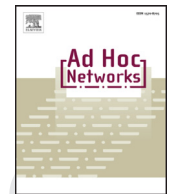




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Dynamic deployment of randomly deployed mobile sensor nodes in the presence of obstacles

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ABSTRACT

For random deployment of wireless sensor networks in a specified geographical location and in the presence of obstacles, optimal network coverage is highly desirable while maintaining network connectivity. In this piece of work, we propose an efficient autonomous deployment scheme, named as Obstacle Avoidance Virtual Force Algorithm (OAVFA), for self-deployment of randomly scattered homogeneous as well as heterogeneous mobile sensor nodes over a squared sensing field to enhance the network coverage and ensure the network connectivity in the presence of obstacles. Our proposed approach is localized in the sense that each decision taken by the sensor node is strictly based on information acquired from its neighbors. The simulation results show that OAVFA provides an efficient self-deployment of mobile sensor nodes in the presence of obstacles.

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1. Introduction

Nowadays, Wireless Sensor Networks (WSNs) have attracted tremendous research interest due to its various applications from environment monitoring, battlefield surveillance, target tracking, wildfire detection, precision agriculture, smart homes and offices, industrial process monitoring and asset management [1]. A mobile sensor network is a collection of inexpensive, low-powered, small size, and multifunctional mobile sensor nodes. The effectiveness of WSNs mainly depends on the network coverage, lifetime and connectivity provided by the sensor deployment strategies such as deterministic and random deployment. Placing sensor nodes manually in predetermined positions on the basis of simple geometric structure (e.g., Hexagon, Square, Rhombus, and Triangular Lattice) is simple and optimal, but this deployment strategy is not suitable in many applications where the application

environment is unknown, hostile or inhospitable. For these applications, sensor nodes are required to be deployed randomly by means of dispersing sensors from aircraft or artillery ordinance.

An efficient self-deployment algorithm is highly required to ensure optimal network coverage while maintaining connectivity for such randomly deployed sensors. Presently, virtual force-based self-deployment strategies are adopted to overcome the limitations exhibited by random deployment [2–10]. In this work, an efficient distributed self-deployment algorithm has been proposed for randomly deployed homogeneous as well as heterogeneous mobile sensor nodes. This algorithm is named as Obstacle Avoidance Virtual Force Algorithm (OAVFA). Experimental results carried out with our proposed algorithm not only maximizes coverage area but also ensures the connectivity between all sensor nodes in the presence of obstacles. A set of sensor nodes with identical speeds, communication ranges, and sensing ranges has been identified as homogeneous sensor nodes while heterogeneous sensor nodes differs only in the sensing ranges which are strictly different for various sensors. It has been assumed that the speeds and the communication ranges

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41 for heterogeneous sensor remain constant during the
42 process.

43 The proposed algorithm is localized and executed at
44 each sensor node. In this algorithm, each sensor node con-
45 sideres all attractive and repulsive virtual forces due to its
46 neighboring sensor nodes, obstacles, and the sensing field
47 boundary to determine its movements to enhance the net-
48 work coverage while maintaining connectivity, prevent the
49 sensor nodes from moving out of sensing field boundary,
50 and avoid the obstacles. Here neighbor sensor nodes of i th
51 sensor s_i means the sensor nodes that are within the com-
52 munication range of s_i .

53 In the next section, a brief but latest literature sur-
54 veys on sensor node deployment has been outlined.
55 Section 3 provides a basic discussion about the network
56 coverage and sensing model. Our proposed deployment
57 algorithm, Obstacle Avoidance Virtual Force Algorithm
58 (OAVFA) has been described in Section 4. In Section 5,
59 simulation results are presented followed by conclusions
60 in Section 6.

61 2. Related work

62 Sensor arrangement is an imperative issue for some
63 essential objectives in WSNs like coverage, lifetime, and
64 connectivity. For randomly deployed sensor networks, an
65 efficient deployment algorithm is required to self-deploy
66 the mobile sensor nodes to maximize coverage area,
67 ensure the network connectivity and prolong the network
68 lifetime. In [2,3], an incremental and greedy algorithm
69 is presented in which nodes are deployed one at a time.
70 The objective is to maximize the coverage under the
71 constraint that nodes maintain line of sight with each
72 other. Howard et al. [4] have presented a centralized
73 deployment approach based on potential field theory to
74 deploy the mobile sensor nodes (mobile robots) in an
75 unknown environment to enhance the network coverage.
76 In [5,6], the sensor nodes are placed in a grid-like manner
77 to ensure coverage and connectivity. A robust and scalable
78 deployment scheme, based on simulated annealing tech-
79 nique for complete coverage is presented in [7]. In [8], Heo
80 and Varshney have proposed a distributed self-deployment
81 algorithm for mobile sensor networks to maximize the
82 coverage and to maintain uniformity in node distribution.
83 Poduri and Sukhatme [9] have proposed a deployment
84 scheme for mobile sensor network to enhance the net-
85 work coverage with maintaining K -connectivity. In [10],
86 Guo et al. have proposed an adaptive coverage algorithm
87 by considering inner repulsion, random disturbance and
88 boundary contraction to maximize the coverage. By com-
89 bining the potential field theory and the plate coverage
90 theory, a centralized deployment algorithm called as a Vir-
91 tual Force Algorithm (VFA) is presented in [11,12]. This VFA
92 cannot quickly converge to a steady state. In [13], the au-
93 thors proposed a sensor deployment optimization strategy
94 based on Target Involved Virtual Force Algorithm (TIVFA)
95 to improve coverage and detection probability. In [14],
96 Wang et al. have proposed several algorithms that identify
97 existing coverage holes in the network and compute the
98 desired target locations where sensor should move in or-
99 der to increase the coverage. In [15], the authors developed

100 a decentralized and scalable algorithm based on potential
101 field theory for motion control of mobile sensor networks
102 to cover the maximum area of the free space in minimum
103 time. A localized algorithm for determining whether every
104 point in the service area of the sensor network is covered
105 by at least k sensors is presented in [16]. Voronoidiagram
106 and Delaunay triangulation are used in [17] to estimate
107 the worst and best case coverage in a sensor network.
108 In [18], the authors used Delaunay triangulation, Gabriel
109 graph and relative neighborhood graph to find the path
110 with best coverage. A few excellent surveys on the present
111 state-of-the-art research on sensor network is presented
112 in [19–23]. In [24], the authors have explored geographic
113 routing in duty-cycled mobile WSNs and proposed two
114 geographic-distance-based connected- k neighborhood
115 (GCKN) sleep scheduling algorithms for geographic rout-
116 ing schemes. In [25], the authors gave necessary and
117 sufficient conditions for 1-coverage and 1-connected
118 wireless sensor grid network. Tian and Georgansa
119 [26] have proved that the communication range is twice of
120 the sensing range is the sufficient condition for complete
121 coverage preservation implies connectivity among active
122 nodes if the original network is connected. The optimal
123 deployment patterns to achieve both full coverage and
124 connectivity for all ranges of R_c/R_s is presented in [27–29].
125 In [30], the authors proposed a self-deployment mecha-
126 nism that allows to maintain network connectivity during
127 the deployment of mobile sensor nodes. This mechanism
128 is robust against message losses during deployment. Shen
129 et al. [31] have proposed a grid scan method to calculate
130 coverage rate for arbitrary sensing radius. The main objec-
131 tive of this approach is to provide a better coverage with
132 less nodes. In [32], the authors developed a mechanism
133 to replace failed sensors in a large-scale static sensor
134 networks by using few mobile robots. The goal of this
135 work is to minimize the motion and the messaging over-
136 head. Chen et al. [33] have proposed two novel algorithms
137 named as Improved Virtual Force Algorithm (IVFA) and
138 Exponential Virtual Force Algorithm (EVFA) to improve
139 the performance of traditional VFA. In [34], the authors
140 presented an efficient deployment algorithm named as
141 Self-Deployment by Density Control (SDDC). In this work,
142 virtual force is decided by density at a sensor node and
143 obstacles and the algorithm is not suitable for sparse
144 initial distribution. In [35], Kribi et al. have proposed
145 Dth_Lmax_Serialized_VFA algorithm to enhance coverage
146 and maintain network connectivity of the sensor networks.
147 A Virtual Force directed Co-evolutionary Particle Swarm
148 Optimization (VFCPSO) is presented in [36]. This algorithm
149 is appropriate for small scale application due to its high
150 computation time. Yu et al. [37] have proposed an algo-
151 rithm base on virtual force and the concept of adjacent
152 relationship of nodes to enhance the coverage rate and
153 reduce the convergence time. A Distributed Virtual Forces
154 Algorithm (DVFA) is proposed in [38] to establish coverage
155 and connectivity. The problem of connectivity optimization
156 in random 3D networks is addressed in [39] where the de-
157 ployment problem considers the maximization of network
158 connectivity satisfying lifetime constraints. Autonomous
159 mobile robots that deploy a wireless sensor network to
160 be used in disasters is introduced in [40]. In [41], the

161 authors proposed a deployment algorithm for hetero- 209
 162 geneous sensor networks based on the circle packing 210
 163 technique to enhance the coverage area. In [42], Xiaoping 211
 164 et al. have analyzed the performance of different virtual 212
 165 force models used in node deployment algorithms. 213

166 In this work, an obstacle avoidance VFA is introduced
 167 for deployment of both heterogeneous as well as homo-
 168 geneous mobile sensor nodes over a squared sensing field
 169 containing different shape of obstacles.

170 3. Coverage and sensing model

171 Coverage is one of the key parameters to evaluate the
 172 performance of deployment algorithms [2–9]. According to
 173 Poduri and Sukhatme [9], there are three categories of cov-
 174 erage: barrier coverage, target or point coverage, and area
 175 coverage. In barrier coverage, sensor nodes have to form a
 176 barrier to detect intruders. Target coverage refers to moni-
 177 toring fixed number of targets in a Region of Interest (ROI).
 178 Area coverage means that every point within ROI must be
 179 monitored by at least one sensor node or by the joint de-
 180 tection of several sensor nodes. Usually, this coverage is
 181 necessary when applications need to monitor the entire
 182 area of interest. In general, area coverage [31] means how
 183 well the ROI is monitored by the sensor network and is
 184 evaluated as in (1).

$$185 \text{Coverage}(C) = \frac{\cup_{i=1, 2, \dots, N} A_{si}}{A_{Tot}} \quad (1)$$

186 Where A_{si} denotes the area covered by the mobile sensor
 187 node s_i , N is the number of mobile sensor nodes deployed
 188 in ROI and A_{Tot} is the area of the entire ROI.

189 Sensor models have direct impact on network coverage
 190 of WSNs [43]. Sensing models as reported in various lit-
 191 eratures can broadly be classified as Binary sensor model
 192 and Probabilistic sensing model [10–13,43]. For the pur-
 193 pose of evaluation of our proposed algorithm, we prefer
 binary sensor model.

194 3.1. Binary sensor model (BSM)

195 In most of the existing work, the disk sensing model
 196 is used for coverage calculation for its simplicity. Accord-
 197 ing to this model [11,12,31], an event is detected by a sen-
 198 sor node s_i with a detection probability 1, if the occurrence
 199 of the event is within the sensing radius R_s of the sensor
 200 node s_i . Otherwise the probability of detection is 0 as given
 201 in (2).

$$202 C_{xy}(p, s_i) = \begin{cases} 1, & \text{if } d(s_i, P) \leq R_s \\ 0, & \text{if } d(s_i, P) > R_s \end{cases} \quad (2)$$

203 Where $d(s_i, P) = \sqrt{(x_i - x)^2 + (y_i - y)^2}$ is Euclidean dis-
 204 tance between the i th sensor node $s_i(x_i, y_i)$ and the
 event occurring point $P(x, y)$.

205 3.2. Coverage ratio calculation

206 For randomly deployed sensor networks, the cover-
 207 age calculation by geometric analysis is too complicated.
 208 Therefore, we adopt a grid scan method [31] to evaluate

the coverage ratio. According to this method, the entire ROI
 is divided into a specified number of uniform grids and
 each grid is denoted by its center point. The grid is covered
 if its center point is within the sensing range of a sensor
 node and the coverage ratio is calculated as in (3).

$$214 \text{Coverage}(C) = \frac{m}{n} \quad (3)$$

215 Where, m represents the number of grids covered by the
 216 sensor nodes and n is the number of total grids in entire
 ROI. For binary model, $m = \text{card}(\cup_{i=1, 2, \dots, N} G_i)$, where G_i
 217 denotes the grid points within the sensing range R_s of the
 218 i thth sensor node. Here, by $\text{card}(\cdot)$ we indicate cardinal-
 219 ity of a set. The accuracy of this method depends upon the
 220 size of the grid, the smaller the grid size the more accurate
 221 the method.

222 4. Obstacle Avoidance Virtual Force Algorithm (OAVFA)

223 The proposed OAVFA is based on the following assump-
 224 tions. They are: (i) all the sensor nodes have locomotion
 225 capability and can move effectively to any direction and
 226 any distance within the sensing boundary, (ii) each sen-
 227 sor node has one unique ID, (iii) all sensors are equipped
 228 with localization system (i.e. GPS), (iv) every sensor node
 229 is able to acquire the relative position of the other sensor
 230 nodes within its communication range, (v) all the sensor
 231 nodes have circular sensing and communication areas, (vi)
 232 the sensing field is a square sized area demarcated with
 233 a clear boundary, (vii) the sensing field contains obstacles
 234 of different shapes and sizes, (viii) every sensor node is
 235 able to detect the shape and position of any obstacles in its
 236 sensing range and can calculate the nearest distance from
 237 the obstacle by using the time-of-flight method.

238 The main objective of our proposed OAVFA is not only
 239 to maximize the coverage area but also to reduce the
 240 moving energy requirement in the presence of different
 241 obstacles in ROI. Each sensor node s_i is subjected to an
 242 attractive or repulsive force (\vec{F}_{ij}) by its neighbor sensor
 243 node s_j , a repulsive force \vec{F}_{iO_m} by an obstacle O_m , and a
 244 repulsive force \vec{F}_{ib} by sensing field boundaries. Therefore,
 245 the net force on a sensor node s_i is evaluated as in (4).

$$246 \vec{F}_i = \sum_{j=1, j \neq i}^K \vec{F}_{ij} + \sum_{m=1}^{N_0} \vec{F}_{iO_m} + \vec{F}_{ib} \quad (4)$$

247 Where K is the number of neighborhood sensor
 248 nodes of s_i ; N_0 is the number of obstacles in ROI.
 249 Depending on the calculated total force \vec{F}_i , the sensor node
 s_i moves to its new location as given in (5).

$$250 x_{inew} = x_{iold} + F_{ix}; \quad y_{inew} = y_{iold} + F_{iy} \quad (5)$$

251 Where x_{iold} and y_{iold} denote the current location of sensor
 252 node s_i ; x_{inew} and y_{inew} denote the next location of sensor
 253 node s_i ; F_{ix} and F_{iy} denote the x and y directional com-
 254 ponents respectively of the displacement s_i goes through
 255 as the same is subjected by the force \vec{F}_i . The maximum
 256 distance traveled by a sensor node in each iteration is
 257 decided by its velocity. So we restrict the upper limit of F_{ix}
 and F_{iy} by introducing two thresholds Th_x and Th_y .

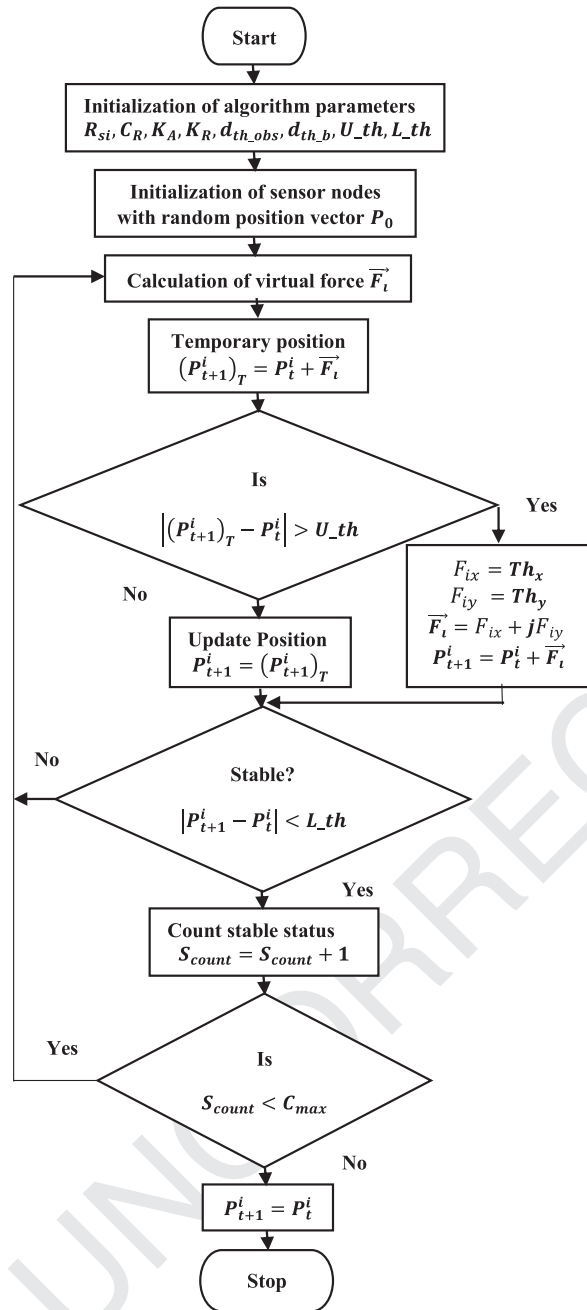


Fig. 1. Obstacle Avoidance Virtual Force Algorithm (OAVFA).

258 Fig. 1 shows the flow chart of our proposed algorithm.
 259 This localized deployment algorithm is executed at each
 260 sensor node s_i and the sensor node ceases its movement
 261 if it moves less than a predefined threshold (L_th) for the
 262 time duration C_{max} .

263 4.1. Virtual force due to sensor node

264 Consider a network of N sensor nodes $s_1, s_2, s_3, \dots, s_N$
 265 at positions $p_1, p_2, p_3, \dots, p_N$ with sensing radius $R_{s1}, R_{s2},$
 266 R_{s3}, \dots, R_{sN} respectively and each sensor node is defined

267 by its communication range C_R . Let d_{ij} represent the Euclidean
 268 distance between the sensor nodes s_i and s_j , i.e. $d_{ij} = \|p_i - p_j\|$
 269 and the force exerted on s_i by the neighbor sensor node s_j
 270 be denoted by \vec{F}_{ij} . The force model is given in (6).
 271

$$\vec{F}_{ij} = \begin{cases} 0 & \text{if } d_{ij} > C_R \\ (K_A(d_{ij} - d_{th}^{ij})) \left(\frac{p_i - p_j}{d_{ij}} \right) & \text{if } C_R \geq d_{ij} > d_{th}^{ij} \\ 0 & \text{if } d_{ij} = d_{th}^{ij} \\ (K_R(d_{th}^{ij} - d_{ij})) \left(\frac{p_i - p_j}{d_{ij}} \right) & \text{if } d_{ij} < d_{th}^{ij} \end{cases} \quad (6)$$

272 Where K_A and K_R are the force coefficients. Usually $K_A \leq$
 273 K_R . The threshold distance d_{th}^{ij} controls the overlapping degree
 274 between the sensor nodes s_i and s_j and for our proposed model
 275 $d_{th}^{ij} = \frac{\sqrt{3}}{2} (R_{s_i} + R_{s_j})$. In case of homogeneous sensor network,
 276 the sensing range is identical for all sensor nodes i.e. $R_{s1} = R_{s2} = \dots = R_{sN} = R_s$
 277 and the threshold distance $d_{th} = \sqrt{3}R_s$.
 278

279 4.2. Force model of obstacle on sensor

280 The obstacles such as walls or buildings exert repulsive forces
 281 on a sensor node. Let d_{iO_j} is the shortest distance between the
 282 sensor node s_i and the obstacle O_j and (x_{O_j}, y_{O_j}) is the nearest
 283 point in the obstacle O_j from sensor node s_i . If the distance d_{iO_j}
 284 is less than a pre-defined threshold distance d_{th_obs} , a repulsive
 285 force is exerted by the obstacle O_j on sensor node s_i and the
 286 force is computed as in (7).
 287

$$\vec{F}_{iO_j} = \begin{cases} 0 & \text{if } d_{iO_j} \geq (d_{th_obs}) \\ (K_{R1} (d_{th_obs} - d_{iO_j}), \alpha_{iO_j} + \pi) & \text{if } d_{iO_j} < (d_{th_obs}) \end{cases} \quad (7)$$

288 Where K_{R1} is a constant parameter that represents the strength
 289 of the repulsive force.

290 4.3. Boundary force on sensor

291 The boundary forces on the sensor reduce the unwanted coverage
 292 outside the ROI. The boundaries of sensing field exert repulsive
 293 forces on a sensor. Let d_{ib} is the perpendicular distance between
 294 the sensor node s_i and the sensing field boundary. If the distance
 295 d_{ib} is less than a pre-defined threshold distance d_{th_b} , a
 296 repulsive force is exerted by the boundary on sensor node s_i
 297 and the force is computed as in (8).
 298

$$\vec{F}_{ib} = \begin{cases} 0 & \text{if } d_{ib} \geq (d_{th_b}) \\ (K_{R2}(d_{th_b} - d_{ib}), \alpha_{ib} + \pi) & \text{if } d_{ib} < (d_{th_b}) \end{cases} \quad (8)$$

299 In a squared area, the boundary forces will be there due to the
 300 four boundaries surrounding the ROI. Thus \vec{F}_{ib} is the combined
 301 force from all boundaries as given in (9).

$$\vec{F}_{ib} = \vec{F}_{ib}^{x_1} + \vec{F}_{ib}^{x_2} + \vec{F}_{ib}^{y_1} + \vec{F}_{ib}^{y_2} \quad (9)$$

302 The above virtual forces guide the mobile sensor nodes to enhance
 303 the area coverage while maintaining connectivity, prevent the sensor
 304 nodes from moving out of sensing field boundary, and avoid the
 305 obstacles. In OAVFA, each

Table 1
Simulation parameters.

Parameters	Value
Field size	100 m × 100 m
Grid size	1 m × 1 m
Max. velocity of sensor node	0.5 m/s
K_A	0.001
K_R	0.2
K_{R1}	0.8
K_{R2}	0.8
U_{th}	0.5
L_{th}	0.001
C_{max}	10
Max_iteration	300

Table 2
Simulation parameters for homogeneous sensor.

Parameters	Value
Sensing Range (R_s)	10 m
Communication Range ($C_R = 2 \times R_s$)	20 m
Threshold distance ($d_{th} = \sqrt{3}R_s$)	17.32 m
$d_{th,obs} = \sqrt{3}R_s/2$	8.66 m
$d_{th,d} = \sqrt{3}R_s/2$	8.66 m

node stops its movement when it has reached its stable position.

In this paper, the performances of distributed deployment algorithms are evaluated by considering two aspects: coverage ratio and moving energy consumption. Coverage ratio is the ratio of the number of grid points that are not in obstacle and have a detection probability of 1 to the total number of grid points in ROI that are not in obstacles and is evaluated as in (3). Moving energy consumption means the energy required for movement of sensor nodes. In this work, the moving energy consumption is considered as the average moving distance of all sensor nodes in each step and is calculated as in (10).

$$D_{avg} = \frac{\sum_{i=1}^N \sqrt{(x(i)_{new} - x(i)_{old})^2 + (y(i)_{new} - y(i)_{old})^2}}{N} \quad (10)$$

5. Simulation results

We have implemented the deployment algorithms in MATLAB environment to demonstrate their performance. In our simulation, the sensor nodes are initially deployed at random over a 100 m by 100 m squared sensing field and grid scan method is used for evaluation of network coverage. The sensing field is treated as 100 by 100 grids when we calculate the coverage. In this paper we assume that the maximum velocity of each mobile sensor node is 0.5 m/s. For simulation, we set the maximum distance that a sensor node can move in each iteration as 0.5 m. The parameters used for simulation are given in Table 1.

5.1. Simulation using homogeneous sensors

The simulation results obtained using homogeneous mobile sensor nodes having sensing range 10 m and communication range 20 m is presented in this section. Here,

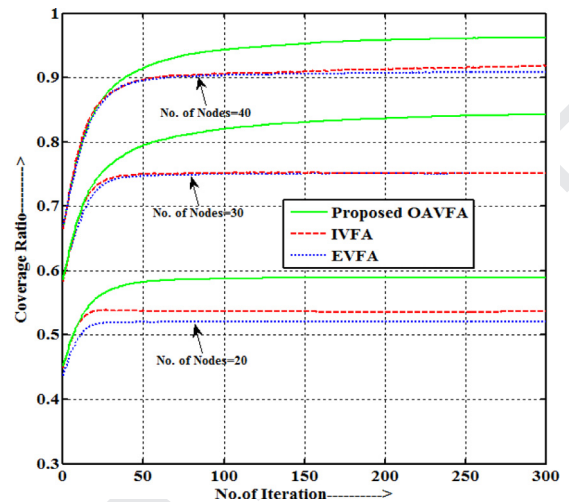


Fig. 2. Binary coverage ratio vs. no. of iterations.

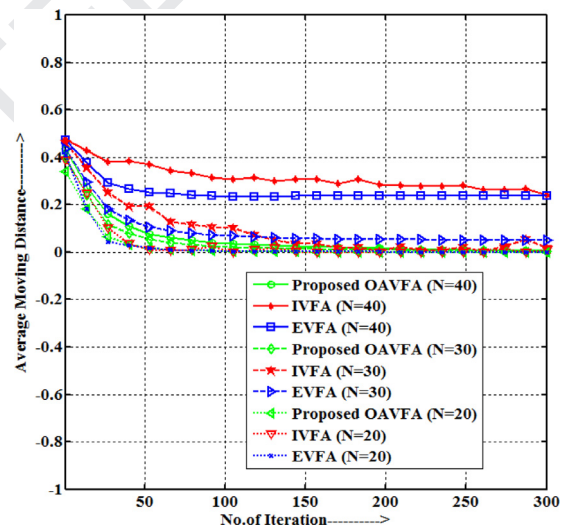


Fig. 3. Average moving distance vs. no. of iterations.

we use statistical methods to analyze the performance of deployment algorithms. In our simulation, 100 different random initial deployments are applied to each deployment algorithm. The parameters used for simulation are given in Tables 1 and 2. Fig. 2 shows the average final binary coverage ratio vs. iterations for IVFA [33], EVFA [33], and our proposed OAVFA without any obstacles when number sensors deployed in ROI is 20, 30 and 40.

From Fig. 2 it is clear that the coverage ratio due to OAVFA is higher than the other two approaches. Fig. 3 shows the average moving distance of sensor nodes in each step. The average moving energy consumption of virtual force algorithms decreases and ours converge very fast as compare to other two algorithms.

The performance of IVFA, EVFA and OAVFA on coverage rate and convergence time for three different network sizes with number of sensor nodes $N = 20, 30$ and 40 is given in Table 3. We observe that OAVFA attain a higher

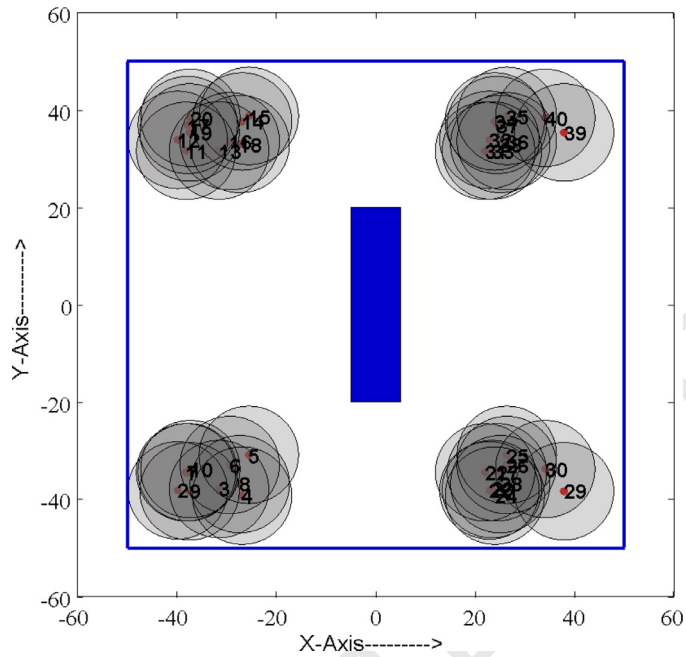


Fig. 4. Initial sensor position with coverage rate 31.61%.

Table 3

Performance summary.

Parameters	IVFA		EVFA		OAVFA				
N	40	30	20	40	30	20	40	30	20
Coverage Rate (%)	91.8	75.1	53.6	90.7	75.1	52	96.2	84.2	60
No. of iterations to achieve steady state	>300	>300	150	>300	>300	90	220	200	80

353 coverage rate compare to IVFA and EVFA for all three
 354 cases. For IVFA and EVFA, in the case of $N = 40$ and 30 ,
 355 some nodes are still subjected to repulsive or attractive
 356 force and move continuously even when the coverage rate
 357 remains constant. For OAVFA, the algorithm converge very
 358 well after 80, 200 and 220 iterations, respectively.

359 We also simulate OAVFA in presence of different shapes
 360 of obstacle at the central area of ROI. Initially, 40 ho-
 361 mogeneous mobile sensor nodes having sensing radius 10
 362 are split into four groups and randomly deployed at the
 363 four corners of the sensing field as shown in Fig. 4. From
 364 Figs. 5–11 illustrate the final sensor locations after execu-
 365 tion of proposed OAVFA.

366 From above results, it is clear that at the end of final
 367 deployment, no mobile sensor node is outside the ROI and
 368 the sensor nodes are self-deployed with avoidance of ob-
 369 stacle to cover the whole sensing field and also maintain
 370 the connectivity. The average final binary coverage rate vs.
 371 number of iterations with and without obstacles for 100
 372 different initial deployments is shown in Fig. 12. It appears
 373 that the coverage rate is as good as with and without the
 374 presence of an obstacle in ROI.

375 5.2. Simulation using heterogeneous sensors

376 The simulation results due to heterogeneous mobile
 377 sensors are presented in this section. Initially, the het-

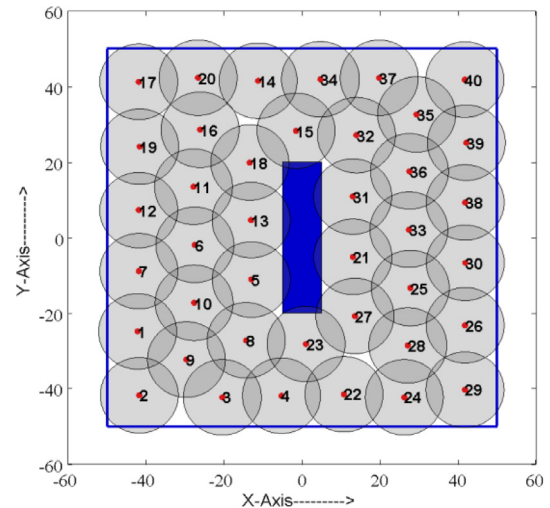


Fig. 5. Final deployment with coverage rate 97.75%.

erogeneous mobile sensors are randomly deployed in a 378
 100 m by 100 m squared sensing field containing different 379
 shape of obstacles. We set maximum sensing range limit 380
 is 10 and communication range of each sensor node is 20 381
 (i.e. $C_R = 2 \times$ maximum sensing range limit) and minimum 382

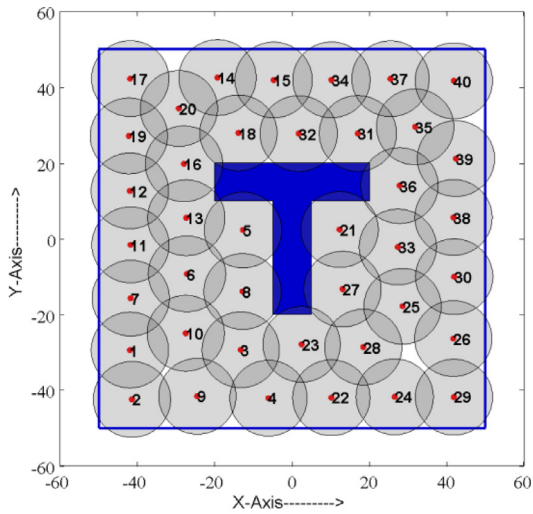


Fig. 6. Final deployment with coverage rate 98.29%.

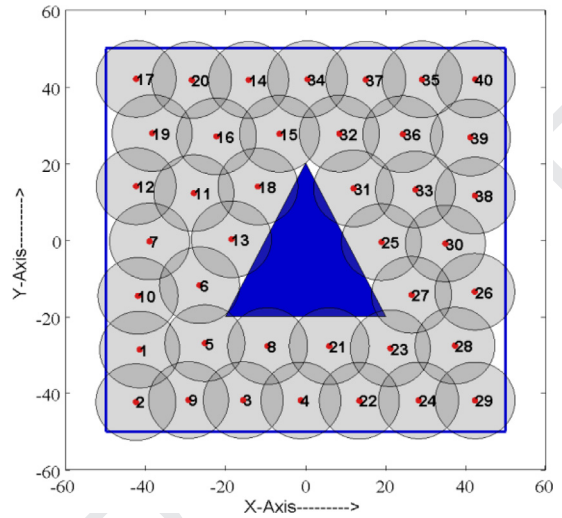


Fig. 9. Final deployment with coverage rate 97.86%.

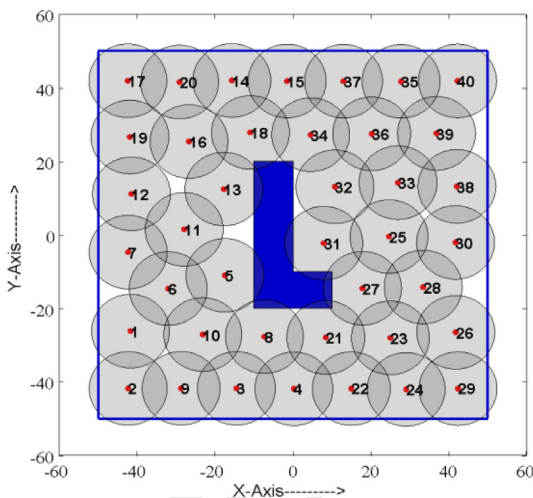


Fig. 7. Final deployment with coverage rate 97.19%.

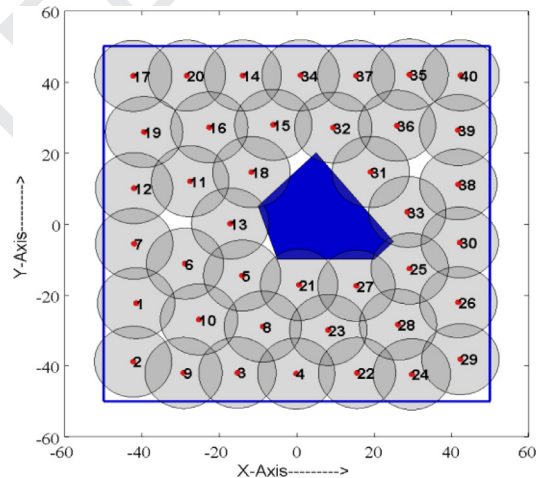


Fig. 10. Final deployment with coverage rate 97.52%.

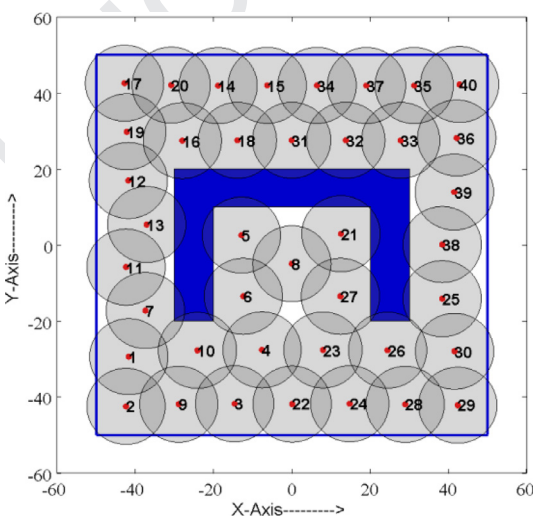


Fig. 8. Final deployment with coverage rate 97.09%.

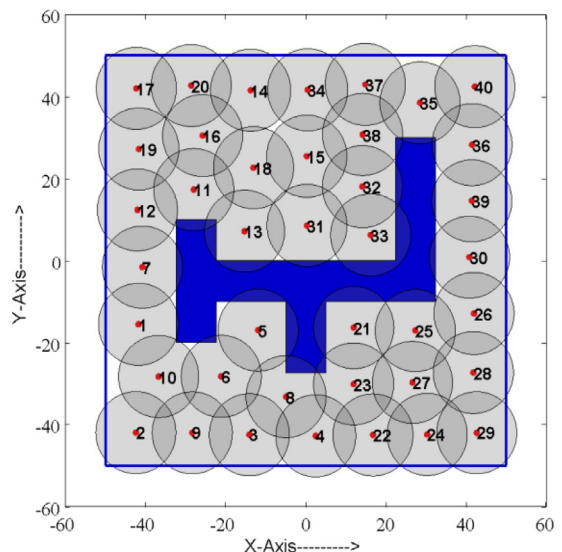


Fig. 11. Final deployment with coverage rate 98.34%.

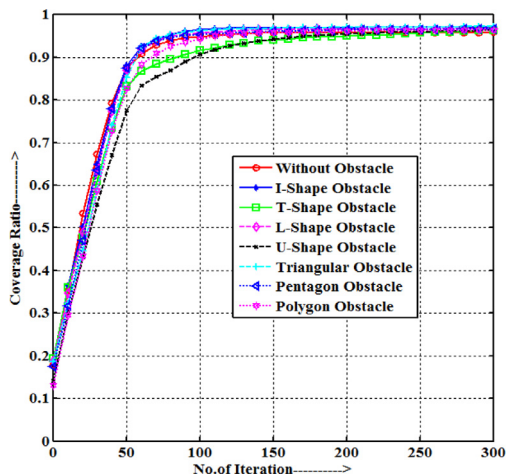


Fig. 12. Coverage rate with and without obstacles.

Table 4

Simulation parameters for heterogeneous sensor.

Parameters	Value
Sensing Range limit (R_{si})	6 m-10 m
Communication Range ($C_R = 2 \times \max(R_{si})$)	20 m
Threshold distance (d_{th}^{ij})	$\frac{\sqrt{3}}{2}(R_{si} + R_{sj})$
$d_{th_obs}(S_i)$	$\sqrt{3}R_{si}/2$
$d_{th_b}(S_i)$	$\sqrt{3}R_{si}/2$

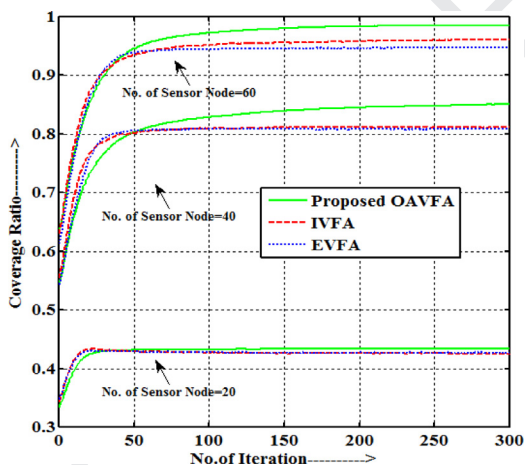


Fig. 13. Binary coverage ratio vs. no. of iterations.

sensing range limit is 6. The parameters used for simulation are given in Tables 1 and 4. Fig. 13 shows the average final binary coverage ratio vs. iterations for IVFA [33], EVFA [33], and our proposed OAVFA without any obstacles when number sensor deployed in ROI is 60, 40 and 20.

Fig. 13, indicates that, our proposed algorithm has better coverage than that of other two VFA approaches. The average moving distance of sensor nodes in each step is shown in Fig. 14. The average moving distance decreases for all three deployment algorithms, but ours converge faster. The performance of IVFA, EVFA and OAVFA on coverage rate for three different network sizes with number

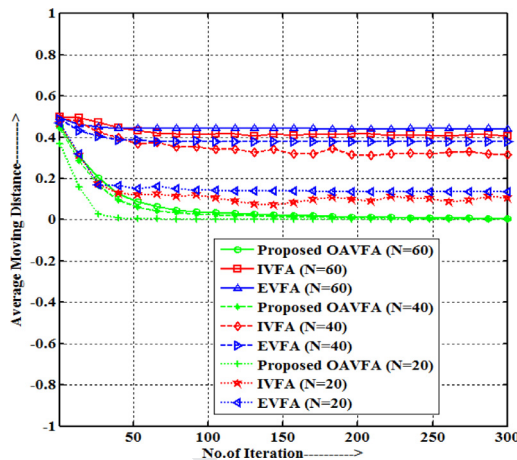


Fig. 14. Average moving distance vs. no. of iteration.

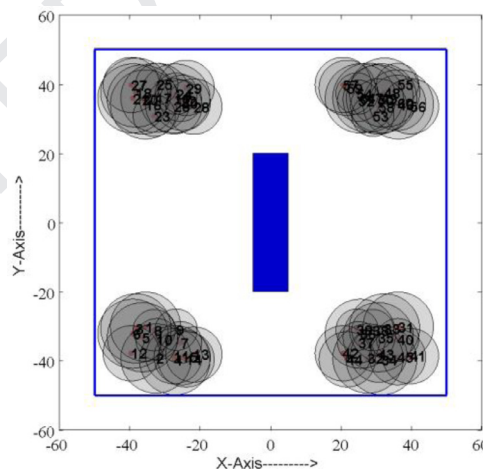


Fig. 15. Initial deployment with coverage rate 29.94%.

of sensor nodes $N = 20, 40$ and 60 is given in Table 5. We observe that OAVFA attain a higher coverage rate compare to IVFA and EVFA for all three cases. For IVFA and EVFA, in the case of $N = 20, 40$, and 60 some nodes are still subjected to repulsive or attractive force and move continuously even when the coverage rate remains constant. For OAVFA, the algorithm converge very well after 70, 220 and 250 iterations, respectively

To demonstrate the performance of proposed OAVFA with different obstacle shapes we have simulated our approach in a 100 m by 100 m squared sensing field cover with a clear boundary with different obstacle shapes at the central area of ROI. Fig. 15 illustrates the initial random deployment of 60 heterogeneous sensor nodes at four corners of ROI in the presence of I-shape obstacle and Figs. 16-22 illustrate the final position of sensor nodes after 300 iterations in the presence of different shape of obstacles in ROI.

From above results, it is clear that at the end of final deployment, no mobile sensor node lies outside the ROI. The sensor nodes are self-deployed with avoidance of obstacle to cover the whole sensing field and the connectivity

Table 5
Performance summary.

Parameters	IVFA			EVFA			OAVFA		
N	60	40	20	60	40	20	60	40	20
Coverage Rate (%)	96.1	81.1	42.6	94.7	80.9	43	98.5	85	43.6

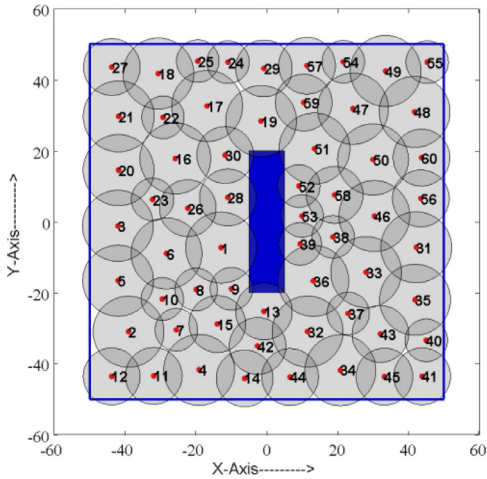


Fig. 16. Final deployment with coverage rate 99.14%.

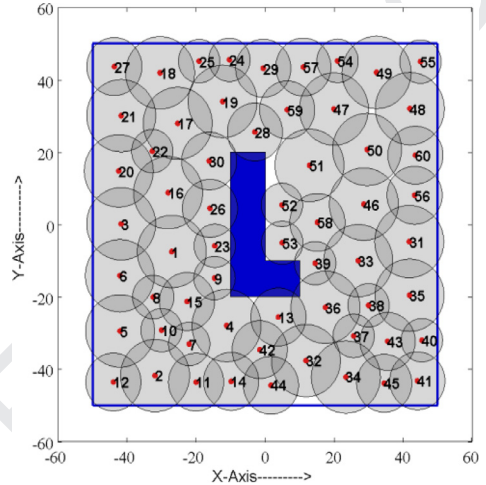


Fig. 18. Final deployment with coverage rate 98.92%.

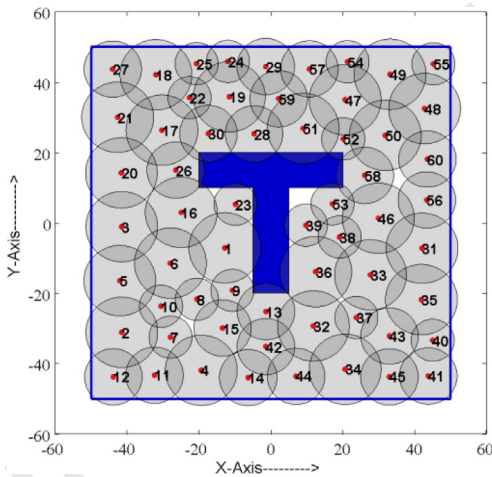


Fig. 17. Final deployment with coverage rate 98.7%.

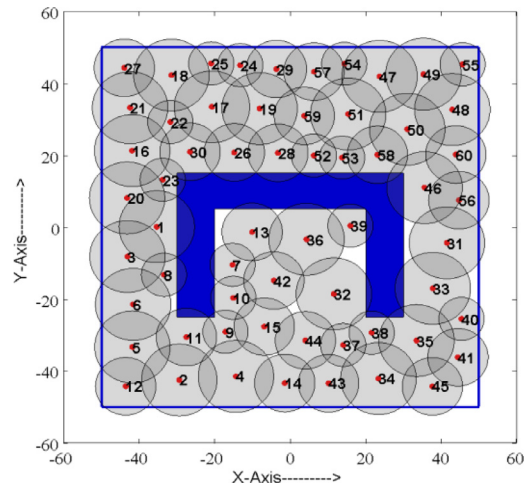


Fig. 19. Final deployment with coverage rate 97.98%.

is also maintained. We studied the impact of the shape of the obstacles on the coverage rate. The average final binary coverage rate vs. number of iterations with and without obstacles for 100 different initial deployments is shown in Fig. 23. It appears that the coverage rate is as good as with and without the presence of an obstacle in ROI

6. Conclusion

In this paper, we propose a localized self-deployment scheme called OAVFA for homogeneous as well as heterogeneous mobile sensor networks with random initial distribution. This algorithm works well in the scenarios of the random initial distribution of mobile sensor nodes to max-

imize the area coverage and minimize the moving energy requirement in the presence of obstacles while maintaining connectivity. To prevent the sensor nodes from moving out of sensing field boundary, we consider a repulsive force exerted by sensing field boundary. We also add repulsive force exerted by obstacles to avoid the presence of obstacles in ROI. Simulation results demonstrate that the proposed approach provides better performance than IVFA and EVFA for deployment of homogeneous as well as heterogeneous sensor nodes in a squared sensing field with and without the obstacles.

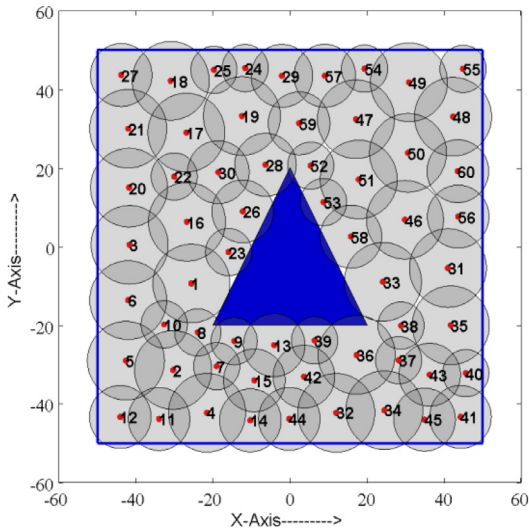


Fig. 20. Final deployment with coverage rate 99.38%.

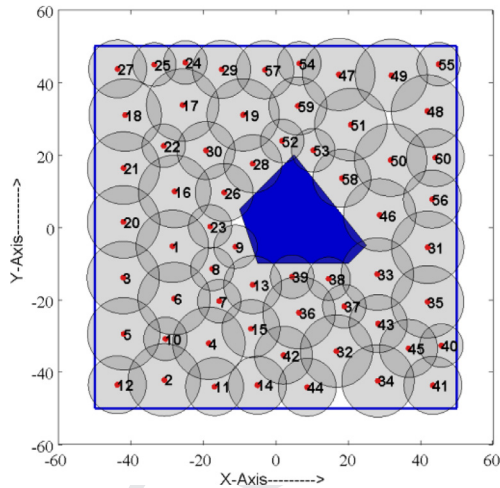


Fig. 21. Final deployment with coverage rate 99.09%.

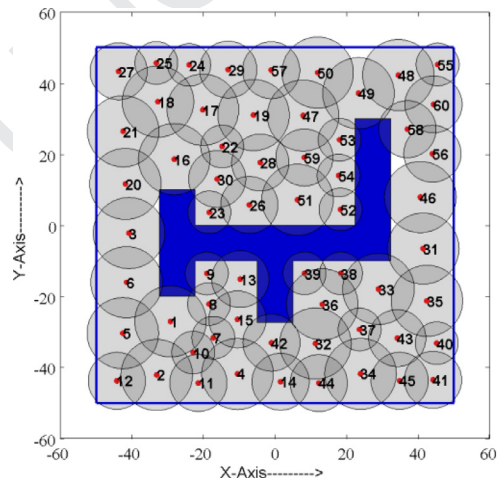


Fig. 22. Final deployment with coverage rate 98.96%.

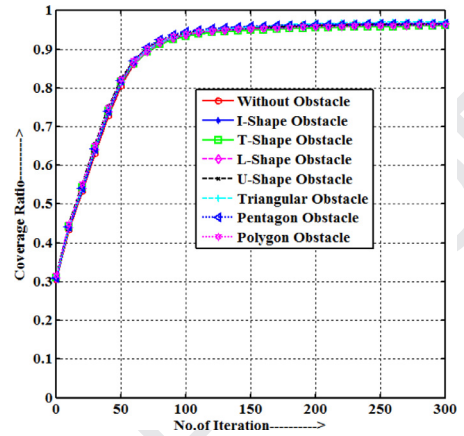


Fig. 23. Binary coverage with and without obstacles.

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