_1}^l, y_c^j) \end{aligned}$$

692 or function of the strategies:

$$IF_l^w(s^l) = \sum_{k \in W_l \cap L^w} f(s^l, s^k) + \sum_{j \in W_l \cap L^z} g(s^l, s^j) \quad (10)$$

693 where:

$$f(s^l, s^k) = \begin{cases} 0, & s^l \neq s^k \text{ and WiFi channel } c_1 \text{ of link } l \\ & \text{does not overlap with WiFi} \\ & \text{channel of link } k. \\ g_{lk} p_w^k, & s^l = s^k \\ w_{c_1 c} g_{lk} p_w^k, & s^l \neq s^k \text{ and WiFi channel } c_1 \text{ of link } l \\ & \text{overlaps with WiFi channel of link } k. \end{cases}$$

694 and:

$$g(s^l, s^j) = \begin{cases} 0, & \text{WiFi channel } c_1 \text{ of link } l \text{ does not overlap} \\ & \text{with ZigBee channel of link } j. \\ g_{lj} p_z^j, & \text{WiFi channel } c_1 \text{ of link } l \text{ overlaps with} \\ & \text{ZigBee channel of link } j. \end{cases}$$

695 We observe that the maximization of utility function
696 U_w corresponds to the minimization of the Interference
697 Function IF^w . Due to the property of *monotone transformation*,
698 if the modified game with utility IF^w is a potential
699 game, then the original BBN-stage SIM game with utility

700 U_w is also a potential game with the same potential
701 function. Then, the BBN-stage SIM game (\mathcal{G}_1) is expressed
702 as follows:

$$(\mathcal{G}_1) : \min_{x_{c_1}^l \in S^l(t)} IF_l^w(x_{c_1}^l, x_{c_1}^{-l}) \forall l \in L^w$$

$$s.t. \sum_{c \in C^w} x_c^l = 1 \quad \forall l \in L^w(t) \quad (11)$$

$$x_{c_1}^l \in \{0, 1\} \quad \forall l \in L^w(t), c_1 \in C^w, \quad (12)$$

703 For convenience, we designate by $-l$ all the players be-
704 longing to W_l . Constraint (11) forces the assignment of
705 a single WiFi channel for a single WiFi link for each
706 player, the connectivity within the sub-BBNs is ensured
707 so that a unique WiFi channel is allocated to every pair
708 of links belonging to the same sub-BBN through the ex-
709 change of polling messages between the delegate player
710 and the other players of each sub-BBN. The convergence
711 of the BBN-stage SIM game to a Nash Equilibrium is given
712 by the following theorem:

714 **Theorem 1.** The BBN-stage SIM game \mathcal{G}_1 is an exact poten-
715 tial game.

716 **Proof.** we construct the potential function as follows:

$$\Phi^w(s^l, s^{-l}) = \frac{1}{2} \sum_{i \in L^w} \sum_{k \in W_i \cap L^w} f(s^i, s^k) + \sum_{i \in L^w} \sum_{j \in W_i \cap L^z} g(s^i, s^j) \quad (13)$$

717 Therefore, when player $l \in L^w$ changes its action at time
718 epoch $t \in \mathcal{T}$, from s^l to s^l , the variation of the poten-
719 tial function subsequent to this player's strategy change is
720 given by:

$$\begin{aligned} \Phi^w(s^l, s^{-l}) - \Phi^w(s^l, s^{-l}) &= \frac{1}{2} \sum_{i \in L^w} \sum_{\substack{k \in W_i \cap L^w \\ i \neq l}} f(s^i, s^k) + \sum_{i \in L^w} \sum_{\substack{j \in W_i \cap L^z \\ i \neq l}} g(s^i, s^j) \end{aligned} \quad (14)$$

$$- \frac{1}{2} \sum_{i \in L^w} \sum_{\substack{k \in W_i \cap L^w \\ i \neq l}} f(s^i, s^k) - \sum_{i \in L^w} \sum_{\substack{j \in W_i \cap L^z \\ i \neq l}} g(s^i, s^j) \quad (15)$$

$$+ \frac{1}{2} \sum_{i \in L^w} f(s^i, s^l) - \frac{1}{2} \sum_{i \in L^w} f(s^i, s^l) \quad (k=l) \quad (16)$$

$$+ \frac{1}{2} \sum_{k \in W_l \cap L^w} f(s^l, s^k) + \sum_{j \in W_l \cap L^z} g(s^l, s^j) \quad (i=l) \quad (17)$$

$$- \frac{1}{2} \sum_{k \in W_l \cap L^w} f(s^l, s^k) - \sum_{j \in W_l \cap L^z} g(s^l, s^j) \quad (i=l) \quad (18)$$

725 We can easily see that (14) + (15) = 0. On the other
726 hand, since each player has only interference with his
727 neighboring set, then $\{i \in L^w : i \neq l\} = \{k \in W_l \cap L^w\}$, and
728 we assume that function f is symmetric so as we con-
729 sider symmetric channel gains ($g_{lk} = g_{kl}$ if $b_{kl} = 0$, Fig. 6),
730 therefore:

$$\begin{aligned} \Phi^w(s^l, s^{-l}) - \Phi^w(s^l, s^{-l}) &= \sum_{k \in W_l \cap L^w} f(s^l, s^k) + \sum_{j \in W_l \cap L^z} g(s^l, s^j) \end{aligned} \quad (19)$$

$$- \sum_{k \in W_l \cap L^w} f(s^l, s^k) - \sum_{j \in W_l \cap L^z} g(s^l, s^j) \quad (20)$$

$$= IF_1^w(s^l, s^{-l}) - IF_1^w(\hat{s}^l, s^{-l}) \quad (21)$$

733 Accordingly we prove that, when a delegate $l \in \mathcal{L}^w$ de-
734 viates from a strategy s^l to an alternate strategy \hat{s}^l , the
735 change in the exact potential function Φ^w exactly mirrors
736 the change in l 's utility. Therefore the BBN-stage SIM game
737 is an exact potential game. \square

738 Thereby, we can rely on the following theorem [30] to
739 confirm the existence of a Nash Equilibrium to our game.

740 **Theorem 2.** Every potential game has at least one pure Nash
741 Equilibrium, namely the strategy s^l that minimizes $\Phi^w(s^l)$.

742 The result of Theorem 2 motivates us to design the Best
743 Response SIM algorithm in Section 5 to resolve the BBN-
744 stage SIM game.

745 4.2. WBAN-stage SIM game

746 We now consider the WBAN-stage game, where each
747 WBAN will be assigned a ZigBee channel to his ZigBee
748 radio interface, that guarantees the minimal interference
749 with his neighbors.

750 (1) ZigBee local interaction game

751 Similarly to the BBN stage, denote Z_h as the set of
752 neighbors of ZigBee link h , including the set of edges be-
753 tween ZigBee link h and interfering WiFi and ZigBee links,
754 using the conflict graph:

$$Z_h(t) = \{j \in \mathcal{L}^z : (h, j) \in E_c(t)\} \cup \{k \in \mathcal{L}^w(t) : (h, k) \in E_c(t)\}$$

755 Hence, we can define the local interaction game of the
756 WBAN stage (\mathcal{G}_2) as follows:

- 757 • *Players:* set \mathcal{N} of WBANs. For the WBAN-stage, the
758 player is assimilated to his ZigBee link h .
- 759 • *Strategies/actions:* $s^h(t) = y_{c_2}^h(t)$, strategy to choose a
760 ZigBee channel c_2 for ZigBee link h from the set of
761 available channels in \mathcal{C}^z .
- 762 • *Utility function:* is, similarly to BBN stage, function of
763 the SIR considering the ZigBee interface which is used
764 for intra-WBAN communications, given by:

$$SIR^z(y_{c_2}^h)(t) = 10\log\left(\frac{g_{hh}p_z^h}{I^{wz}(y_{c_2}^h) + I^z(y_{c_2}^h)}\right), \quad (22)$$

765 $I^{wz}(y_{c_2}^h)$ represents the cross-technology interference caused
766 by mobile terminals using WiFi channels that interfere
767 with the ZigBee channel c_2 on which WBAN link h is
768 tuned.

$$I^{wz}(y_{c_2}^h) = \sum_{\substack{k \in \mathcal{L}^w \\ b_{kl}=0}} \sum_{c \in \mathcal{C}^w} a_{cc_2} x_c^k y_{c_2}^h g_{hk} p_w^k(t). \quad (23)$$

769 $I^z(y_{c_2}^h)$ accounts for the co-channel interference of nearby
770 WBANs sharing the same ZigBee channel c_2 of player h .

$$I^z(y_{c_2}^h) = \sum_{k \in \mathcal{L}^z} y_{c_2}^k y_{c_2}^h g_{hk} p_z^k(t). \quad (24)$$

771 Conversely to the BBN stage (Eq. (5)), in Eq. (22) only
772 cross and co-channel interference components are consid-
773 ered at the denominator, since all ZigBee channels are

completely orthogonal among each other, i.e. no mutual
interference is there. In case of sharing the same Zig-
Bee channel, i.e., expression (24), the corresponding ex-
perimental scenario in [29] measures 18% of packet losses,
which led to the conclusion that the impact of ZigBee
co-channel interference may be significant. Therefore, we
model our game so that selecting different and non-
overlapping ZigBee channels for intra-WBAN communica-
tions emerges as the best strategy for all players. Unlike
BBN-stage game where a unique WiFi channel is required
by a sub-BBN, in WBAN stage, WBANs of the same sub-
BBN use different ZigBee channels for intra-WBAN commu-
nications. Yet, to ensure a fair sharing of available ZigBee
resources within BBNs, we consider local interaction be-
haviors among players interacting within the same neigh-
boring set, which is translated in the utility function by a
local cooperation quantity as a tradeoff to the player self-
ish attitude. Thus, we define the utility function of player h
for the WBAN-stage game as follows:

$$U_z(y_{c_2}^h) = SIR^z(y_{c_2}^h) + \sum_{k \in Z_h} SIR^z(y_c^k) \quad (25)$$

$$= 10\log(g_{hh}p_z^h) + \sum_{k \in Z_h} 10\log(g_{kk}p_z^k) - IF_h^z(y_{c_2}^h)$$

where: $IF_h^z(y_{c_2}^h) = I_h(y_{c_2}^h) + \sum_{k \in Z_h} I_k(y_{c_2}^h)$ 793

and: $I_k(y_{c_2}^h) = 10\log(I^{wz}(y_c^k) + I^z(y_c^k)), \forall c \in \mathcal{C}^z : y_c^k = 1$ 794

$I_k(s^h)$, with $s^h = y_{c_2}^h$, is the total interference suffered by
link k of a neighboring WBAN when link h plays strategy
 $y_{c_2}^h$. 795-797

As in [31], using the monotone transformation property,
the WBAN-stage SIM game is expressed as follows: 798-799

$$(\mathcal{G}_2) : \min_{y_{c_2}^h \in S^h(t)} IF_h^z(y_{c_2}^h, y_{c_2}^{-h}) \quad \forall h \in \mathcal{L}^z$$

$$s.t. \sum_{c \in \mathcal{C}^z} y_c^h = 1 \quad \forall h \in \mathcal{L}^z(t) \quad (26)$$

$$y_c^h \in \{0, 1\} \quad \forall h \in \mathcal{L}^z, c \in \mathcal{C}^z \quad (27) \quad 800$$

Constraint (26) forces the assignment of a single ZigBee
channel for a ZigBee link, for each player. 801-802

(2) Convergence of WBAN-stage game: Nash Equilibrium 803

The property of the proposed local interaction game is
characterized by the following theorem: 804-805

Theorem 3. \mathcal{G}_2 is an exact potential game which has at least
one pure strategy NE, and the optimal solution of its potential
function constitutes a pure strategy NE. 806-807-808

Proof. we construct the potential function as follows: 809

$$\Phi^z(s^h, s^{-h}) = \sum_{k \in \mathcal{L}^z} I_k(s^h, s^{-h})$$

if we compute the variation of the utility function when
player $h \in \mathcal{L}^z$ changes its action at time epoch $t \in T$, from
 s^h to \hat{s}^h , we obtain: 810-811-812

$$IF_h^z(s^h, s^{-h}) - IF_h^z(\hat{s}^h, s^{-h}) = I_h(s^h, s^{-h}) - I_h(\hat{s}^h, s^{-h})$$

$$+ \sum_{k \in Z_h} [I_k(s^h, s^{-h}) - I_k(\hat{s}^h, s^{-h})] \quad (28)$$

813 On the other hand, the variation of the potential func-
814 tion subsequent to this player's strategy change is given
815 by:

$$\begin{aligned} \Phi^z(s^h, s^{-h}) - \Phi^z(\hat{s}^h, s^{-h}) &= \sum_{k \in L^z} I_k(s^h, s^{-h}) - \sum_{k \in L^z} I_k(\hat{s}^h, s^{-h}) \\ &= I_h(s^h, s^{-h}) - I_h(\hat{s}^h, s^{-h}) + \sum_{k \in Z_h} [I_k(s^h, s^{-h}) - I_k(\hat{s}^h, s^{-h})] \\ &\quad + \sum_{k \in L^z \setminus Z_h, k \neq h} [I_k(s^h, s^{-h}) - I_k(\hat{s}^h, s^{-h})] \end{aligned} \quad (29)$$

816 Yet, with the local cooperative nature of WBAN-stage
817 game, h player's action only affects players in its interfer-
818 ence range, thus we have:

$$I_k(s^h, s^{-h}) - I_k(\hat{s}^h, s^{-h}) = 0 \quad \forall k \in L^z \setminus Z_h, k \neq h$$

819 This leads to the following equation:

$$IF_h^z(s^h, s^{-h}) - IF_h^z(\hat{s}^h, s^{-h}) = \Phi^z(s^h, s^{-h}) - \Phi^z(\hat{s}^h, s^{-h})$$

820 Accordingly we prove that, when a player $h \in L^z$ de-
821 viates from a strategy s^h to an alternate strategy \hat{s}^h , the
822 change in the exact potential function Φ^z exactly mirrors
823 the change in h 's utility.

824 Therefore the WBAN-stage SIM game is an exact poten-
825 tial game. \square

826 4.3. A discussion on social interactions of WBANs 827 in the SIM games

828 The social information in the BBN and WBAN level
829 games can be collected by using a signaling protocol, like
830 one of those presented in [27,32], to allow mobile termi-
831 nals to exchange control messages (on proximity informa-
832 tion) among each other in order to build (and maintain)
833 the network topology and the conflict graph, and then
834 compute in a completely distributed fashion the chan-
835 nel assignment that minimizes the (mutual and cross-
836 technology) interference (or maximizes the SIR at WiFi and
837 ZigBee radio interfaces), based on local information.

838 More in detail, we recall that our WiFi and ZigBee util-
839 ity functions rely on the neighboring sets of a WBAN MT's
840 WiFi and ZigBee pair of links (l, h) , defined as:

$$\begin{aligned} W_l(t) &= \{k \in L^w(t) : (l, k) \in E_c(t)\} \cup \{j \in L^z : (l, j) \in E_c(t)\} \\ Z_h(t) &= \{j \in L^z : (h, j) \in E_c(t)\} \cup \{k \in L^w(t) : (h, k) \in E_c(t)\} \end{aligned}$$

841 Link-state messages are used to spread topology infor-
842 mation to the entire network. A link-state message con-
843 tains two lists of WiFi and ZigBee neighbors, each iden-
844 tified by its WBAN and BBN identifiers. Such messages are
845 used by the BBN players to build the network topology and
846 the conflict graph. Then, WBANs' MTs send beacon mes-
847 sages to their neighbors, recognized in their neighboring
848 sets $(W_l(t), Z_h(t))$.

849 For example, a WiFi beacon message is only sent to the
850 delegates of neighboring BBNs, since a single WiFi chan-
851 nel should be selected by each BBN. Such message contains
852 the identifier of the WBAN, a list of neighbors (from which
853 control traffic has been recently received), and his local in-
854 formation, needed for the utility functions of his neighbors,
855 i.e., $x_{c_1}^k$ and $y_{c_2}^j$, where c_1 and c_2 are the WiFi and Zig-
856 Bee channels selected by his WiFi and ZigBee links (k, j) . In

contrast, the ZigBee beacon message is sent to his neigh- 857
858 boring WBANs, within the same BBN, evenly, and contains
859 in addition his SIRz value needed by the local interaction
860 game, as explained hereafter.

861 Upon receiving a beacon message, the interference mit-
862 igation algorithm (BR-SIM) extracts the information neces-
863 sary to update the utility function. In particular, for
864 each WBAN receiving a ZigBee beacon message from a
865 neighboring WBAN, BR-SIM extracts the SIRz advertised in
866 the beacon message, and updates his utility function, by
867 adding this SIRz value to the local cooperation quantity,
868 as a tradeoff to the player selfish attitude (Eq. (25)). For a
869 detailed description of the information exchange protocol,
870 please refer to our previous work [27].

871 5. Best-Response algorithm for SIM game (BR-SIM)

872 Potential games have two appealing properties: they
873 admit at least one pure-strategy NE which can be ob-
874 tained through a best-response dynamics carried out by
875 each player, and they have the Finite Improvement Prop-
876 erty (FIP) [33], which ensures the convergence to a NE
877 within a finite number of iterations. In the following, we
878 propose an iterative algorithm (Algorithm 1) that imple-
879 ments a best response dynamics for our proposed game.

880 Algorithm 1 takes as input the current time epoch
881 $t \in \mathcal{T}$, the set \mathcal{N} of WBANs, the conflict graph $G_c(V_c(t),$
882 $E_c(t))$, the available WiFi and ZigBee channels (C^w, C^z) ,
883 the channel gain, the mutual and cross-technology chan-
884 nel overlapping, and the network connectivity matrices
885 $(\mathcal{G}, \mathcal{W}, \mathcal{A}, \mathcal{B}(t))$. It gives as output the channel allocation
886 matrices $X_w(t)$ and $Y_z(t)$, the minima of the WiFi and Zig-
887 Bee Interference Functions obtained at the Nash Equilib-
888 rium, and the number of iterations NE_{iter} needed to con-
889 verge to a NE point.

890 Algorithm 1 starts by forming the coalitions of sub-
891 BBNs whose delegates are representative WiFi links situ-
892 ated in the center with symmetric gains. The delegates
893 and the underlying WBANs are initialized to random WiFi
894 and ZigBee channels with respect to the connectivity crite-
895 rion within BBNs. Then, the algorithm iteratively examines
896 whether there exists any player that is unsatisfied, and in
897 such case a greedy selfish step is taken so that such player
898 l changes his current strategy $s^l(\tau)$, $\tau < t$, to a better strat-
899 egy $s^l(\tau + 1)$ with respect to the current action profile of
900 all other players, as follows:

$$\begin{aligned} s^l(\tau + 1) &= \arg \min_{s^l \in C^w} IF_l^w(s^l, s^{-l}) \quad s.t. \\ s^{-l} &= \{s^1(\tau + 1), s^2(\tau + 1), \dots, s^{l-1}(\tau + 1), \\ &\quad s^{l+1}(\tau), \dots, s^{l^w(t)}(\tau)\} \end{aligned} \quad (30)$$

901 where s^1, s^2, \dots, s^{l-1} have been updated to their best-
902 responses at iteration $\tau + 1$ and do not change from their
903 selected strategies during the current iteration.

904 Alike the WiFi Best-response procedure, players itera-
905 tively update the ZigBee channels that minimize their In-
906 terference Functions, with respect to their WiFi channels
907 selected at the BBN- (or WiFi-) stage step. Thus, for a Zig-
908 Bee player h , the strategy domain of the ZigBee channel
909 selection process is delimited to the set of available Zig-
910 Bee channels $C_h^z(t)$, i.e., not overlapping with his assigned

Algorithm 1: SIM Best Response NE (BR-SIM).

Input : $t \in \mathcal{T}, \mathcal{N}, G_c(V_c(t), E_c(t)), \mathcal{C}^w, \mathcal{C}^z, \mathcal{G}, \mathcal{W}, \mathcal{A}, \mathcal{B}(t)$
Output: $X_w(t), Y_z(t), IF_{min}^w(t), IF_{min}^z(t), NE_{iter}$

- 1 **Initialization**
- 2 Grouping of sub-BBNs and election of the set of delegates: $L_{delegates}^w$;
- 3 Set randomly WiFi and ZigBee action-tuples at $t=0$,
 $S^w(0) = \{s_0^1, s_0^2, \dots, s_0^{|L_{delegates}^w|}\}$ and $S^z(0) = \{s_0^1, s_0^2, \dots, s_0^{|L^z|}\}$;
- 4 **end Initialization**
- 5 **while** $S^w(\tau)$ is not a Nash Equilibrium **do**
- 6 **for** $l \in L_{delegates}^w$
- 7 better response update $s^l(\tau + 1)$: select the WiFi channel that minimizes its Interference Function according to (30);
- 8 **end for**
- 9 Set the delegates action profile to $S^w(\tau + 1) = \{s^1(\tau + 1), s^2(\tau + 1), \dots, s^{|L_{delegates}^w|}(\tau + 1)\}$;
- 10 Calculate $IF^w(\tau + 1) = \{IF_1^w(\tau + 1), \dots, IF_{|L_{delegates}^w|}^w(\tau + 1)\}$;
- 11 $\tau = \tau + 1$;
- 12 $NE_{iter}++$;
- 13 **end while**
- 14 $S^w(t) = S^w(\tau)$ is a Nash Equilibrium, delegates communicate their WiFi channel selections to WBANs;
- 15 Set the BBN-stage action profile
 $S^w(t) = \{s^1(t), s^2(t), \dots, s^{|L^w|}(t)\}$ and $X_w(t)$ matrix;
- 16 **while** $\min IF^z(\tau)$ is not reached **do**
- 17 Repeat steps 6–11 for $h \in L^z$ to select the ZigBee channels that minimize the players Interference Function according to (31);
- 18 $NE_{iter}++$;
- 19 **end while**
- 20 Set the WBAN-stage action profile
 $S^z(t) = \{s^1(t), s^2(t), \dots, s^{|L^z|}(t)\}$ and $Y_z(t)$ matrix.

911 WiFi channel at time epoch t . Therefore, the best-response
 912 strategy of ZigBee player h is expressed by:

$$s^h(\tau + 1) = \arg \min_{s^h \in \mathcal{C}_h^z(t)} IF_h^z(s^h, s^{-h}) \quad s.t. \quad (31)$$

$$s^{-h} = \{s^1(\tau + 1), s^2(\tau + 1), \dots, s^{h-1}(\tau + 1), s^{h+1}(\tau), \dots, s^{|L^z(t)|}(\tau)\} \quad (31)$$

913 Due to the FIP property, such algorithm is guaranteed
 914 to converge in a finite number of iterations to a BBN-stage
 915 NE, and then to a local interaction ZigBee NE where no
 916 player has an incentive to deviate from his best-response
 917 choice.

918 6. Sub-Optimal Randomized Trials for SIM 919 game (SORT-SIM)

920 In large-scale networks with several BBNs, especially in
 921 real-time-constrained applications, the exhaustive search
 922 of NE can be extremely time consuming. Therefore, we
 923 propose, as an alternative solution, the SORT-SIM algo-
 924 rithm to deal with this specific issue. SORT-SIM is based on
 925 the principle of ensuring feasible SIR values for all players

while allowing them to play simultaneously, and reducing
 the probability of channel selection conflicts. 926 927

Algorithm 2 takes the same inputs as **Algorithm 1**, and
 gives the same outputs, i.e., the channel allocation matrices
 $X_w(t)$ and $Y_z(t)$, the minima of the Interference Functions,
 and the number of iterations $SORT_{iter}$ needed to reach the
 sub-optimal solution. 928 929 930 931

Algorithm 2: SIM Sub-Optimal Randomized Trials (SORT-SIM).

Input : $t \in \mathcal{T}, \mathcal{N}, G_c(V_c(t), E_c(t)), \mathcal{C}^w, \mathcal{C}^z, \mathcal{G}, \mathcal{W}, \mathcal{A}, \mathcal{B}(t)$
Output: $X_w(t), Y_z(t), IF^w(t), IF^z(t), SORT_{iter}$

- 1 Grouping of sub-BBNs and election of the set of delegates $L_{deleg}^w(t)$
- 2 **for** delegate WiFi link $l \in L_{deleg}^w(t)$
- 3 Calculate the set of neighbors \mathcal{W}_l ;
- 4 Calculate the set of free WiFi channels $\mathcal{C}_{free}^w(l)$;
- 5 **end for**
- 6 **while** $IF^w(\tau)$ is not a sub-optimal solution **do**
- 7 **for** delegate WiFi link $l \in L_{deleg}^w(t)$
- 8 **if** $\mathcal{C}_{free}^w \neq \emptyset$ **then** Randomly select WiFi channel c_1
 from $\mathcal{C}_{free}^w(l)$;
- 9 **else** Randomly select WiFi channel c_1 such as
 $SIR^w(x_c^l) > SIR_{th}^w$; **end if**
- 10 **end for**
- 11 Delegates communicate their WiFi channels
 selections to the underlying WBANs;
- 12 Set the BBN-stage channel allocation matrix $X_w(t)$;
 Calculate $IF^w(\tau) = \{IF_1^w(\tau), \dots, IF_{L^w}^w(\tau)\}$;
- 13 $\tau = \tau + 1$;
- 14 $SORT_{iter}++$;
- 15 **end while**
- 16 **for** ZigBee links $h \in L^z(t)$
- 17 Calculate the set of available ZigBee channels for
 link h , $\mathcal{C}^z(h)$;
- 18 Calculate the set of neighbors \mathcal{Z}_h ;
- 19 Calculate the set of free ZigBee channels \mathcal{C}_{free}^z from
 $\mathcal{C}^z(h)$;
- 20 **end for**
- 21 **while** $IF^z(\tau)$ is not a sub-optimal solution **do**
- 22 **for** ZigBee links $h \in L^z(t)$
- 23 **if** $\mathcal{C}_{free}^z(h) \neq \emptyset$ **then** Randomly select ZigBee
 channel c_2 from $\mathcal{C}_{free}^z(h)$;
- 24 **else** Randomly select ZigBee channel $c_2 \in \mathcal{C}^z(h)$
 such as $SIR^z(y_c^h) > SIR_{th}^z$; **end if**
- 25 **end for**
- 26 Set the WBAN-stage channel allocation matrix $Y_z(t)$;
 Calculate $IF^z(\tau) = \{IF_1^z(\tau), \dots, IF_{L^z}^z(\tau)\}$;
- 27 $\tau = \tau + 1$;
- 28 $SORT_{iter}++$;
- 29 **end while**

At the beginning, **Algorithm 2** describes the main steps
 relative to the grouping of sub-BBNs, the election of their
 representative links and the calculation of their corre-
 sponding set of neighbors. Then, the WiFi channel alloca-
 tion is performed, for each delegate l , as follows: 932 933 934 935 936 937

- (i) First, select randomly a WiFi channel from the list of free WiFi channels, if available, i.e., not allocated in neighboring set of link l (step 8).

$$C_{free}^w(l) = \{c \in C^w : \forall k \in \mathcal{W}_l(t) \cap \mathcal{L}^w(t), x_c^k = 0\}$$

- (ii) If no free channel is available, calculate at step 9 the utility (SIR^w) for each delegate and select randomly from the list, WiFi channels that provide an SIR^w above the threshold value (SIR_{th}^w).

$$c_1 = \begin{cases} \text{Rand}(C_{free}^w(l)), & \text{if } C_{free}^w(l) \neq \emptyset \\ \text{Rand}\{c \in C^w : SIR^w(x_c^l) > SIR_{th}^w\}, & \text{otherwise.} \end{cases} \quad (32)$$

- (iii) To ensure a fair sharing of resources, a WBAN should release his WiFi channel after at most θ_s . θ is defined as the maximum time of reservation of the wireless channel, and is assumed as a configurable parameter.
- (iv) Finally, the WBANs belonging to the same sub-BBN are tuned on the WiFi channel selected by their leader.

The previous operations are iteratively repeated until reaching a number of trials where no WBAN has an incentive to deviate from his channel choice, presenting, thus, a sub-optimal solution for the SIM problem.

Since multiple ZigBee channels could be used within the same sub-BBN, the channel allocation problem is relaxed in the WBAN stage and the aforementioned operations are processed indifferently for each ZigBee link $h \in \mathcal{L}^z(t)$, omitting the last operation (iv.), except some restrictions on the available ZigBee channels. Indeed, for each sub-BBN provided with WiFi channel c_1 , we should delimit the set of available ZigBee channels $C^z(h)$ eliminating those that overlap with c_1 :

$$C^z(h) = \{c \in C^z : a_{cc_1} = 0\} \quad \forall (l, h) \subset \mathcal{L}(t), c_1 \in C^w : x_{c_1}^l = 1$$

Hence, the algorithm calculates the set of available ZigBee channels for each sub-BBN (step 17), as well as the list of free ZigBee channels (step 19), which is computed with respect to the set $C^z(h)$.

$$C_{free}^z(h) = \{c \in C^z(h) : \forall k \in \mathcal{Z}_h \cap C^z, y_c^k = 0\}$$

Finally, the ZigBee channel c_2 is computed similarly to the WiFi part (step 23, 24), as follows:

$$c_2 = \begin{cases} \text{Rand}(C_{free}^z(h)), & \text{if } C_{free}^z(h) \neq \emptyset \\ \text{Rand}\{c \in C^z(h) : SIR^z(y_c^h) > SIR_{th}^z\}, & \text{otherwise.} \end{cases} \quad (33)$$

We also keep the condition on the fair sharing of resources, so that a WBAN should release his ZigBee channel after at most θ_s .

Although the proposed SORT-SIM algorithm does not provide the optimal solution for SIM game, it guarantees, at the worst cases, an appropriate strategy with feasible SIR value, i.e. $SIR > SIR_{th}$, while reducing the probability to select the same channel by neighboring WBANs. Furthermore, the simplicity of implementation of SORT-SIM algorithm is a major feature for such highly constrained BBN environment.

7. Performance evaluation

This section illustrates and discusses the numerical results obtained in different network scenarios of both algorithms BR-SIM and SORT-SIM, which have been implemented using the Scilab software package [34]. Then, we compare our algorithms with two existing power control approaches [8,35], which handle almost the same problem we tackle in this work, i.e., the interference mitigation for nearby WBANs.

The mobile WBANs, which number varies in the range [20,50], are randomly deployed in a $1000 \times 1000 \text{ m}^2$ area, and grouped into four overlapping BBNs. The mobility is simulated using the common *random way-point model* [36] (Fig. 7). We consider the first five overlapping WiFi channels of the ISM band ($C^w = \{1, 5\}$) and the whole band of ZigBee channels ($C^z = \{11, 26\}$) in order to simulate the WiFi mutual interference and the cross-technology scenarios. To compute channel gains, we refer to the BBN-specific channel gain model in [23]. The WiFi and ZigBee transmission powers are set to 100 mW and 1 mW, respectively. To prove and compare the effectiveness of our two distributed solutions, we successively evaluate the effect of the WBANs density on the dynamics of the BR-SIM channel selection algorithm and then on the performance of the SORT-SIM algorithm. More specifically, we evaluate the WiFi and ZigBee signal-to-interference ratios for each BBN, proving that the BR-SIM algorithm guarantees a fair sharing of wireless resources, while SORT-SIM presents quickness benefits in some BBN scenarios. SIR^w and SIR^z , in Eqs. (5) and (22), respectively, are indeed our original utility functions that are obtained after the computation of the WiFi and ZigBee Interference Functions.

7.1. BR-SIM versus SORT-SIM

The curves in Figs. 8 and 9 illustrate, respectively, the dynamics of the BR-SIM algorithm for different BBN densities, namely for the number of WBANs $N = 20$ and $N = 40$. More specifically, Fig. 8a and Fig. 8b show the average WiFi SIR and ZigBee SIR, respectively, for $N = 20$. Fig. 8c further shows the convergence of the SIR at the ZigBee interface of a subset of players under the BR-SIM algorithm. Similarly, Fig. 9a, Fig. 9b and Fig. 9c display, respectively, the evolution of the average SIR and the actual SIR values for a subset of players by each BBN, so as to show the effect of the network density on the convergence of the BR-SIM algorithm. As expected, increasing the BBN density results in increasing the network overall interference and the number of iterations to reach an equilibrium.

Besides, we notice at the Nash Equilibrium that the worst WiFi SIR (21 dB for $N = 20$ and 9 dB for $N = 40$), measured with the standard transmission power of 20 dBm (100 mW) is always above the receiver sensitivity of most commercial cards (the lowest receiver sensitivity for the Atheros chipset is -95 dB), even considering other effects like fading and thermal noise. The same conclusions are observed for the worst ZigBee SIR measured by all four BBNs (i.e., the WBAN that experiences the worst SIR in a BBN), which varies between 25 and 30 dB for $N = 20$ and $N = 40$ respectively. Note that the worst SIR measured

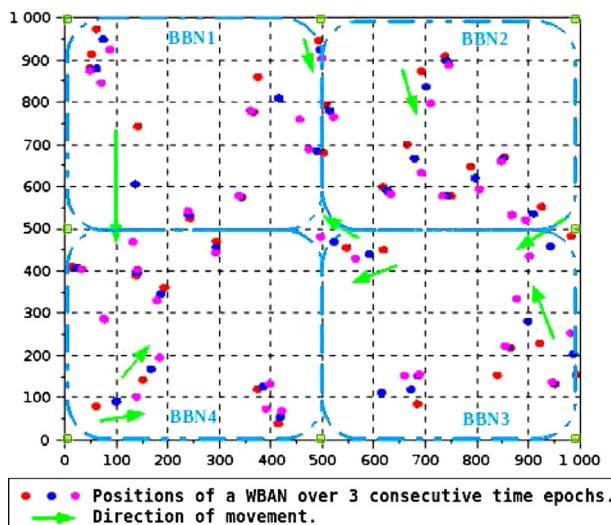


Fig. 7. Simulation scenario for $N = 40$ WBANs.

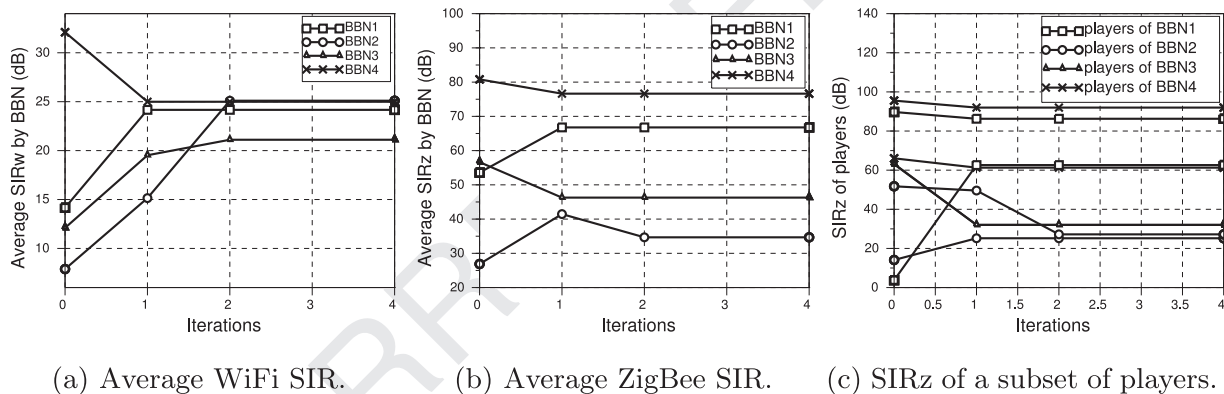


Fig. 8. Dynamics of the BR-SIM algorithm for each BBN, with $N = 20$ WBANs.

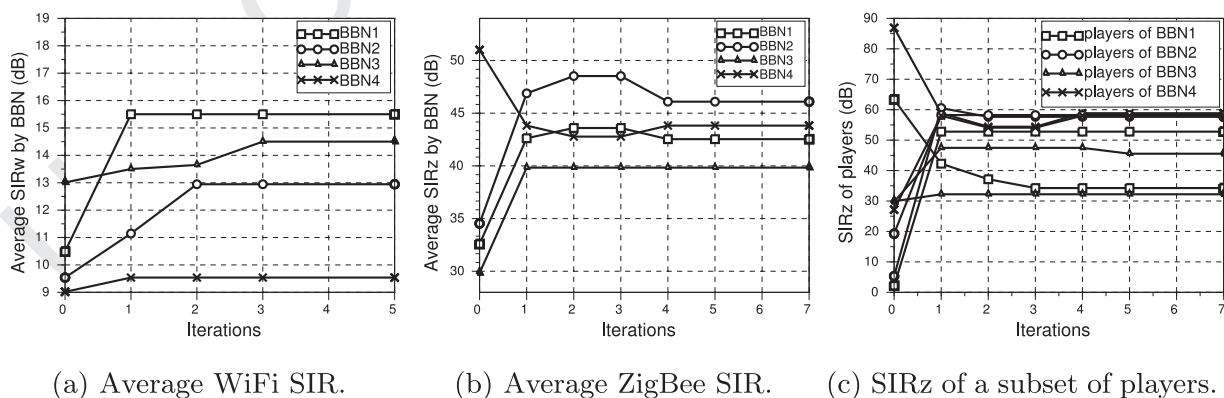


Fig. 9. Dynamics of the BR-SIM algorithm for each BBN, with $N = 40$ WBANs.

1042 at the ZigBee interface is higher than the value measured
 1043 at the WiFi interface due to the restricted number of
 1044 overlapping WiFi channels used in the simulation in order
 1045 to enable mutual and cross-technology interference, thus
 1046 resulting in conflicting transmissions using the WiFi tech-

1047 nology. Naturally, within a BBN only WiFi transmissions
 1048 coming from surrounding BBNs are considered in the
 1049 computation of the WiFi interference, since we assume the
 1050 utilization of a coordination scheme for intra-BBN commu-
 1051 nications, whereas the ZigBee interface of any WBAN

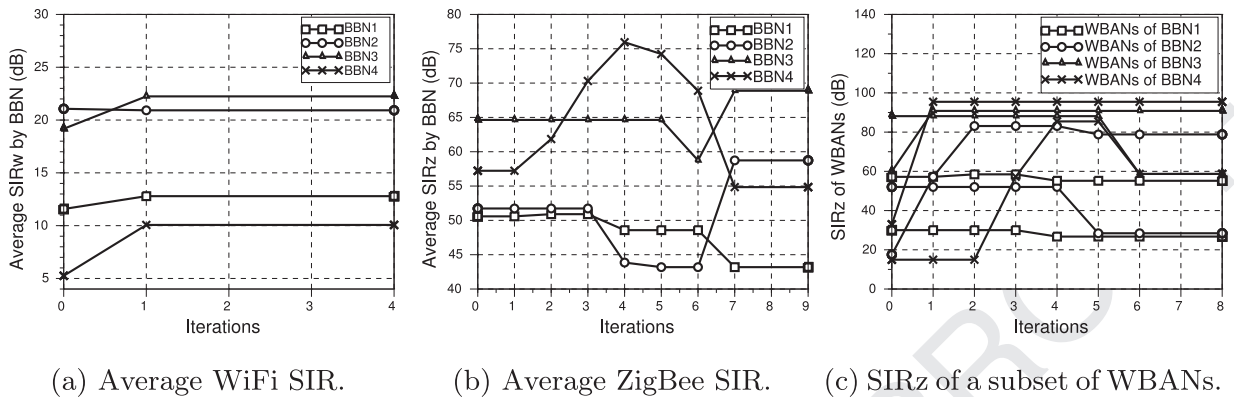


Fig. 10. Iterations of the SORT-SIM algorithm for each BBN, with $N = 20$ WBANs.

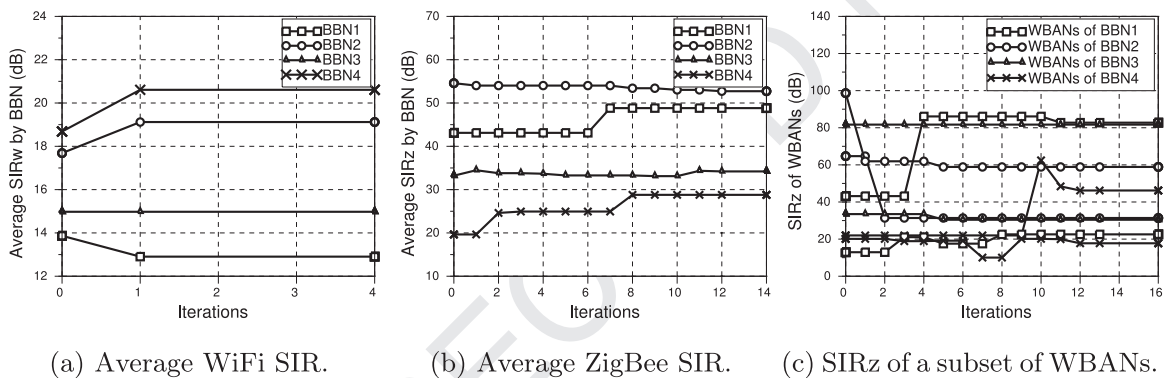


Fig. 11. Iterations of the SORT-SIM algorithm for each BBN, with $N = 40$ WBANs.

1052 experiences both intra-BBN and inter-BBN interference.
 1053 Thereby, further experiments with non-overlapping WiFi
 1054 channels would reverse the previous conclusions and
 1055 assess higher values of WiFi SIR versus ZigBee SIR.

1056 Yet, the performance of BR-SIM is ensured since it
 1057 provides a rather fair, socially-aware channel allocation,
 1058 so that both WiFi and ZigBee signal-to-interference ratios
 1059 tend to be quite close to a mean value at the Nash Equi-
 1060 librium. Nevertheless, a noticeable decrease in the range
 1061 of SIR values (mainly SIRz), at the NE point, is observed
 1062 when the density of the WBANs is high ($N = 40$), and the
 1063 SIR curves are tightly close. Indeed, higher densities occa-
 1064 sion a more fair spreading of players over the neighboring
 1065 BBNs, that will suffer from relatively fair interference envi-
 1066 ronment. This explains why, for lower densities, the aver-
 1067 age SIR values for each BBN are spread out over a larger
 1068 range of values.

1069 On the other hand, Figs. 10 and 11 illustrate the signal-
 1070 to-interference ratios at WiFi and ZigBee interfaces ob-
 1071 tained by the SORT-SIM algorithm for the same topology
 1072 configurations (i.e., $N = 20$ and $N = 40$). Almost the same
 1073 conclusions can be made for SORT-SIM, as far as BR-SIM
 1074 results, in terms of the evolution of SIR metrics as a func-
 1075 tion of WBANs density, wherein we can observe the degra-
 1076 dation of both WiFi and ZigBee SIR values while increas-
 1077 ing the BBN density. However, if we observe the average SIR
 1078 of the whole network we can notice the main differences be-

1079 tween the behavior of the two algorithms. Indeed, Fig. 16a
 1080 and b shows a more accentuated steepness of SORT-SIM
 1081 curves compared to that of BR-SIM, which means that the
 1082 effectiveness of SORT-SIM is more density-sensitive, while
 1083 BR-SIM seems to be more robust to density changes. In fact
 1084 with higher densities, i.e., beyond $N = 30$ players, SORT-
 1085 SIM presents more severe degradation in SIR values for
 1086 both WiFi and ZigBee transmission links, whereas BR-SIM
 1087 shows a smooth decrease while preserving good SIR ratios.

1088 Now, if we observe the performance of each algorithm
 1089 separately, we notice rather similar behaviors at low
 1090 densities (Figs. 8 and 10), where few players are spread
 1091 out over the simulation area. Both algorithms compete in
 1092 allocating feasible, near optimal, WiFi and ZigBee chan-
 1093 nels to all players. However, for high densities we notice
 1094 that BR-SIM curves merge around the average SIR, while
 1095 SORT-SIM still presents great divergences among players'
 1096 SIR values. This can be explained by the usefulness of the
 1097 cooperative component of BR-SIM, where the local interac-
 1098 tions among neighbors allow it to fairly share the wireless
 1099 resources. Whereas, SORT-SIM proceeds in a completely
 1100 non-cooperative manner, thus some players get maximal
 1101 SIR values, while others settle for channel allocations with
 1102 minimal SIR values, just above the threshold.

1103 Yet, the SIR values at both WiFi and ZigBee inter-
 1104 faces under the BR-SIM and SORT-SIM algorithms are il-
 1105 lustrated in detail in Figs. 12 and 13, respectively. More

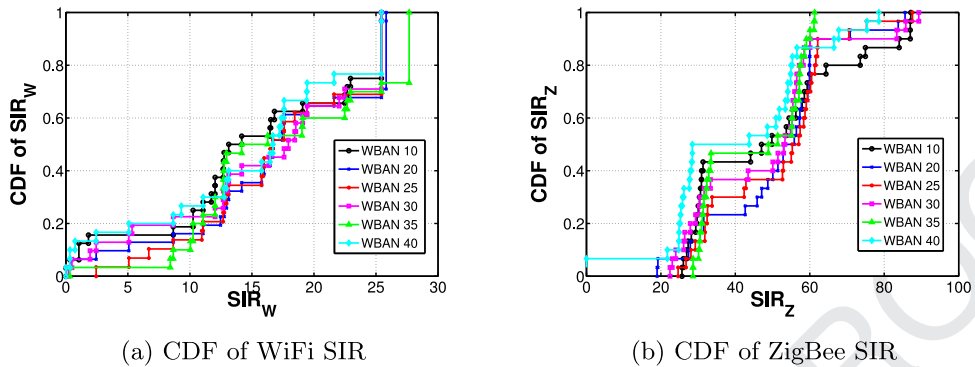


Fig. 12. BR-SIM: empirical Cumulative Distribution Function (CDF) of the SIR measured at WiFi and ZigBee interface of all WBANs in the BBN scenario of 40 WBANs with 30 time epochs of 10 s each.

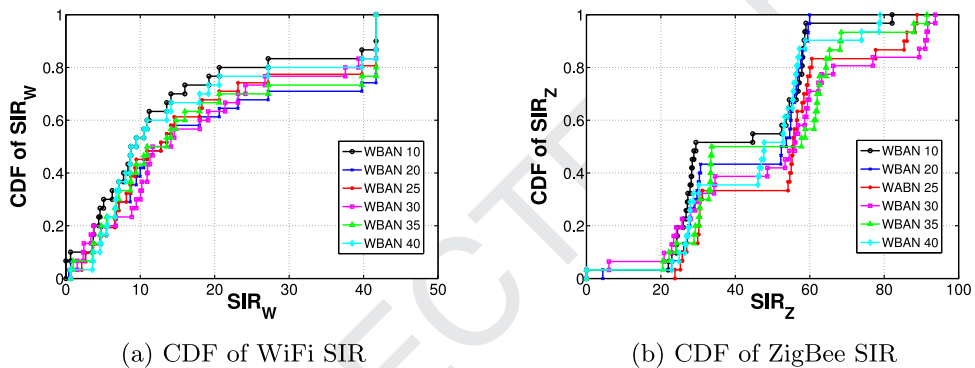


Fig. 13. SORT-SIM: empirical Cumulative Distribution Function of the SIR measured at WiFi and ZigBee interface of all WBANs in the BBN scenario of 40 WBANs with 30 time epochs of 10 s each.

1106 specifically, these figures show the empirical Cumulative
 1107 Distribution Function (CDF) of the SIR when the total number
 1108 of WBANs $N = 40$ and for a time duration of 300 s,
 1109 which is divided in 30 time epochs of 10 s each. Let us
 1110 first focus on the SIR metric for WiFi obtained with BR-
 1111 SIM (Fig. 12a) and SORT-SIM (Fig. 13a). It can be observed
 1112 that the SIR values under both algorithms are quite similar
 1113 and range from 0 to ≈ 40 dB. However, it is not hard to
 1114 see that BR-SIM guarantees for the majority of the players
 1115 fair values of SIR (in the range [10,25]), while SORT-SIM
 1116 performs WiFi channel assignment to transmission links in
 1117 a much more aggressive way, where some players enjoy
 1118 high values of SIR while others suffer from very low val-
 1119 ues. Similarly, for the SIR value measured at the ZigBee
 1120 interface, Figs. 12b and 13 b show that in more than 50% of
 1121 the scenarios, the SIR is higher than approximately 50 dB.
 1122 However, note that in the case of SORT-SIM and for the
 1123 6 considered WBANs the percentage of players getting a
 1124 value of SIR below 20 dB is larger than the one obtained
 1125 with BR-SIM. Hence, this trend confirms the fact that BR-
 1126 SIM guarantees at the same time some fairness along play-
 1127 ers and good performance.

1128 Besides, we calculate with Scilab the computation time
 1129 (CPU time) for both algorithms and we find noticeable dif-
 1130 ference between them. Indeed, the BR-SIM computation
 1131 time is about four times larger than that of the SORT-SIM
 1132 execution instance. For example, the maximum computa-

1133 tion time we measured to solve the BR-SIM algorithm over
 1134 30 consecutive time epochs was approximately equal to
 1135 1060 s, for $N = 50$ WBANs. Conversely, SORT-SIM takes
 1136 less than 228 s to find the sub-optimal solutions for the
 1137 SIM problem, under the same network instances and pa-
 1138 rameters' settings. Furthermore, it can be observed that the
 1139 BR-SIM algorithm converges to a stable operational point
 1140 in few iterations, in particular, all BBNs converge to their
 1141 best WiFi and ZigBee channel allocations in at most 3 and
 1142 5 iterations, respectively, while SORT-SIM performs with
 1143 greater number of iterations (up to 15), but within less
 1144 computation time.

1145 Finally, BR-SIM outperforms in terms of fairness and ro-
 1146 bustness the SORT-SIM algorithm, especially at higher den-
 1147 sities, thus representing a practical solution for interfer-
 1148 ence mitigation in realistic BBN scenarios. However, SORT-
 1149 SIM presents simplicity and rapidity benefits which makes
 1150 it useful, under specific BBN scenarios, mainly at low den-
 1151 sities and low QoS requirements.

7.2. Comparison with power control approaches

1152
 1153 In this section, we compare our BR-SIM and SORT-SIM
 1154 algorithms to the distributed power control algorithm pro-
 1155 posed in [8] and to the joint relay selection and transmit
 1156 power control algorithm proposed in [35].

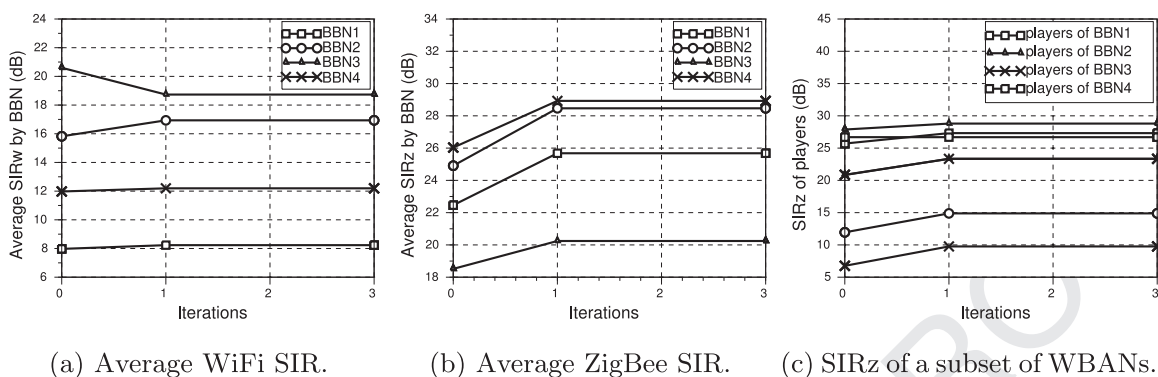


Fig. 14. Dynamics of the PAPU algorithm for each BBN, with $N = 40$ WBANs.

1157 Authors in [8] formulated a power control game con-
 1158 sidering interference between neighboring WBANs and
 1159 energy-efficiency. They derived a distributed power control
 1160 algorithm, called the ProActive Power Update (PAPU) algo-
 1161 rithm, to reach a unique Nash Equilibrium (NE) represen-
 1162 ting the best tradeoff between energy-efficiency and net-
 1163 work utility. As in our model, PAPU assumes a TDMA-based
 1164 MAC protocol to deal with intra-WBAN interference avoid-
 1165 ance, and uses the SINR metric to define the utility func-
 1166 tion of the power control game. However, neither WBAN
 1167 mobility is considered, nor wireless technologies are
 1168 specified.

1169 Alike our SIR metrics defined in our paper by expres-
 1170 sions (5) and (22), respectively, for WiFi and ZigBee re-
 1171 ceived signals, the SINR was defined in [8] without con-
 1172 sideration of heterogeneous wireless technologies. This will
 1173 be reflected in the final SINR values, as we will show
 1174 hereafter.

1175 Indeed, we have implemented the PAPU algorithm with
 1176 the same network configuration of our BR-SIM and SORT-
 1177 SIM algorithms, and with the following definition of the
 1178 power best-response performed by each WBAN/player:

$$b_i(p_{-i}) = \frac{1}{c_i} - \frac{\sum_{j \neq i} h_{ji} p_j + n_0}{h_{ii}} \quad (34)$$

1179 where p_j is the transmission power of player j , h_{ji} rep-
 1180 represents the channel gain between transmitter j and re-
 1181 ceiver i , h_{ii} the intra-network gain, n_0 is the background
 1182 white noise power (which is ignored in our simulations
 1183 since we calculate the SIR), and c_i the power price. The ob-
 1184 tained (average) SIR values are reported in Figs. 14 and 16.

1185 First, it can be observed from Fig. 14 that PAPU is rather
 1186 efficient with respect to WiFi SIR maximization; results are
 1187 almost in the same range as the BR-SIM and SORT-SIM al-
 1188 gorithms. This can be explained by the fact that PAPU's
 1189 WiFi SIR does not consider the cross-technology inter-
 1190 ference from ZigBee on WiFi links. Only intra-WBAN chan-
 1191 nel gains are involved, whereas in real BBN scenarios the
 1192 cross-technology channel gains introduce further inter-
 1193 ference components to the SIR denominator.

1194 However, the difference mainly appears in the second-
 1195 stage game (Fig. 16b), where PAPU provides less efficient
 1196 SIR values for the ZigBee signal. Whilst BR-SIM and SORT-
 1197 SIM provide ZigBee SIR values over 20 dB (up to 80 dB),

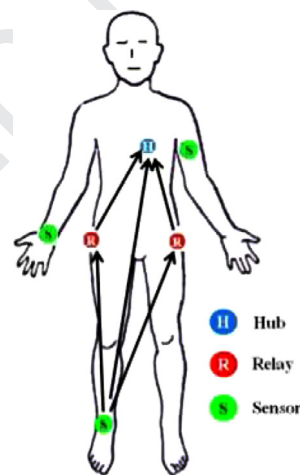


Fig. 15. WBAN configuration for the RSPC algorithm [35].

1198 PAPU's maximum ZigBee SIR is around 20 dB (up to 1199
 1200 40 dB for lower network densities). Yet, as its authors ex- 1201
 1202 plained, PAPU requires limited information exchange be- 1203
 1204 tween WBANs, and as a consequence the player strategy 1205
 1206 is purely selfish, without any consideration of neighboring 1207
 1208 WBANs' utilities. With local interactions of our SIM game, 1209
 1210 BR-SIM and SORT-SIM achieve better SIR values, and thus 1211
 1212 stronger wireless signal. This also explains the regularity 1213
 1214 of PAPU curves, whereas the negotiations among players 1215
 1216 are better observed on the BR-SIM and SORT-SIM curves. 1217
 1218 It is worth noting that the reduced number of iterations 1219
 1220 of the PAPU algorithm within our network configuration, 1221
 1222 compared to that of the original paper, is also due to the 1223
 1224 local interaction behavior among players, which allows a 1225
 1226 rapid convergence to the NE.

1227 We now compare BR-SIM and SORT-SIM to the *joint Re- 1228
 1229 lay Selection and transmit Power Control algorithm*, referred 1229
 1230 to hereafter as *RSPC algorithm*, proposed in [35]. 1231

1232 In [35], each WBAN has the following configuration (see 1233
 1234 Fig. 15): a hub at the chest, two relays at the right and left 1235
 1236 hips, and three sensors at other suitable locations. The hub, 1237
 1238 the sensor and the two relays are denoted as H , S , R_1 and 1239
 1240 R_2 , respectively. Time division multiple access (TDMA) and 1241
 1242 asynchronous TDMA are respectively used as intra- and 1243

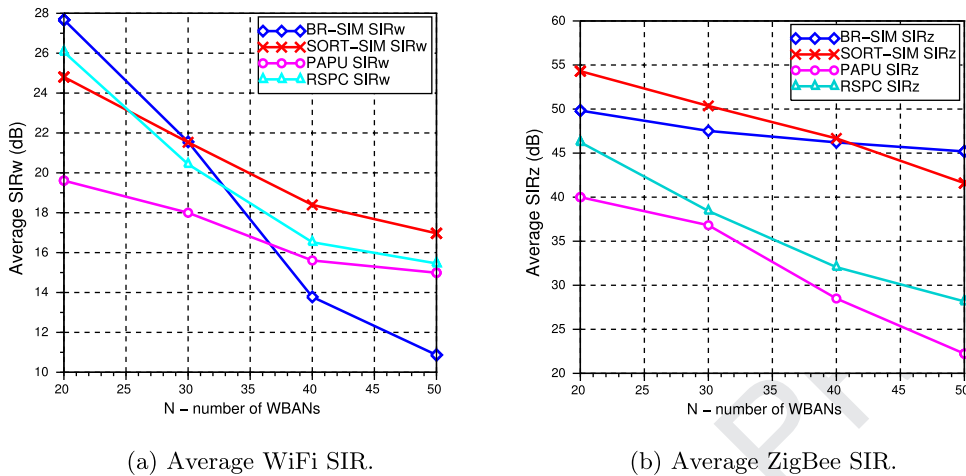


Fig. 16. BR-SIM and SORT-SIM vs PAPU and RSPC. Average WiFi and ZigBee SIR as a function of network density.

inter-WBAN access schemes, since it has been shown in [37] that they provide better interference mitigation than other access schemes in terms of power consumption and channel quality.

The major contribution of the RSPC algorithm is the use of opportunistic relaying with no cooperation between WBANs to provide inter-body channel gain measurements, in order to improve reliability (decrease the outage probability) and reduce the power consumption. RSPC uses the on-body and inter-body channel data sets in [38], obtained through exhaustive scenarios performed in realistic environments, over several hours of normal everyday activities. In each experiment, sensors transmit in a round-robin fashion with 5 ms separation between each other.

Thereby, the RSPC algorithm can be summarized in the three following steps:

1. Power control at the sensor level: the sensor performs power control on a channel at time epoch τ using the value at time epoch $\tau - 1$, and selects the one-hop relay: StoH (Sensor-to-Hub), StoR1 (Sensor-to-Relay1) or StoR2 (Sensor-to-Relay2).
2. Power control at the relay level: select the relay transmit power to the hub, in the transmit range.
3. Branch selection at the hub: the hub selects the path (StoH, StoR1-R1toH or StoR2-R2toH) that gives the best SINR.

The authors in [35] assert that relay-assisted communications can reduce co-channel interference from neighboring WBANs, by increasing the SINR of the packets transmitted by the sensor node and received at the WBAN coordinator (the hub/the MT in our model), expressed by:

$$\text{SINR} = \frac{T_x \times |h_{TxRx}|^2}{\sum T_{x_{int,i}} |h_{int,i}|^2} \quad (35)$$

where T_x is the sensor/relay transmit power obtained by the Power Control function (step 1 or 2 of the RSPC algorithm). $|h_{TxRx}|$ represents the average channel gain across the duration of the sensor/relay transmitted signal, while $|h_{int,i}|$ is the channel gain between the interferer int , which

is the neighboring WBAN sensor, and the sensor or selected relay i . Finally, $T_{x_{int,i}}$ denotes the interfering power of neighboring WBAN sensor int to the sensor/relay i . The instantaneous noise at the receiving node has been omitted, since we compare SIR metrics.

For the one-hop relay selection, we consider the WBAN configuration given in Fig. 15. Since TDMA is used as access scheme, sensors cannot transmit simultaneously within a WBAN. Yet, to adapt the RSPC algorithm to our network model, we focus on a WBAN's sensor-of-interest, and we assimilate the neighboring interferer sensor to its corresponding MT. The one-hop relay process will be considered while selecting the intra-WBAN transmit power, i.e. in the ZigBee stage. We further assume that WBANs use a WiFi channel for inter-WBAN exchanges. Power control will also be performed for WiFi transmissions in a way to maximize the MT WiFi SIR, using the ZigBee power vectors of neighboring WBANs, computed at the previous time epoch.

We run our simulations and we calculate the WBAN's SIR (SIRw and SIRz), considering the aggregate interference due to transmit powers of the neighboring WBANs.

It can be observed from Fig. 16a that, in general, the RSPC WiFi SIR curve lies between BR-SIM and SORT-SIM curves. Even though RSPC does not perform iterations to reach the best SIR, unlike the game models, it optimizes once the sensor/relay transmit power with its Power Control algorithm and achieves rather efficient SIR values.

These results can be explained by analyzing, as we do hereafter in Fig. 17, the aggregate interference, calculated as the sum of interference suffered by the hub/MT, due to WiFi and ZigBee transmissions of neighboring WBANs.

In Fig. 17, we notice an important gap between the RSPC aggregate interference and the one obtained by our algorithms (BR-SIM and SORT-SIM) and PAPU. Specifically, IN_{BR-SIM} and $IN_{SORT-SIM}$ are always lower than those of PAPU and RSPC, even though sometimes the WiFi SIR of RSPC is higher than the one achieved by BR-SIM or SORT-SIM (Fig. 16a). This can be explained as follows:

- The aggregate interference values of the BR-SIM and SORT-SIM algorithms are considerably lower than

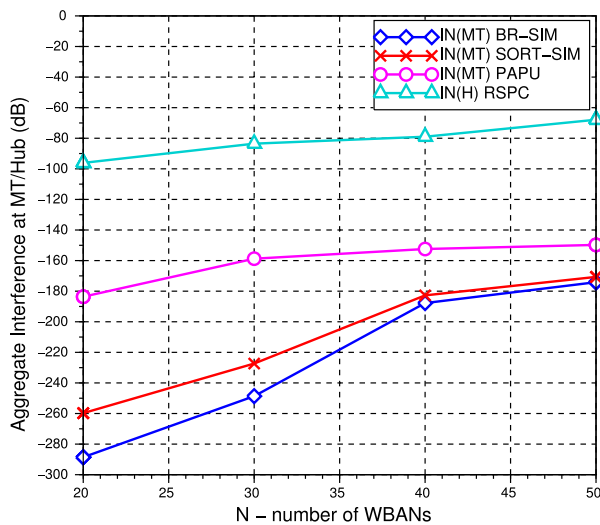


Fig. 17. Aggregate interference at the Hub/MT.

those of PAPU and RSPC, because in our interference mitigation model we assign WiFi/ZigBee channels to wireless links in a way to reduce the co-channel and cross-interference components. Therefore, neighboring interfering WiFi/ZigBee links are omitted (by allocating them orthogonal channels) or reduced by the w_{mm} scalar, to ensure minimum mutual interference.

- The gap is less important for the SIR values, because the MT/Hub channel gains and transmit powers are far larger than the interference component in the four algorithms, either with power control (PAPU and RSPC), or with constant transmit power (BR-SIM and SORT-SIM). Indeed, the four algorithms achieve efficient interference mitigation, ensuring feasible SIR values. However, the advantage of BR-SIM and SORT-SIM mainly appears when we compare the aggregate interference (Fig. 17) and the ZigBee SIR (Fig. 16b). This can be explained by the fact that our algorithms give some privilege to ZigBee links w.r.t. WiFi links; WiFi interference on ZigBee links is considered more crucial than ZigBee interference on WiFi links. In other words, our algorithms make sure that WiFi links (which use a transmit power 100 times higher than that of ZigBee) will not prevent ZigBee transmissions and deteriorate the BBN system performance.

Although the aggregate interference $IN(MT)$ of BR-SIM and SORT-SIM is significantly lower than that of RSPC $IN(H)$, it increases more rapidly for higher densities, because the use of orthogonal channels is no more possible, and BR-SIM and SORT-SIM start using channels with minimum mutual interference, with constant WiFi and ZigBee powers. However, RSPC maintains approximately the same level of interference by adjusting the transmit power of the sensor/relay nodes. Hence, it would be interesting in future work to consider a control power mechanism together with the channel assignment to further improve the efficiency of the SIM game.

8. Conclusion

In this paper we studied the distributed interference mitigation problem in BBN scenarios from a game theoretical perspective. In particular, our work made three main contributions. First, we formulated the problem as a game considering the SIR, which accurately models the channel capacity that can be achieved in the presence of mutual and cross-technology interference. Second, we studied the properties of our game proving the existence of a Nash Equilibrium, which represents channel allocations that minimize the mutual and cross-technology interference. Third, we proposed a two-stage algorithm (called BR-SIM) based on the best-response dynamics to compute the Nash Equilibria in a distributed fashion. We further developed an alternative approach (SORT-SIM) that reaches a sub-optimal solution in less computational time than BR-SIM. Finally, we evaluated and compared our SIM game theoretical approaches to (relay-assisted) power control schemes (i.e., PAPU and RSPC) in realistic BBN scenarios. We first showed that the BR-SIM algorithm converges quickly and achieves feasible values for the utility functions, while SORT-SIM presents some practicability benefits under specific network scenarios. Then, we demonstrated that BR-SIM and SORT-SIM outperform PAPU and RSPC in terms of SIR and Aggregate Interference in several cases, and especially when the network density is quite low.

Besides, numerical results we gathered in the present work show that BBN scenarios require the definition of distributed scheduling algorithms to avoid simultaneous transmissions that might affect the channel quality and completely prevent communications among network nodes.

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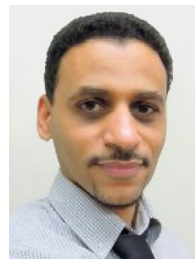
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