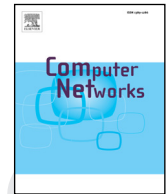




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Energy efficient power allocation in cognitive radio network using coevolution chaotic particle swarm optimization

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ABSTRACT

In this paper, the trade-off between utility and energy consumption in orthogonal frequency division multiplexing (OFDM)-based cognitive radio (CR) network is investigated. Energy efficiency problem is very important in the field of CR network, where the utility is maximized and the energy consumption is minimized in such a CR network. Since the trade-off between them has been paying more attentions in literature, this study summarizes the power allocation as an optimization problem that maximizes the energy efficiency via a new energy efficiency metric defined by this paper. The formulated problem is a large-scale nonconvex problem, which is very difficult to solve. In this paper, we present an improved particle swarm optimization (PSO) algorithm to solve the difficult large-scale optimization problem directly. Given the weak convergence of the original PSO around local optima, an improved version that combines the chaos theory is proposed in this study, where chaos theory can help PSO search for solutions around the personal and global bests. In addition, for the purpose of accelerating the convergence process when facing with such a large-scale optimization, the original problem is decomposed into a number of small ones by employing the coevolutionary methodology, and then divide-and-conquer strategy is used to avoid producing infeasible solutions. Simulations demonstrate that the proposed coevolution chaotic PSO needs a smaller number of iterations and can achieve more energy efficiency than the other algorithms.

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1. Introduction

Nowadays almost all wireless spectrum has been licensed to existing wireless communications applications. With the increasing demand for wireless data service, spectrum scarcity will become a big problem in future development of wireless communications networks. One promising solution to overcome the spectrum scarcity problem is to use opportunistic spectrum access techniques such as cognitive radio [1], which lets unlicensed users (called secondary users or cognitive users) temporarily utilize a licensed spectrum band, if the licensed users

(called primary users) are idle or the interference received at primary users from secondary transmissions is tolerable (in other words, secondary transmissions do not affect the transmission quality of primary users). Due to its potential to largely improve the spectrum utilization efficiency, CR has received much attention from academia, industry, and spectrum regulation agencies [2].

In a cognitive radio (CR) network, primary users have the highest priority to use the spectrum. Secondary users are aware of the transmission environments, and can adapt their transmission/reception patterns to the varying spectrum environments. As an example, consider that a secondary user uses a licensed spectrum band. If the corresponding primary user is back (i.e., the primary user needs to use the spectrum band), the secondary user needs to stop using the spectrum band, and try to find

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other idle spectrum bands to continue its wireless access service. Dynamic spectrum allocation is a main challenge in the design of cognitive radio networks, which enables wireless devices to opportunistically access portions of the spectrum as they become available.

In this paper, we consider an orthogonal frequency division multiplexing (OFDM)-based CR network, and we focus on the network utility (to be defined in Section 3). We define “utility per Joule” as the energy efficiency metric, which can effectively characterize the trade-off between utility and energy. A power allocation problem is formulated, which maximizes the energy efficiency. Since there are base stations (BSs) in the system, the optimization problem is centralized and nonconvex, and is hard to be transformed to a convex problem. So in this paper, we adopt PSO (particle swarm optimization) algorithm, which can solve the nonconvex optimization efficiently. PSO algorithms are modern heuristic algorithms based on bird flocking, there is no theory proof for PSO to get the global optimum, but they have demonstrated their potential in solving complicated optimization problems [3–6] and network optimization problems [7–9]. The advantages of PSO algorithms include: they have simple theoretical structure with good convergence properties; they are easy to implement; they do not require the objective functions to be continuous. PSO methods have been popularly used in wireless networks. For example, Zhao et al. [7] uses PSO to optimize CR parameters based on the spectrum environments and user needs; a PSO-based distributed resource allocation algorithm in wireless mesh networks is proposed in Ref. [8]; and Lin [9] applies PSO to deal with router node placement problem in a dynamic wireless mesh network such that the network connectivity and client coverage are maximized.

However, it is very likely that traditional PSO algorithms may be trapped into local optimal solutions (which are not global optimal). Therefore, in the literature, chaos, which has the features of randomness, ergodicity and regularity, has been used in PSO algorithms recently [10–14]. Chaotic PSO algorithms can maintain the population diversity, which is a nice property. Liu et al. [10] applies chaotic dynamic in PSO algorithms, using the chaotic local searching behavior. Coelho and Herrera [11] considers fuzzy identification, which enhances PSO algorithms with chaotic Zaslavskii map sequences and efficient Gustafson–Kessel clustering. The chaotic PSO algorithm is shown to be effective in building a good TS fuzzy model. In [12], the authors consider prediction of silicon content in hot metal, in which PSO algorithms are enhanced with chaotic under the logistic equation. A binary PSO is used in [13] to predict operon in bacterial genomes, and chaotic sequence is introduced when updating inertia weight. In [14], the authors apply a PSO algorithm to estimate the unknown parameters for a hybrid-forecasting model, in which initial values of unknown constants in particle velocity and position equations are generated by chaotic mapping. Due to the nice features of chaos theory in PSO algorithms, we adopt a chaotic PSO algorithm in this paper.

Besides the chaos, we also apply the cooperative coevolution theory, since cooperative coevolution theory is very suitable for large scale optimization problems. An applica-

tion of cooperative coevolution theory in PSO can be found in [15], in which PSO position update rule relies on Cauchy and Gaussian distributions. And in our recent paper [16], chaos theory is combined into cooperative coevolving PSO, as chaos theory can help PSO search for solutions around the personal and global bests, thus avoiding being trapped into local optimal points. And a belief space is used to store the experiences for individuals to learn from each other indirectly. In this paper for CR networks, coevolution and chaos theory are all combined with PSO, referred to as CCPSO. But there is no need to set a belief space. Two populations of PSO are included in the CCPSO, and the problem is solved by using max–min approach.

The rest of the paper is organized as follows. Related work is given in Section 2. The system model and problem description are presented in Section 3. The proposed CCPSO algorithm is given in Section 4. Numerical results are provided in Section 5, followed by conclusions in Section 6.

2. Related work

The utility maximization in a multi-cell CR network under a total transmit power constraint is considered in [17]. A cooperative secure resource allocation in CR Networks was considered in [18], since the problem is NP hard, the problem is transformed into a generalized geometric programming model. Since secondary users are usually powered by battery, energy consumption in their wireless transmissions is an important issue. Further, large energy consumption is often due to large transmission power, which actually generates large interference to users in the vicinity and degrades service of those users. Accordingly, in this paper, we consider energy efficient CR networks that employ orthogonal frequency division multiplexing (OFDM) technology [19]. The reason we consider OFDM is that OFDM is very suitable for high speed broadband wireless access due to its immunity to inter-symbol-interference (ISI).

Energy efficiency has been well-investigated in traditional wireless networks [20,21]. Code division multiple access (CDMA) networks are considered in [22], which develops a cross-layer algorithm for energy efficiency. It is proved that the algorithm is Pareto-optimal under certain conditions. Meshkati et al. [23] also considers a CDMA network, which studies the trade-off between energy efficiency and delay. A game theoretical approach is presented to maximize the utility by selecting the transmit power under a delay requirement.

Due to the popularity of CR research, energy efficiency in CR has also received a lot of attention. Buzzi and Saturnino [24] consider a cognitive CDMA wireless network, and presents a game-theoretic algorithm to achieve energy efficiency in a one-shot fashion. Wu and Tsang [25] investigates the sensing and transmission durations of secondary users. A nonconvex optimization problem is formulated to achieve energy-efficient power allocation, which is solved by analyzing three special cases. In [26], the authors develop an energy efficient power control algorithm for OFDM-based CR networks. The objective function is based on the “throughput per Joule” metric. The formulated optimization problem, which is nonconvex, is transformed

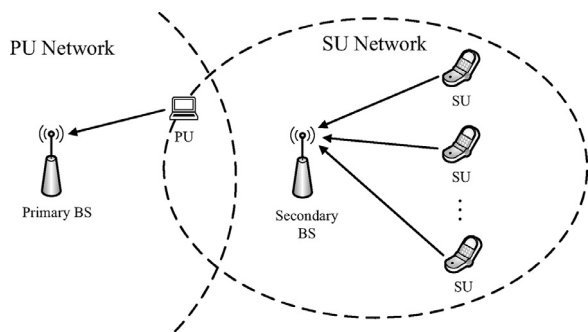


Fig. 1. The network system.

148 into a convex problem based on parametric programming.
 149 Mao et al. [27] considers multiple-input multiple-output
 150 cognitive systems and achieves energy efficient spectrum
 151 optimization under three constraints (on the total power,
 152 the interference power, and the minimum system through-
 153 put). The formulated optimization problem, which is also
 154 nonconvex, is transformed to a one-dimension problem
 155 that has a quasi-concave objective function. Zhong and
 156 Wang [28] uses “rate per Joule” as the energy efficiency
 157 metric, which stands for the number of information bits
 158 that are successfully transmitted per Joule of energy over
 159 a normalized bandwidth.

160 We consider network utility as sigmoid form, which is
 161 more suitable for a practical network. Accordingly, we de-
 162 fine energy efficiency as “utility per Joule”, to character-
 163 the trade-off between the utility and the energy consump-
 164 tion. The formulated optimization problem is nonconvex
 165 and has a large scale, which is difficult to solve. Previous
 166 optimization works usually try to transform nonconvex
 167 problems to convex ones based on some assumptions or
 168 try to get suboptimal solutions. In this paper, we develop
 169 power allocation algorithms based on an improved CCPSO.
 170 We solve the difficult large-scale optimization problem
 171 directly, without any assumptions or transforms. In this
 172 paper, chaos technique is combined into PSO, whose
 173 randomizing and erogeneity characteristics can help avoid
 174 being trapped into local optimal points. Coevolution idea
 175 is used as well to PSO, to exploit its nice feature in dealing
 176 with large scale problems. The proposed algorithm has two
 177 populations, and we solve the problem by using max–min
 178 approach. Simulation results show that the proposed algo-
 179 rithm can solve the large scale nonconvex power allocation
 180 problem effectively and efficiently. The chaotic initializa-
 181 tion and search can help the algorithm jump out of local
 182 optima and the divide-and-conquer strategy used in the
 183 coevolution theory can help the algorithm avoid producing
 184 infeasible solutions, the statistical test results reveal that
 185 the proposed method outperforms other existing methods
 186 and has stronger robustness than other methods.

187 3. System model and problem description

188 3.1. System model

189 As shown in Fig. 1, consider a CR network consisting of
 190 N secondary users (SUs) as transmitters and one secondary
 191 base station (BS) as the common secondary receiver. Here

192 uplink communication is considered (downlink communi-
 193 cation can be treated similarly). The secondary transmis-
 194 sers are allowed to use the licensed spectrum band of a
 195 primary transceiver pair: a primary user (PU) as the trans-
 196 mitter and a primary BS as the receiver. OFDM technology
 197 is used in both secondary and primary transmissions. To
 198 avoid degrading the transmission quality of primary users,
 199 it is required that the interference from secondary trans-
 200 mitters to the primary BS is below a threshold denoted I_{th} .

201 The licensed spectrum band is divided into K sub-
 202 channels using OFDM [19]. For the wireless links (desired
 203 communication links and interference links), block fading
 204 model is assumed. In specific, for each wireless link over
 205 each subchannel, its link gain is fixed for a time slot, and
 206 changes independently in the next time slot. For each link,
 207 the link gains over different subchannels are independent
 208 from each other. In this paper, we consider power alloca-
 209 tion of the secondary transmitters in each time slot.

210 For a specific time slot, denote P_k^i as the transmit
 211 power of SU i over subcarrier k ($k = 1, 2, \dots, K$), and G_k^{ii} ($i =$
 212 $1, 2, \dots, N$) as the link gain from SU i to the secondary BS.
 213 Then the signal to interference ratio (SINR) for SU i 's com-
 214 munication over subchannel k , denoted γ_k^i , is given as

$$215 \gamma_k^i = \frac{G_k^{ii} P_k^i}{\sum_{j \neq i} G_k^{jj} P_k^j + \sigma^2}, \quad (1)$$

216 where σ is the background noise.

217 3.2. Utility function

218 For SU i 's communication over subchannel k , next we
 219 define a utility function denoted as U_k^i . The utility function
 220 is expected to reflect the satisfaction level of the service
 221 quality. In this paper, we adopt the sigmoidal form utility
 222 introduced in [30], which means that the user is more and
 223 more satisfied with the service as the quality improves.
 224 The sigmoidal function can capture the value of the service
 225 to the user quite naturally, and be defined as

$$226 U_k^i = \frac{1}{1 + e^{-a_i(\gamma_k^i - b_i)}} \quad (2)$$

227 where a_i is slope parameter (a large a_i means that the ap-
 228 plication has a soft quality-of-service requirement) and b_i
 229 is a shift parameter that is actually the required average
 230 data rate of the application.

231 Targeting at higher utility function with less energy
 232 consumption, we define an energy efficiency metric as the
 233 ratio of the utility to the transmission power, given as

$$234 E = \frac{\sum_{i=1}^N U_k^i}{\sum_{i=1}^N P_k^i + P_c} \quad (3)$$

235 where P_c is static power consumption other than wireless
 236 transmissions (for example, circuit power consumption). In
 237 other words, the energy efficiency metric is actually “utility
 238 per Joule”.

239 3.3. Constraints

240 Recall that over each subchannel, the interference re-
 241 ceived by the primary BS from all secondary transmis-
 242 sions should be bounded by a threshold I_{th} . Denote G_k^{0i} ($k =$
 243 $1, 2, \dots, K$) as the link gain from SU i to the primary BS.

240 1, \dots, K) as the link gain from the SU i to the primary BS.
241 Thus, the interference power constraint is given as

$$\sum_{i=1}^N G_k^{0i} P_k^i \leq I_{th}. \quad (4)$$

242 It is also required that the total transmission power
243 over each subchannel is bounded by P_{max} . So we also have
244 a total transmit power constraint, given as

$$\sum_{i=1}^N P_k^i \leq P_{max}. \quad (5)$$

245 3.4. Problem description

246 In this work, we consider the maximal energy effi-
247 ciency, subject to the interference power constraint and
248 the total transmit power constraint. An optimization prob-
249 lem can be formulated as follows:

$$\begin{aligned} \text{maximize } E &= \frac{\sum_{i=1}^N U_k^i}{\sum_{i=1}^N P_k^i + P_c} \\ \text{subject to } \sum_{i=1}^N G_k^{0i} P_k^i &\leq I_{th} \\ \sum_{i=1}^N P_k^i &\leq P_{max}, \\ P_k^i &\geq 0. \end{aligned} \quad (6)$$

250 For the objective function of problem (6), the following
251 theorem is in order.

252 **Theorem 1.** The objective function $E = \frac{\sum_{i=1}^N U_k^i}{\sum_{i=1}^N P_k^i + P_c}$, where U_k^i
253 is given in (2), is a nonconvex function of P_k^i .

254 **Proof.** See Appendix. \square

255 **Theorem 1** indicates that problem (6) is nonconvex.
256 Since it is hard to transform the problem to a convex prob-
257 lem, we resort to PSO methods to solve it.

258 4. The proposed optimization algorithm

259 4.1. Standard particle swarm optimization

260 Similar to genetic algorithm (GA) [31,32] and differen-
261 tial evolution (DE) [33], PSO is a population-based opti-
262 mization method which was first proposed by Kennedy
263 and Eberhart [29]. The system is initialized with a popu-
264 lation of random solutions and it searches for optima by
265 updating generations. There are two learning processes in
266 the generation of PSO: cognitive learning process based
267 on individuals history, and social learning process based
268 on a swarm's history accumulated by sharing information
269 among all the particles in the swarm. Particles fly around
270 in multidimensional search space.

271 During flight, each particle adjusts its position accord-
272 ing to its own experience and the experience of neighbor-
273 ing particles, making use of the best position encountered
274 by itself and its neighbors. The direction of movement of
275 a particle is defined by the set of particles in the target

276 particle's vicinity and the target particle's history experi-
277 ence. Each particle keeps track of its coordinates in the
278 problem space, which are associated with the best solu-
279 tion achieved so far, denoted as $pbest$. Another best value
280 tracked by the global version of the optimizer is the overall
281 best value, and its location, obtained so far by any particle
282 in the population, denoted as $gbest$. At each time step, the
283 particle swarm optimization consists of velocity changes of
284 each particle toward $pbest$ and $gbest$ locations. Acceleration
285 is weighted by a random term, which separates random
286 numbers being generated for acceleration toward $pbest$ and
287 $gbest$ locations.

288 The a th particle's coordinates (position) is denoted as
289 $X_a = (x_{a1}, \dots, x_{ad})$, $d = 1, 2, \dots, D$, where D is the dimen-
290 sion of the optimal solution, and $V_a = (v_{a1}, \dots, v_{ad})$ denotes
291 the corresponding flight speed (velocity). Let $pbest_a =$
292 $(x_{a1}^{pbest}, \dots, x_{ad}^{pbest})$ and $gbest = (x_1^{gbest}, \dots, x_d^{gbest})$ be the best
293 position of individual a and its neighbors' best position so
294 far, respectively. Using the information, the updated veloc-
295 ity of individual a is modified under the following equation
296 in PSO:

$$V_a^{t+1} = \omega V_a^t + c_1 (pbest_a^t - X_a^t) + c_2 (gbest^t - X_a^t) \quad (7)$$

$$X_a^{t+1} = X_a^t + V_a^{t+1} \quad (8)$$

297 where t is the iteration number, c_1 and c_2 are constants,
298 which represent the weighting of the stochastic accelera-
299 tion terms that pull each particle toward $pbest$ and $gbest$
300 positions, and ω is the inertia weight parameter.

301 PSO is very efficient in solving complex optimization
302 problems. But it is easy to fall into local optimal solutions.
303 Then the inertia weight parameter is adjusted and chaos
304 theory is combined to PSO.
305

306 4.2. Adaptive inertia weight factor(AIWF)

307 It is clear that Eq. (7) represents the influence of pre-
308 vious velocity, which provides the necessary momentum
309 for particles to roam across the search space. The inertia
310 weight ω is the modulus that controls the impact of pre-
311 vious velocity on the current one. So, the balance between
312 exploration and exploitation in PSO is dictated by the value
313 of ω . Shi and Eberhart in [34,35] made a significant im-
314 provement in the performance of the PSO with a linearly
315 varying inertia weight over the generations. ω varies adap-
316 tively in response to the objective values of the particles.
317 In particular, AIWF is determined as follows.

$$\omega = \begin{cases} \omega_{min} + \frac{(\omega_{max} - \omega_{min})(f - f_{min})}{f_{avg} - f_{min}} & f \leq f_{avg}, \\ \omega_{max} & \text{otherwise,} \end{cases} \quad (9)$$

318 where ω_{max} and ω_{min} denote the maximum and mini-
319 mum of ω , respectively, f is the current objective value
320 of the particle, and f_{avg} and f_{min} are the average and
321 minimum objective values of all particles, respectively.
322 Under the guidance of Eq. (9), ω varies depending on the
323 objective value of the particle so that particles with low
324 objective values can be protected. AIWF provides a good
325 way to maintain population diversity and to sustain fast
326 convergence.

327 4.3. Chaos

328 Chaos theory demonstrates sensitive dependence on
329 initial conditions and also includes infinite unstable peri-
330 odic motions. Due to non-repetitive nature of chaos, it can
331 carry out overall searches at higher speeds than stochas-
332 tic esodic searches. The combination of optimization meth-
333 ods and chaotic systems is an important issue in non-linear
334 science. Here, logistic equation is employed to obtain chaos
335 queues, denoted as z_1, z_2, \dots , as follows:

$$z_{n+1} = \mu z_n (1 - z_n), n = 0, 1, 2, \dots \quad (10)$$

336 in which $0 \leq z_0 \leq 1$, and μ is the control parameter. When
337 $\mu = 4$, the system of (9) has been proved to be entirely
338 chaotic.

339 The essential procedure of chaotic particle swarm opti-
340 mization (CPSO) is as follows:

- 341 Step 1: Chaos initialization for particle.
342 Step 2: Evaluate the fitness function of each particle.
343 Step 3: Update each particle's velocity and position.
344 Step 4: Optimize the global best value by chaos search.
345 Step 5: If the stopping criteria satisfied, then output the
346 optimum solution, otherwise, loop to Step 2.

347 4.4. The large scale optimal algorithm using coevolution
348 chaotic particle swarm optimization (CCPSO)

349 Since the large scale and nonconvex nature makes our
350 formulated problem very complex, cooperative coevolu-
351 tion is applied, which has been proposed as a promising
352 framework for tackling large scale optimization problems
353 [36,37]. It can be regarded as an automatic approach to im-
354 plement the divide-and-conquer strategy. The detailed pro-
355 cedure of the proposed coevolution chaotic particle swarm
356 optimization (CCPSO) is as follows.

357 Particles' positions within the population in the CCPSO
358 represent the candidate solutions for solving the control
359 problem. That's to say, the procedure of particle searching
360 for the best position is to search for the power of prob-
361 lem (6). The CCPSO runs at the secondary base station
362 in the SU network, which is a centralized one. The sec-
363 ondary base station collects the necessary powers' infor-
364 mation from the secondary users and updates the memory
365 to perform the calculation, and then broadcasts the solu-
366 tion to the secondary users.

367 The objective is to maximize aggregate source net util-
368 ity per power consumption subject to the constraints in
369 CR networks. Constraints are handled based on the penalty
370 functions in the search space [38].

$$L(P_k^i, \mu_i, \kappa_i, v_i) = E + \mu_i(P_{max} - P_k^i) + \kappa_i \left(P_k^{th} - \sum_{i=1}^N G_k^{0i} P_k^i \right) + v_i(\gamma_k^{itar} - \gamma_k^i) + G(P_k^i) \quad (11)$$

371 where μ_i, κ_i and γ_i are the Lagrange multipliers for the
372 constraints. Since the optimization problem is nonconvex,
373 $G(P_k^i)$ is the penalty term added to the Lagrangian, which
374 can assure the max-min problem and the original problem
375 to be equal [38].

376 The optimization problem in CR is now in the form of
377 augmented lagrangian. Therefore, the problem is solved by

using max-min approach. Two populations of PSO are in- 378
cluded in the CCPSO. In the first PSO, the variable is P_k^i 379
and μ_i, κ_i and γ_i are set to be constant. The fitness de- 380
fines how well the position vector of each particle satisfies 381
the requirements of the optimization problem. The fitness 382
function for P_k^i is represented as 383

$$F_1(P_k^i) = \max(L(P_k^i, \mu_i, \kappa_i, \gamma_i)). \quad (12)$$

And in the second PSO, P_k^i is set to be constant, while 384
 μ_i, κ_i and γ_i are all set to be variables. The fitness func- 385
tion for μ_i, κ_i and γ_i is given as 386

$$F_2(\mu_i, \kappa_i, \gamma_i) = \min(L(P_k^i, \mu_i, \kappa_i, \gamma_i)). \quad (13)$$

The cooperation among particles is established through 387
the "history" $pbest_a$ and $gbest$, which are updated if better 388
fitness is obtained. The general procedure of CCPSO algo- 389
rithm can now be described by the following pseudocode:

Procedure of CCPSO

- 1: Initialization of two PSOs
- 2: Run the first PSO for generation 1
- 3: Re-evaluate the $pbest_a$ values for the second PSO if it is not in the first cycle
- 4: Run the second PSO for generation 2
- 5: Re-evaluate the $pbest_a$ values for the first PSO
- 6: If the termination condition is not met, go to Step 2;
- 7: Otherwise, output $gbest$.
- 8: End procedure

And the detailed improved PSO procedure is as follows. 390

391 Step 1: Chaos initialization is adopted to locate the 392
positions of particles and to increase the diversity of the 393
population and the ergodicity in the course of searching 394
without changing the randomness of algorithm when 395
initializing the particles. Some initial population with good 396
performances is chosen from the initial group with a large 397
number of population. Initialize a vector $z_1, z_2, \dots z_a$ each 398
component of which is set as a random value in the range 399
[0, 1], and generate chaos queues $z_1, z_2, \dots z_a$ by the itera- 400
tion of logistic Eq. (10). Transfer the chaos queues into the 401
range of the parameters according to following equation: 402

$$x_a^t = s_a + (t_a - s_a)z_a^t \quad (14)$$

where $[s_a, t_a]$ is the value range of each particle. 403

404 Step 2: Compute the fitness values of the vectors and
405 choose the best M solutions as the initial solutions of M
406 particles. Randomly initialize the velocity of M particles.

407 Step 3: Using the global best and the individual best
408 of each particle, each particle's velocity and position are
409 updated according to Eq. (7).

410 Step 4: Evaluate the fitness of each particle and com-
411 pare the evaluated fitness value of each particle to its in-
412 dividual best $pbest_a$. If $pbest_a$ is better than current value,
413 update $pbest_a$ as current position.

414 Step 5: If current value of the fitness function is bet-
415 ter than the global best $gbest$, update $gbest$ as the current
416 position.

417 Step 6: Optimize $gbest$ by chaos search. Firstly, scale
418 $gbest$ into [0,1] according to

$$x_a = (Gbest - s_a)/(t_a - s_a) \quad (15)$$

Table 1
Parameters of PSO.

Algorithm	Parameter settings	References
SA	$T_0 = 500^\circ\text{C}$, $C = 0.95$, $\Gamma = 9$	[41]
GA	$CRA = 0.55$, $MR = 0.005$, $CO = 0.35$, $MO = 0.25$	[32]
DE	$F = 0.5$, $CR = 0.9$	[33]
PSO	$\omega = 0.57$, $c_1 = c_2 = 2$, $v_{max} = 4$	[29]
CPSO	$\omega = 0.9$, $c_1 = c_2 = 2$, $v_{max} = 4$	[39]
CCPSO	$\omega: [0.4, 0.9]$, $c_1 = c_2 = 2$, $v_{max} = 4$	–

419 and generate chaos queues x_a^t by iteration of logistic equation
420 then transfer the chaos queues into the optimization
421 variable $gbest^t$ according to the following equation

$$gbest^t = s_a + (t_a - s_a)x_a^t \quad (16)$$

422 based on which the solution set is obtained: $gbest^t$. Com-
423 pute the fitness value of each feasible solution $gbest^t$ in the
424 problem space during chaotic search, and get the best so-
425 lution P_k^{i*} .

426 Step 7: When the constraints are violated, we pay an
427 extra charge proportional to the amount of violation with
428 the penalty value. And when the maximum iteration is
429 reached, then stop, we can get the global optimum P_k^{i*} that
430 are the solutions of the power allocation problem. Other-
431 wise, loop to Step 3.

432 5. Numerical results

433 In this section, we present numerical results for the
434 proposed power allocation algorithm in a network as
435 shown in Fig. 1. The purpose of the simulations is to show
436 that CCPSO can solve the large scale nonconvex optimiza-
437 tion problem. The total number of secondary users N is
438 set to 200. Without loss of generality, 100 tested users
439 are randomly selected for the simulations by 20 times. Ac-
440 cordingly, the numerical results are averaging results, and
441 the number of the decision variables for the optimization
442 problem is 100. Packet size is 1000 bytes and nodes are
443 equipped with a single transmitter/receiver, which has a
444 radio range of 500 m. An area of size $1000\text{ m} \times 1000\text{ m}$ is
445 considered. The maximum total transmit power is set as
446 $P_{max} = 800$ mW. The background noise is set as $\sigma^2 = 5 \times$
447 10^{-13} dB. No forward error correction is considered. The
448 link gains follow the path loss model: $G_{ii} = K/d_{ii}^\alpha$, and $G_{oi} =$
449 K/d_{oi}^α , where d_{ii} and d_{oi} are the distance from SU i to the
450 secondary BS and primary BS, respectively, α is a path loss
451 exponent and set as 4, and $K = 0.097$. The performance of
452 the proposed algorithm is compared with simulated an-
453 nealing (SA) [40,41], genetic algorithms (GA) [32], differ-
454 ential evolution (DE) [33], the standard PSO [29] and CPSO
455 [39]. All empirical experiments are conducted with a popu-
456 lation of 50, except for SA. The parameters of all the meth-
457 ods are all selected optimally in the simulations in Table 1
458 as selected in corresponding reference For SA method, T_0 ,
459 C and Γ are the initial temperature, cooling rate and stop-
460 ping parameter, respectively. For GA method, CRA , MR , CO
461 and MO are the crossover rate, mutation rate, crossover op-
462 erator, and mutation operator, respectively. For DE method,
463 F is the weighting factor and CR is the crossover constant.
464 Simulations were implemented on a PC with Intel Core™

i7-5820 CPU and memory capacity of 16G(8G*2) running
Matlab 7.12.0.

5.1. Performance comparisons

5.1.1. Convergence comparisons

465 In this example, the performance of the proposed algo-
466 rithm is compared with those of CPSO and PSO. The utility
467 function is considered as $U_k^i = \frac{1}{1+e^{-a_i(y_k^i - b_i)}}$, which is more
468 reasonable for real networks. In this case, the maximum
469 number of iterations of the algorithms is set to be 1000. N
470 is set to be 100, so the energy consumption is fixed with
471 the fixed number of N . Fig. 2 shows the average conver-
472 gence of the best individuals of each iteration for the sys-
473 tem with different methods, where the system parameters
474 a_i is varied and b_i is fixed. In this case from the figure
475 we can find that the proposed algorithm based on CCPSO
476 converges less than 100 iterations. Compared to CPSO pro-
477 posed in [39], which just uses the logistic mapping, the
478 proposed algorithm in this paper can get better perfor-
479 mance. It is because of the coevolutionary theory, which
480 is more suitable for the large-scale-global characteristic of
481 the problem. And the system is more stable when the pa-
482 rameter a_i is large with fixed b_i . Fig. 3 shows the conver-
483 gence of the system with fixed a_i and varying b_i , and in
484 this case a large b_i corresponds to a high total utility. As
485 seen from Figs. 2 and 3, we can see that the parameters a_i
486 and b_i can be used to tune the steepness and the center of
487 the utility, respectively. And the proposed approach based
488 on CCPSO shows better performance in the process of the
489 algorithm and has faster convergence speed than other
490 approaches.

491 In order to verify the trade-off of utility vs energy. In
492 this case, the energy efficiency with the increase of the
493 value of N is given in Fig. 4. The utility function is con-
494 sidered as $U_k^i = \frac{1}{1+e^{-a_i(y_k^i - b_i)}}$ with $a = 5$, $b = 20$. It is ob-
495 served that energy efficiency initially increases due to the
496 node number increases. This is because that the utility of
497 the system increases relatively higher than the consump-
498 tion of power increases. It indicates a trade-off between
499 system utility and total energy consumption. From the fig-
500 ure, we can find that N is found to be 80 that identifies
501 the minimum energy consumption and the maximum util-
502 ity function.

5.1.2. Performance comparisons with different methods

503 Tables 2–5 list the best, worst, mean value of the op-
504 timal energy efficiency solved by GA, SA, DE, PSO, CPSO
505 and CCPSO in the 1000 runs with different a and b with
506

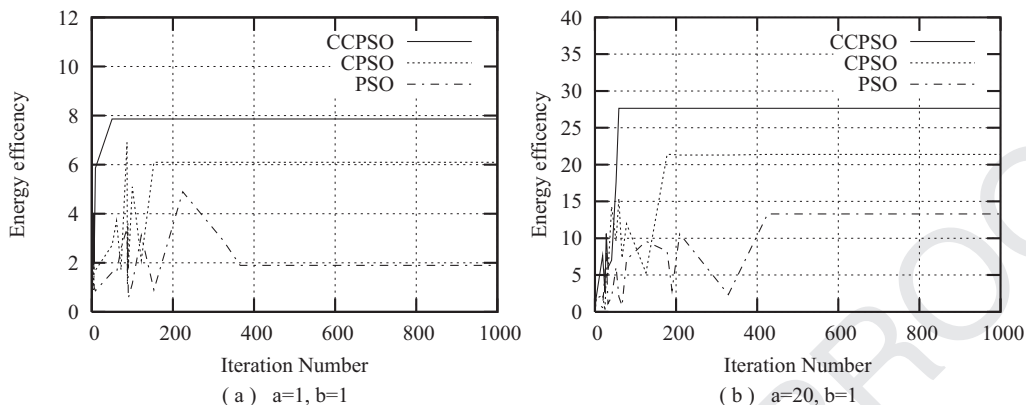


Fig. 2. Energy efficiency with varying a and fixed b when (a) $a = 1, b = 1$ (b) $a = 20, b = 1$.

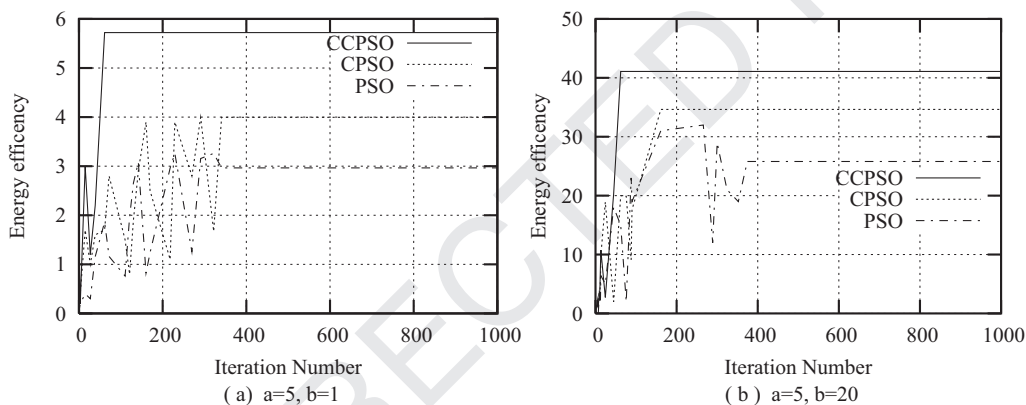


Fig. 3. Energy efficiency with varying b and fixed a when (a) $a = 5, b = 1$ (b) $a = 5, b = 20$.

Table 2
Energy efficiency when $a = 1, b = 1$ with fixed b .

Algorithms	Best	Worst	Average	SD	Average CPU time (s)
SA	1.9699	1.1084	1.502	0.1694	10.04
GA	2.1014	1.209	1.6568	0.1911	10.69
DE	5.6148	4.8811	5.0426	0.0694	11.46
PSO	2.3611	1.3408	1.8961	0.1999	7.24
CPSO	7.045	7.4876	7.1092	0.0102	10.36
CCPSO	7.8652	7.8652	7.8652	0	12.47

Table 3
Energy efficiency when $a = 20, b = 1$ with fixed b .

Algorithms	Best	Worst	Average	SD	Average CPU time (s)
SA	12.3072	9.7767	10.8538	0.4630	22.32
GA	16.4327	10.6643	11.9427	0.5495	23.65
DE	22.3606	19.3033	21.0428	0.5665	26.74
PSO	17.5826	11.9132	13.3112	0.6009	16.53
CPSO	24.5542	22.323	23.394	0.4604	22.46
CCPSO	27.6501	27.3951	27.4776	0.0355	29.38

511 $N = 100$. In order to verify the statistical performance of
 512 the proposed algorithm, we also give the standard deviation
 513 (SD) in the tables. It can be seen that the CCPSO algorithm
 514 can provide better “Best”, “Worst”, “Mean”, “SD” and
 515 “Average CPU time” results for the test functions. SD means
 516 the volatility of the data, from the tables we can see that

some SD of the proposed algorithm are 0, which shows
 that the proposed algorithm is more stable. Because the
 coevolution theory solve the proposed large and complex
 problem using a divide-and-conquer strategy, which avoids
 producing infeasible solutions, so these statistical test
 results reveal that the proposed method outperforms other

Table 4
Energy efficiency when $a = 5$, $b = 1$ with fixed a .

Algorithms	Best	Worst	Average	SD	Average CPU time (s)
SA	2.5152	1.9473	2.3087	0.0888	6.34
GA	3.1082	2.3675	2.6652	0.1280	7.02
DE	4.1462	4.0615	4.1028	0.0178	8.64
PSO	3.6724	2.5032	2.9659	0.1989	5.19
CPSO	4.654	4.112	4.2978	0.2821	6.89
CCPSO	5.7209	5.7143	5.7189	0.0009	9.65

Table 5
Energy efficiency when $a = 5$, $b = 20$ with fixed a .

Algorithms	Best	Worst	Average	SD	Average CPU time (s)
SA	18.6461	10.9909	13.2166	0.9567	23.98
GA	23.6888	14.0908	18.6109	1.9430	24.66
DE	32.0232	28.7112	31.0655	0.4117	33.04
PSO	29.1253	17.2218	22.2320	1.0890	22.07
CPSO	39.099	37.20878	38.627	0.4694	24.47
CCPSO	41.0927	41.0927	41.0927	0	35.68

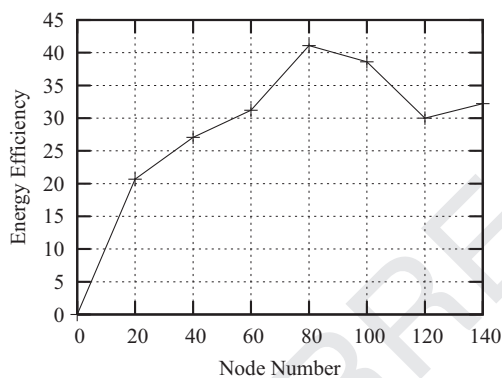


Fig. 4. Energy efficiency with the increase of node number.

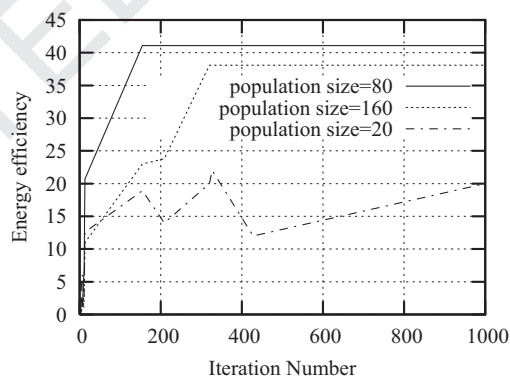


Fig. 5. Effect of population size.

523 existing methods and has stronger robustness than other
 524 methods. Since the proposed algorithm is a dual-swarm
 525 one, the convergence time of the proposed algorithm is
 526 more than the other algorithms, but the results are compar-
 527 able. The chaotic initialization and search have strong
 528 ability to jump out of local optima, which can help the al-
 529 gorithm reduce the search time.

530 5.2. Influence of population size

531 The population size is an important factor which influ-
 532 ences the performance of the stochastic search algorithm.
 533 Too small population may not be able to reach the maxi-
 534 mum value and achieve an optimum, while too large
 535 population makes the proposed algorithm slow and com-
 536 putationally inefficient. Tests are carried out for population
 537 of 20, 80, and 160. We consider the utility function in the
 538 case of $a = 5$, $b = 20$. As shown from Fig. 5, we can see that
 539 the convergence to the optimum is hardly achieved for the
 540 proposed optimization algorithm when the population is
 541 set as 20, and the convergence speed is slower when the
 542 population is set as 160. The size 80 is found to be optimal.

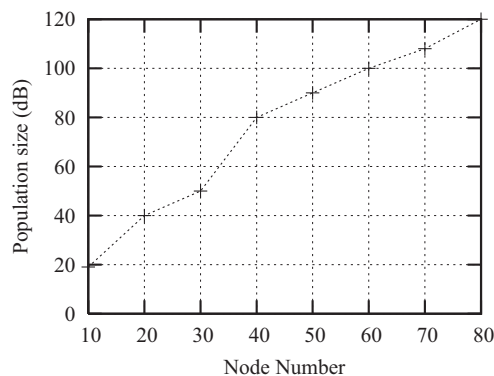


Fig. 6. The relationship between the population size and the number of nodes.

And the relationship between the number of nodes 543
 and particles is studied in this case. As shown from Fig. 6, 544
 when the number of nodes varies from 10 to 100, the 545
 number of particles needs to increase from 40 to 120. 546
 This is because when the network is in a large scale, 547

548 the particles need to be vectors with higher dimensions,
549 which leads to a larger search space.

550 **6. Conclusion**

551 In this paper, a new metric that reflects the trade-off
552 between the utility and energy consumption is defined in
553 CR networks. Since the optimization problem is a large
554 scale and nonconvex one, our proposed algorithm exploits
555 the coevolutionary and chaotic ideas for the dynamic
556 power allocation problem in CR networks. Strict assump-
557 tions such as continuity, differentiability, and convexity
558 of the objective function are not necessary. The formu-

$$\frac{\partial E}{\partial P_k^i} = \frac{\frac{\partial U_k^i}{\partial P_k^i} (P_k^i + \frac{P_c}{N}) - U_k^i}{(P_k^i + \frac{P_c}{N})^2} \tag{18}$$

$$= \frac{\frac{\partial U_k^i}{\partial Q} \frac{\partial Q}{\partial \gamma_k^i} \frac{\partial \gamma_k^i}{\partial P_k^i} (P_k^i + \frac{P_c}{N}) - U_k^i}{(P_k^i + \frac{P_c}{N})^2} \tag{19}$$

$$= \frac{1}{(P_k^i + \frac{P_c}{N})^2} \left[\frac{a_i Q}{(1+Q)^2} \frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} (P_k^i + \frac{P_c}{N}) - \frac{1}{1+Q} \right] \tag{20}$$

$$\frac{\partial^2 E}{(\partial P_k^i)^2} = \frac{\partial}{\partial P_k^i} \left(\frac{\partial E}{\partial P_k^i} \right) \tag{21}$$

$$= \frac{\left(\frac{\partial U_k^i}{\partial P_k^i} (P_k^i + \frac{P_c}{N}) - U_k^i \right) (P_k^i + \frac{P_c}{N})^2 - 2 \left(\frac{\partial U_k^i}{\partial P_k^i} (P_k^i + \frac{P_c}{N}) - U_k^i \right) (P_k^i + \frac{P_c}{N})}{(P_k^i + \frac{P_c}{N})^4} \tag{22}$$

$$= \frac{\left(\frac{\partial U_k^i}{\partial P_k^i} \right) (P_k^i + \frac{P_c}{N})^3 - 2 (P_k^i + \frac{P_c}{N}) \left[\frac{a_i Q}{(1+Q)^2} \frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} (P_k^i + \frac{P_c}{N}) - \frac{1}{1+Q} \right]}{(P_k^i + \frac{P_c}{N})^4} \tag{23}$$

$$= \frac{\left(\frac{a_i Q}{(1+Q)^2} \frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} \right) (P_k^i + \frac{P_c}{N})^3 - 2 (P_k^i + \frac{P_c}{N}) \left[\frac{a_i Q}{(1+Q)^2} \frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} (P_k^i + \frac{P_c}{N}) - \frac{1}{1+Q} \right]}{(P_k^i + \frac{P_c}{N})^4} \tag{24}$$

$$= \frac{a_i^2 Q(Q-1)}{(1+Q)^3} \left(\frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} \right)^2 (P_k^i + \frac{P_c}{N})^3 - 2 (P_k^i + \frac{P_c}{N}) \left[\frac{a_i Q}{(1+Q)^2} \frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} (P_k^i + \frac{P_c}{N}) - \frac{1}{1+Q} \right]}{(P_k^i + \frac{P_c}{N})^4} \tag{25}$$

$$= \frac{a_i^2 Q(Q-1)}{(1+Q)^3} \left(\frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} \right)^2 (P_k^i + \frac{P_c}{N})^2 - \frac{2a_i Q}{(1+Q)^2} \frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} (P_k^i + \frac{P_c}{N}) + \frac{2}{1+Q}}{(P_k^i + \frac{P_c}{N})^3} \tag{26}$$

559 lated optimization problem is solved by using max–min
560 approach, where two populations of PSO are included. The
561 performance of the proposed algorithm is compared with
562 those of related methods in the literature. It is observed
563 that the proposed algorithm is indeed capable of quickly
564 achieving energy-efficient solutions. Future research topics
565 may include dynamic PSO algorithms for cooperative CR
566 networks in which one SU may help relay other SUs' signal
567 to the secondary BS such that cooperative diversity can be
568 achieved. Then each SU may need to distribute its power
569 budget in transmitting its own signal and in relaying other
570 SUs' signals. The problem is much more complex, and
571 deserves further investigation.

572 **Appendix. Proof of Theorem 1**

573 **Proof.** We use proof by contradiction. Suppose that $E =$
574 $\frac{\sum_{i=1}^N U_k^i}{\sum_{i=1}^N P_k^i + P_c}$ is a convex function of P_k^i , then

$$\frac{\partial^2 E}{(\partial P_k^i)^2} \leq 0 \tag{17}$$

575 and Eq. (17) is a necessary condition. For convenience, we
576 set $e^{-a_i(\gamma_k^i - b_i)}$ as Q , thus:

$$\frac{a_i^2 Q(Q-1)}{(1+Q)^3} \left(\frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} \right)^2 (P_k^i + \frac{P_c}{N})^2 - \frac{2a_i Q}{(1+Q)^2} \frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} (P_k^i + \frac{P_c}{N}) + \frac{2}{1+Q} = 0 \tag{27}$$

We set

$$\frac{a_i^2 Q(Q-1)}{(1+Q)^3} \left(\frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} \right)^2 (P_k^i + \frac{P_c}{N})^2 = a \tag{28}$$

$$\frac{2a_i Q}{(1+Q)^2} \frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} (P_k^i + \frac{P_c}{N}) = b \tag{29}$$

$$\frac{2}{1+Q} = c \tag{30}$$

Then

$$b^2 - 4ac = \left(\frac{2a_i Q}{(1+Q)^2} \frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} \right)^2 - \frac{8}{1+Q} \frac{a_i^2 Q(Q-1)}{(1+Q)^3} \left(\frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} \right)^2 \tag{31}$$

$$= \left[\left(\frac{2a_i Q}{(1+Q)^2} \right)^2 - \frac{8}{1+Q} \frac{a_i^2 Q(Q-1)}{(1+Q)^3} \right] \left(\frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} \right)^2 \quad (32)$$

586

$$= \frac{4a_i^2 Q(2-Q)}{(1+Q)^4} \left(\frac{G_k^{ii}}{\sum_{j \neq i} G_k^{ij} P_k^j + \sigma^2} \right)^2 \quad (33)$$

587 When $Q > 2$, $b^2 - 4ac < 0$; when $Q < 2$, $b^2 - 4ac > 0$,
 588 which cannot assurance that $\frac{\partial^2 E}{(\partial P_k^i)^2} \leq 0$. So E is a noncon-
 589 vex function. \square

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